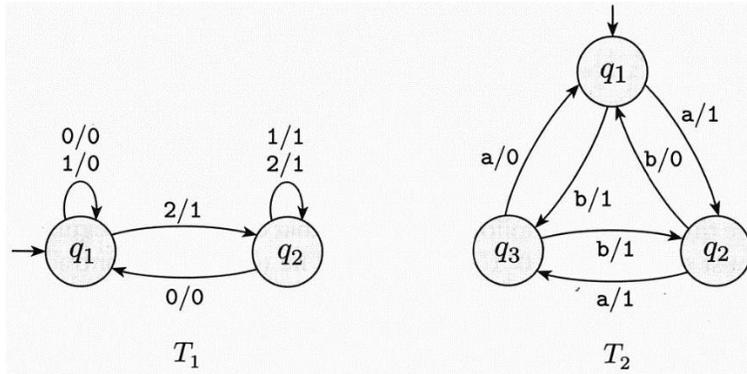


**Theory of Computation, CSCI 438 spring 2022**  
**Definition of Regular languages, pg. 40-44, Jan. 14**

**Exercise 1.24 a, b, f & g, 1.25 & 1.26 (page 87)**

1.24 Given the two transducers  $T_1$  and  $T_2$



Give the sequence of states entered and the output produced in each of the following parts.

a.  $T_1$  on input 001 Enters the sequence of states  $q_1, q_1, q_1, q_1$  and outputs 000

b.  $T_1$  on input 211 Enters the sequence of states  $q_1, q_2, q_2, q_2$  and outputs 111

f.  $T_2$  on input bbab Enters the sequence of states  $q_1, q_3, q_2, q_3, q_2$  and outputs 1111

g.  $T_2$  on input  $\epsilon$  Enters the sequence of states  $q_1$  and outputs nothing

1.25 Read the informal definition of the finite state transducer given in Exercise 1.24. Give a formal definition of this model, following the pattern in Definition 1.5 (page 35). Assume that an FST has an input alphabet  $\Sigma$  and an output alphabet  $\Gamma$  but not a set of accept states.

A finite state transducer, FST, is a 5-tuple  $(Q, \Sigma, \delta, \Gamma, q_0)$  where

1.  $Q$  is a finite set of states
2.  $\Sigma$  is a finite set called the input alphabet
3.  $\Gamma$  is a finite set called the output alphabet
4.  $\delta: Q \times \Sigma \rightarrow Q \times \Gamma$  is the transition function
5.  $q_0 \in Q$  is the start state

Include a formal definition of the computation of an FST.

Computation of a FST begins in state  $q_0$ , reading symbols from  $\Sigma^*$  and outputting symbols from  $\Gamma^*$ . For each symbol read from the input, the FST applies  $\delta$  to the current state and symbol, moving to the new state, and outputting the indicated symbol from  $\Gamma$ . The FST quits when all of the input has been consumed.

Alternatively, describing the computation more formally:

Given a finite state transducer, FST,  $T = (Q, \Sigma, \delta, \Gamma, q_0)$ , where

$\delta^*$  is the application of  $\delta$  multiple times and has the signature

$$\delta^*: Q \times \Sigma^* \rightarrow Q \times \Gamma^*$$

On input  $w$ ,  $w \in \Sigma^*$ , the FST outputs  $o \in \Gamma^*$ , when

$$\delta^*(q_0, w) = (q_x, o)$$

1.26 Using the solution you gave to Exercise 1.25, given a formal description of the machines T1 and T2 depicted in Exercise 1.24.

T<sub>1</sub> is  $M = (\{q_1, q_2\}, \{0, 1, 2\}, \{0, 1\}, \delta, q_1)$  where  $\delta$  is defined as:

$$\delta(q_1, 0) \rightarrow q_1, 0$$

$$\delta(q_1, 1) \rightarrow q_1, 0$$

$$\delta(q_1, 2) \rightarrow q_2, 1$$

$$\delta(q_2, 0) \rightarrow q_1, 0$$

$$\delta(q_2, 1) \rightarrow q_2, 1$$

$$\delta(q_2, 2) \rightarrow q_2, 1$$

T<sub>2</sub> is  $M = (\{q_1, q_2, q_3\}, \{a, b\}, \{0, 1\}, \delta, q_1)$  where  $\delta$  is defined as:

$$\delta(q_1, a) \rightarrow q_2, 1$$

$$\delta(q_1, b) \rightarrow q_3, 1$$

$$\delta(q_2, a) \rightarrow q_3, 1$$

$$\delta(q_2, b) \rightarrow q_1, 0$$

$$\delta(q_3, a) \rightarrow q_1, 0$$

$$\delta(q_3, b) \rightarrow q_2, 1$$