

## Overview

- Recursion
- A method calling itself
- A new way of thinking about a problem
- A powerful programming paradigm
- Examples:
- Last time:
- Factorial, binary search, H-tree, Fibonacci
- Today:
- Greatest Common Divisor (GCD)
- Brownian Motion
- Sorting things


## Greatest Common Divisor

- GCD
- Find largest integer $d$ that evenly divides $p$ and $q$
- e.g. $\operatorname{gcd}(4032,1272)=24$
- $4032=2^{6} \times 3^{2} \times 7^{1}$
- $1272=2^{3} \times 3^{1} \times 53^{1}$
- $\operatorname{gcd}=2^{3} \times 3^{1}=24$
- Applications
- Simplify fractions: 1272/4032 $=53 / 168$
- RSA cryptography


## Greatest Common Divisor

- GCD
- Find largest integer $d$ that evenly divides $p$ and $q$
- Euclid's algorithm (300 BC)


```
gcd}(p,q)={\begin{array}{ll}{p}&{\mathrm{ if }q=0}\\{\operatorname{gcd}(q,p%q)}&{\mathrm{ otherwise }}\\{\mathrm{ & reduction step,}}
``` converges to base case
\(\begin{aligned} \operatorname{gcd}(4032,1272) & =\operatorname{gcd}(1272,216) \\ & =\operatorname{gcd}(216,192)\end{aligned}\) \(=\operatorname{gcd}(192,24)\)
\(=\operatorname{gcd}(24,0)\)
\(=24\)

\section*{Greatest Common Divisor}
- GCD
- Find largest integer \(d\) that evenly divides \(p\) and \(q\)


\section*{Brownian motion}
- Physical process that models many natural and artificial phenomenon
- Price of stocks
- Rugged shapes of mountains and clouds
- Fractal landscape and textures for computer graphics


\section*{Greatest Common Divisor}
- GCD
- Find largest integer \(d\) that evenly divides \(p\) and \(q\)
\[
\operatorname{gcd}(p, q)=\left\{\begin{array}{ll}
p & \text { if } q=0 \\
\operatorname{gcd}(q, p \% q) & \text { otherwise }
\end{array}\right. \text { ๘ base case }
\]
converges to base case
```

public static int gcd(int p, int q)
if (q == 0)
return p
else
return gcd(q, p % q);

```

\section*{Simulating Brownian Motion}
- Midpoint displacement method
- Track interval \(\left(\mathrm{x}_{0}, \mathrm{y}_{0}\right)\) to \(\left(\mathrm{x}_{1}, \mathrm{y}_{1}\right)\)
- Choose \(\delta\) randomly from Gaussian distribution
- Divide in half, \(x_{m}=\left(x_{0}+x_{1}\right) / 2\) and \(y_{m}=\left(y_{0}+y_{1}\right) / 2+\delta\)
- Recur on the left and right intervals


\section*{Simulating Brownian Motion}

\section*{Plasma cloud}
- Same idea, but in 2D
- Each corner of square has some greyscale value
- Divide into four sub-squares
- New corners: avg of original corners, or all 4 + random
- Recur on four sub-squares




\section*{Divide and conquer}
- Divide and conquer paradigm
- Break big problem into small sub-problems
- Solve sub-problems recursively
- Combine results
"Divide et impera. Vendi, vidi, vici."
-Julius Caesar
- Used to solve many important problems
- Mergesort, sorting things, \(\mathrm{O}(\mathrm{N} \log \mathrm{N})\)
- Parsing programming languages
- Discrete FFT, signal processing
- Multiplying large numbers
- Traversing multiply linked structures (stay tuned)


\section*{Divide and conquer: sorting}
- Goal: Sort by number, ignore suit, aces high


\section*{Approach}
1) Split in half (or as close as possible)
2) Give each half to somebody to sort
3) Take two halves and merge together


Unsorted pile \#1


Unsorted pile \#2

\section*{Approach}
1) Split in half (or as close as possible)
2) Give each half to somebody to sort
3) Take two halves and merge together


Sorted pile \#1


Sorted pile \#2

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Sorted pile \#2
How many operations to do the merge? Linear in the number of cards, \(\mathrm{O}(\mathrm{N})\) But how did pile 1 and 2 get sorted?

\section*{Recursively of course!}

Split each pile into two halves, give to different people to sort.


\section*{Summary}
- Recursion
- A method calling itself:

- Sometimes just once, e.g. binary search
- Sometimes twice, e.g. mergesort
- Sometimes multiple times, e.g. H-tree
- All good recursion must come to an end
- Base case that does NOT call itself recursively
- A powerful tool in computer science
- Allows elegant and easy to understand algorithms
- (Once you get your head around it)```

