Chapter 11

## Scheme

## Scheme - Functional Language

- Historical context
- Problem area addresses
- Data structures & how implemented
- Built-in items & libraries
- Sample code
- Translation process
- New features
- Strengths & weaknesses
- Conclusion

# Historical Context – Functional Languages

- 1940's computers programed using machine language
- FORTRAN, first high-level programming language (1954)
- Lisp, first functional language (1958)
- Also in 1958
  - Flow-Matic precursor to COBOL
  - ALGOL standard method of describing algorithms, influencer of all imperative languages

### Historical Context - Scheme

- Scheme general semantics & syntax came from Lisp, lexical (static) scoping and block structure came from ALGOL
- Developed from MIT AI Memos known as the Lambda Papers (1975-1980)
- Versions in 1975, 1978, 1984, standardized in 1990, revised 1998, 2006, 2007
- o Influenced R, Rust and Swift

### Problem Area Addressed

LISP follows the mathematical paradigm of algorithms much more than Fortran

Fortran defined for number crunching LISP for symbolic manipulation

## Why use a Functional Language

Functional languages are useful for:

- o game Al portion nice list manipulations
- mathematical computations easy to read
- o concurrency since no state
- want to prove correctness

## Problem Area Addressed by Scheme

Contrary to Lisp, Scheme is functional yet:

- Uses lexical scoping
- Includes imperative features

## Scheme Data Structures

Scheme (and Lisp) only have two data structures

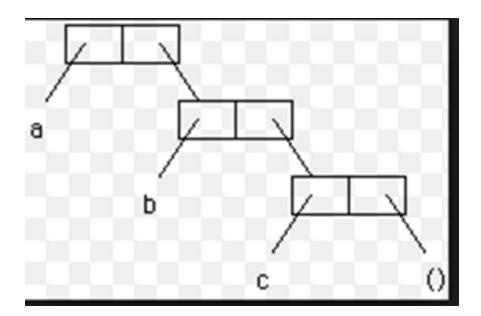
- Atom symbol
- Lists

Atoms: X, hello, #t, #f, 0.566, ½, 3.5 + 4i

Lists: (A B C D) , (A (B C) D (E (F G)))

## Internal Implementation

Data structure: (a b c)



## Major Built-In Procedures: car and cdr

car – takes a list are returns the first element in the list

$$(car'(123)) => 1$$

cdr – takes a list and returns the rest of the list

$$(cdr'(123)) => '(23)$$

#### **Built-In Procedures**

**cons** – takes an element and a list and inserts the element into the first position of the list

**list** – takes any number of elements and constructs a list from them

append – takes two lists and combines them into a single list

length – takes a list and returns it's length

## **Built-In Procedures**

**define** – takes a name and a literal and binds the name to the literal

lambda – takes a list of operands and a function definition and optional a list of values, and creates a function

```
Example:

(define x '(a b c)) ; x is a list

(define x (lambda ; x is a procedure

(n1 n2) (+ n1 n2)

)
```

### **Built-In Predicates**

Predicate a procedure that returns #t or #f. By convention, predicates end with '?'.

```
null?
list?
boolean?
number?
equal?
procedure?
zero? - expects a number
even? - expects an integer
odd? - expects an integer
=, <, <=, >, >= - expect numbers
```

## Scheme Data Libraries

#### Racket libraries:

- Draw basic drawing tools, including drawing contexts such as bitmaps and PostScript files.
- Gui GUI widgets such as windows, buttons, checkboxes, and text fields. The library also includes a sophisticated and extensible text editor.
- Pict functional abstraction layer over draw, useful for creating slide presentations and images

## Sample Code

Take a start and end and return a list of the numbers between, inclusive.

```
Example (myCount 15) returns (12345)
```

## Scheme Translation

Scheme is interpreted

Scheme interpreter cycles continuously through

- Read
- Evaluate
- Print

# New Features of Functional Programs

Ability to create functions "on the fly" – lambda

Higher order functions

## Higher Order Functions

Higher order functions are functions that takes functions as arguments

Example: map

(define square (lambda (x) (\* x x)))

(map square '(2 3 4)) returns '(4 9 16)

## Create Functions on-the-fly

```
From before:
(define square (lambda (x) (* x x)))
(map square '(2 3 4)) returns '(4 9 16)
```

```
Equivalently,
(map (lambda (x) (* x x))) '(2 3 4))
```

# Strengths of Functional Languages

Without understanding functional programming, you can't invent MapReduce, the algorithm that makes Google so massively scalable. The terms Map and Reduce come from Lisp and functional programming.

Purely functional programs have no side effects and are thus trivially parallelizable which is good for concurrency

# Weaknesses of Functional Languages

Performance - Using only immutable values (no side-effects) and recursion can lead to performance problems – high RAM use and speed

Debugging - Writing pure functions is easy, but when combining them, debugging gets hard

Pure functions and I/O don't mix

For some, recursion isn't natural

## Conclusion – Why Study Scheme?

- Writing programs in Scheme helps you think recursively
- Recursion is a powerful problem solving skill
- Good for demonstrating certain language features such as higher order functions

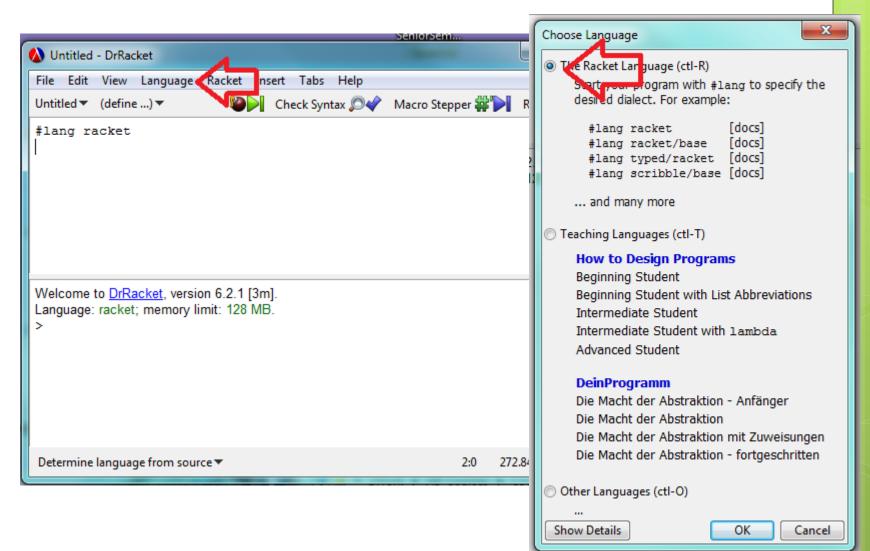
## Conclusion

Experience programming in Scheme will make you a stronger programmer and more knowledgeable about programming languages in general.

## Scheme via Dr. Racket on Windows



## Dr. Racket



## Scheme on katie

On kaite: mzscheme located at /usr/bin/mzscheme.

```
To invoke: mzscheme
Welcome to Racket v5.2.1.
>
(load "filename.mz")
> (exit)
```

### Read-Evaluate-Print

Scheme interpreter cycles continuously through

- Read
- Evaluate
- Print

Sometimes called REPL for Read-Evaluate-Print Loop (pronounced REP-ple)

Items are automatically evaluated.

Single quote - "don't evaluate."

## Basic Types

Literals (numbers, Booleans, strings, characters) evaluate to themselves. For example:

## Scheme – Don't Evaluate

- Single quote lists since don't want evaluation to take place
- Once a list is quoted, don't quote elements within it

```
> (123) => application: not a procedure; expected procedure, given: 1
```

#### **Built-In Procedures**

if – takes a predicate, then-expression, and else-expression and evaluates one of the expressions based on the predicate

when – takes a predicate and expression to eval

**unless** – takes a predicate and expression to evaluate when the predicate is false

#### **Built-In Procedures**

cond – takes any number of pairs where each pair consists of a predicate followed by an expression. It searches down the list of pairs, and evaluates the expression associated with the first predicate that is true. The final predicate can be the word else.

## Built-In Procedure - map

map - takes an function and a list. It applies the function to each element of the list and returns the resulting list.

#### Built-In Procedure - member

**member** (or **memq**) - takes an element and a list. It searches the list and if it finds the element, it returns the tail of the list beginning with that element. Otherwise it returns #f

```
(member 'a '(1 b a d 3)) => (a d 3)
(member 'a '(1 b d 3)) => #f
```

#### Built-In Procedure - assoc

**assoc** - takes an element and a list of pairs. It looks through the list and returns the first pair for which the first element matches

```
(assoc 'a '((c 4) (b 3) (a 5) (d 6)) => (a 5)
(assoc '(q2 0) '((q2 4) q3) ((q3 1) q2) ((q2 0) q5)) =>
((q2 0) q5)
(assoc 'a '((c 4) (b 3)) => #f
```

## Equality

= only for numbers.equal? for any itemeq? most discriminating

Example:
(define a '(1 2 3))
(equal? a '(1 2 3)) => #t
(eq? a '(1 2 3)) => #f

### Built-In Procedure - let

**let** - takes two arguments, a list of pairs and an expression. Each pair in the first list consists of an element and an expression. The 2<sup>nd</sup> argument of let is evaluated using the constant assignments given in the first argument. In other words, the variables are bound by the let for the duration of the expression.

```
Example:

(let ( ; first argument

(a 10)

(b 15)

)

(+ a b) ; second argument
```

#### Built-In Procedure - let

let - the initial values are computed and results are bound to variables

**let\*** - from left to right, compute initial value and bind to variable, so earlier bindings can be used in later bindings

```
Example:
```

## I/O Scheme Functions

I/O Scheme functions:

read – read one Scheme object from the standard input and return it (thus, this is a function with a side effect, a new constant exists in the environment)

(define x (read))

write – print a representation of its argument to standard output so data could be read back in

(write (+ 3 5)) => 8

display – similar to write but datatypes are written as raw bytes or characters (not meant to be written back in)

### Scheme Function read

```
read - read one Scheme object
     Example:
(define x (read))
x
```

Welcome to <u>DrRacket</u>, version 6.2.1 [3m]. Language: racket; memory limit: 128 MB.

hello

eof

Returns 'hello

## Scheme Function read

```
read - read one Scheme object
     Example:
(define x (read))
x
```

Welcome to <u>DrRacket</u>, version 6.2.1 [3m]. Language: racket; memory limit: 128 MB.

(a b c)

eof

Returns '(abc)

#### Scheme Function read-char

```
read-char - read a single character
    Example:
(define x (read-char))
x
```

Welcome to <u>DrRacket</u>, version 6.2.1 [3m]. Language: racket; memory limit: 128 MB.

(a b c)

eof

Returns #\(

## Scheme Function read-line

```
read-line – read up to end of file
Example:
(define x (read-line))
x
```

Welcome to <u>DrRacket</u>, version 6.2.1 [3m]. Language: racket; memory limit: 128 MB.

Hello this is a whole line

eof

Returns "Hello this is a whole line"

## I/O Scheme File Functions

I/O Scheme functions:

open-input-file – as expected

Example: (open-input-file "table.csv")

Example:

(read-line (open-input-file "table.csv"))

eof-object? - as expected

Example:

(if (eof-object? (read-line

(open-input-file "table.csv")))

### **Ports**

Have the usual ports – input, output and error

open-input-file – takes a string pathName, opens the file for input and returns an input-port, or triggers exn:fail:filesystem exception

open-output-file – similarly, but will create a file, so it can't already exist

close-input-port close-output-port custodian-shutdown-all