

Theory of Computation, CSCI 438 spring 2022
Multitape Turing Machines, pg. 176-178, March 27th

2-tape Turing Machine: Like an ordinary Turing Machine but has two tapes, each with its own read/write head. The input is on tape 1 followed by blanks and tape 2 is completely blank. The read/write heads are positioned in the first positions of each tape. For flexibility, say that the read/write heads can move separately, that is allow stays.

$$M=(Q,\Sigma, \Gamma_1,\Gamma_2,\delta,q_0,q_{\text{accept}}, q_{\text{reject}})$$

$$\delta:Q \times \Gamma_1 \times \Gamma_2 \rightarrow Q \times \Gamma_1 \times \Gamma_2 \times \{L,R,S\} \times \{L,R,S\}$$

Multitape Turing Machine: Like an ordinary Turing Machine but has k tapes so k read/write heads. Say initially, input on tape 1 is as usual and the other tapes are blank. Formally define the machine:

$$M=(Q,\Sigma,\Gamma,\delta,q_0,q_{\text{accept}}, q_{\text{reject}})$$

$$\delta:Q \times \Gamma^k \rightarrow Q \times \Gamma^k \times \{L,R,S\}^k$$

$$\delta:Q \times \underbrace{\Gamma \times \Gamma \dots \times \Gamma}_{k \text{ times}} \rightarrow Q \times \underbrace{\Gamma \times \Gamma \dots \times \Gamma}_{k \text{ times}}, \underbrace{\{L,R,S\} \times \{L,R,S\} \dots \times \{L,R,S\}}_{k \text{ times}}$$