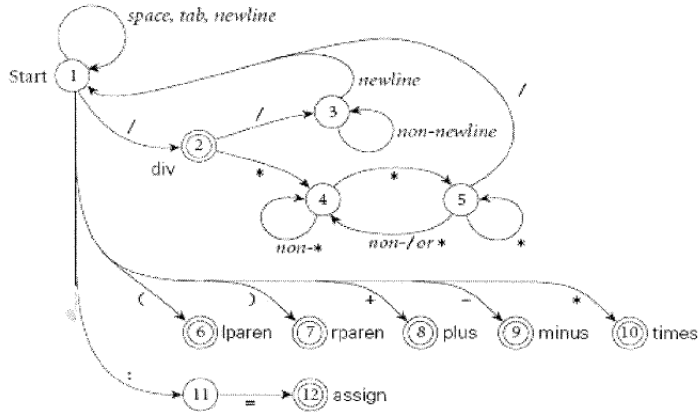


**Concepts of Programming Languages, CSCI 305, Fall 2021
Homework #4, complete by Oct. 8**

1. Create a scanning table for the following minimal dfa.



State	space	tab	newline	/	*	()	+	-	:	=
1	1	1	1	2	10	6	7	8	9	11	-
2	-	-	-	3	4	-	-	-	-	-	-
3	3	3	1	3	3	3	3	3	3	3	3
4	4	4	4	4	5	4	4	4	4	4	4
5	4	4	4	1	5	4	4	4	4	4	4
6	-	-	-	-	-	-	-	-	-	-	-
7	-	-	-	-	-	-	-	-	-	-	-
8	-	-	-	-	-	-	-	-	-	-	-
9	-	-	-	-	-	-	-	-	-	-	-
10	-	-	-	-	-	-	-	-	-	-	-
11	-	-	-	-	-	-	-	-	-	-	12
12	-	-	-	-	-	-	-	-	-	-	-

2. Hand execute the table-driven lexical analyzer pseudo code (Figure 2.11, page 66) on the very simple program.

```
sum :=29.57
write sum+4
```

For each call to the lexical analyzer, list all accesses into the scan_tab, in order as the code is executed. End the call by showing what is returned to the parser.

Each “move” action requires two scan_tab accesses, one to get the action and a second one to get the new_state. Only one access needs to be listed.

First call to the lexical analyzer:

```
scan_tab(s, 1)
scan_tab(u, 16)
scan_tab(m, 16)
scan_tab(space, 16)
return <id, sum>
```

Next call to the lexical analyzer:

```
scan_tab(:, 1)
scan_tab(=, 11)
scan_tab(2, 12)
return <assign, :=>
```

Next call to the lexical analyzer:

```
scan_tab(2, 1)
scan_tab(9, 14)
scan_tab(., 14)
scan_tab(15, 5)
scan_tab(7, 15)
scan_tab(newline, 15)
return <number, 29.57>
```

Next call to the lexical analyzer:

```
scan_tab(newline, 1)
scan_tab(w, 17)
scan_tab(w, 1)
scan_tab(r, 16)
scan_tab(i, 16)
scan_tab(t, 16)
scan_tab(e, 16)
```

```
scan_tab(space, 16)
return <write, write>
```

Next call to the lexical analyzer:

```
scan_tab(space, 1)
scan_tab(s, 17)
scan_tab(s, 1)
scan_tab(u, 16)
scan_tab(m, 16)
scan_tab(+, 16)
return <id, sum>
```

Next call to the lexical analyzer:

```
scan_tab(+, 1)
scan_tab(4, 8)
return <plus, +>
```

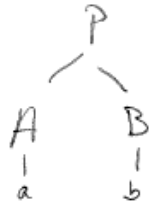
Next call to the lexical analyzer:

```
scan_tab(4, 1)
scan_tab(newline, 14)
return <number, 4>
```

Next call to the lexical analyzer:

```
scan_tab(newline, 1)
EOF when in stat
```

3. Consider the parse tree below.



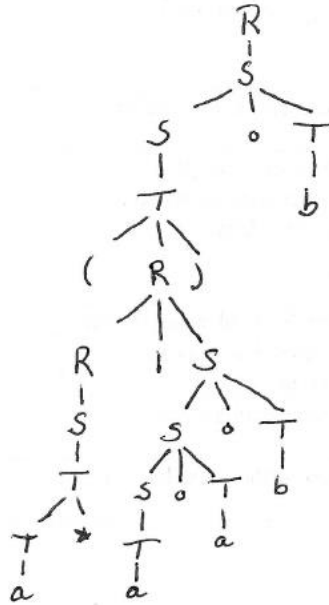
Give a rightmost derivation of the tree:

$$P \Rightarrow AB \Rightarrow Ab \Rightarrow ab$$

Give a leftmost derivation of the tree:

$$P \Rightarrow AB \Rightarrow aB \Rightarrow ab$$

4. Consider the parse tree below.



Give a rightmost derivation of the tree:

$$\begin{aligned}
 R &\Rightarrow S \Rightarrow S^\circ T \Rightarrow S^\circ b \Rightarrow T^\circ b \Rightarrow (R)^\circ b \Rightarrow (R|S)^\circ b \Rightarrow (R|S^\circ T)^\circ b \Rightarrow \\
 &(R|S^\circ b)^\circ b \Rightarrow (R|S^\circ T^\circ b)^\circ b \Rightarrow (R|S^\circ a^\circ b)^\circ b \Rightarrow (R|T^\circ a^\circ b)^\circ b \Rightarrow \\
 &(R|a^\circ a^\circ b)^\circ b \Rightarrow (S|a^\circ a^\circ b)^\circ b \Rightarrow (T|a^\circ a^\circ b)^\circ b \Rightarrow (T^*|a^\circ a^\circ b)^\circ b \Rightarrow \\
 &(a^*|a^\circ a^\circ b)^\circ b \Rightarrow
 \end{aligned}$$

5. Create a grammar for the context-free language

$$L = \{x_1 w^R x_2 \# w \mid x_1, x_2, w \in \{a, b\}^*\}$$

where w^R is the reverse of w

Strings in L	Strings not in L
#	ϵ
aaabaa#b	a
bbaaaa#abb	b
bbabbbb#bbab	bb#aa
	ba#ba
	aaaaba#bab

$$S \rightarrow X P$$

$$X \rightarrow aX \mid bX \mid \epsilon$$

$$P \rightarrow aPa \mid bPb \mid X\#$$