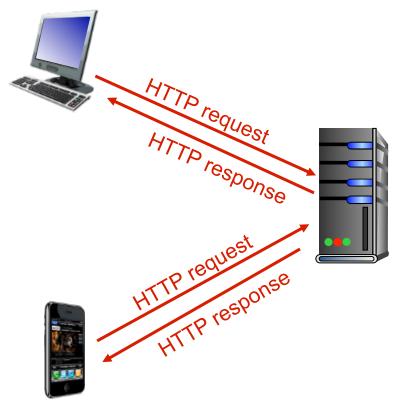
# HTTP, cookies, and caching

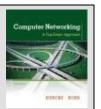




Computer Networking: A Top Down Approach 6th edition

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### Overview

- Chapter 2: Application Layer
  - Many familiar services operate here
    - Web, email, Skype, P2P file sharing
  - Socket programming

#### HTTP

- Statelessness
- Non-persistent vs. persistent connections
- Cookies
- Caching

application
transport
network
link
physical

### HTTP: TCP + stateless

#### **Uses TCP:**

- Client initiates TCP connection (creates socket) to server, port 80
- Server accepts TCP connection from client
- HTTP messages

   (application-layer messages) exchanged between browser (HTTP client) and web server
   (HTTP server)
- TCP connection closed

#### HTTP is "stateless"

 Server maintains no information about past client requests

# Protocols that maintain "state" are complex!

- Past history (state) must be maintained
- If server/client crashes, their views of "state" may be inconsistent, must be reconciled

# Trying out HTTP for yourself

1. Telnet to your favorite Web server:

telnet cis.poly.edu 80

opens TCP connection to port 80 (default HTTP server port) at cis.poly.edu. anything typed in sent to port 80 at cis.poly.edu

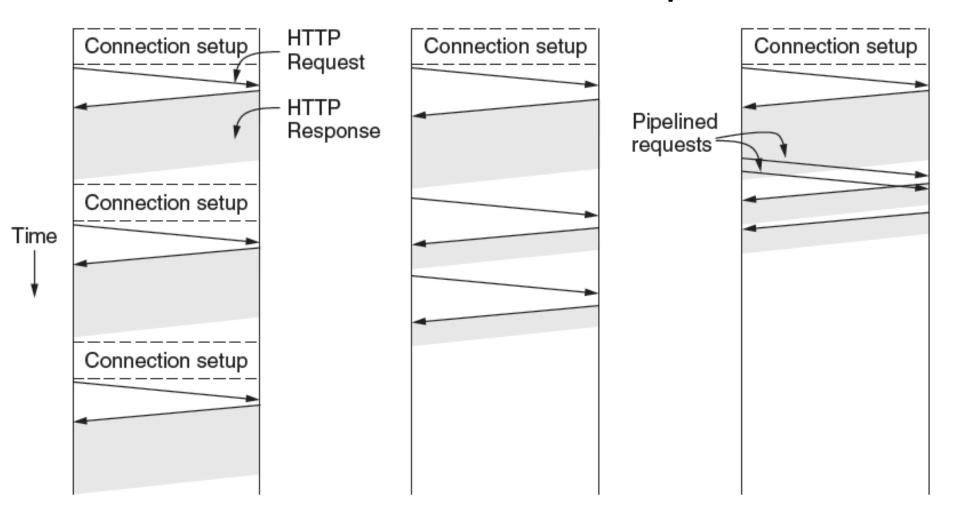
2. Type in a GET HTTP request:

GET /~ross/ HTTP/1.1 Host: cis.poly.edu by typing this in (hit carriage return twice), you send this minimal (but complete) GET request to HTTP server

3. Look at response message sent by HTTP server!

(or use Wireshark to look at captured HTTP request/response)

# HTTP communication options



Multiple connections and sequential requests.

Persistent connection and sequential requests.

Persistent connection and pipelined requests.

### Persistent HTTP

### Non-persistent HTTP

- At most one object sent over TCP connection
  - connection then closed
- Downloading multiple objects required multiple connections

#### Persistent HTTP

 Multiple objects can be sent over single TCP connection between client, server

# Non-persistent HTTP

#### **User enters URL:**

www.someSchool.edu/someDepartment/home.index

(contains text, references to 10 jpeg images)

- 1a. HTTP client initiates TCP connection to HTTP server (process) at www.someSchool.edu on port 80
- 2. HTTP client sends HTTP request message (containing URL) into TCP connection socket. Message indicates that client wants object someDepartment/home.index
- 1b. HTTP server at host
  www.someSchool.edu waiting for
  TCP connection at port 80.
  "accepts" connection, notifying
  client
- 3. HTTP server receives request message, forms response message containing requested object, and sends message into its socket

time

# Non-persistent HTTP (cont.)



4. HTTP server closes TCP connection.

- 5. HTTP client receives response message containing html file, displays html. Parsing html file, finds 10 referenced jpeg objects
- 6. Steps 1-5 repeated for each of 10 jpeg objects

time

# Non-persistent HTTP: response time

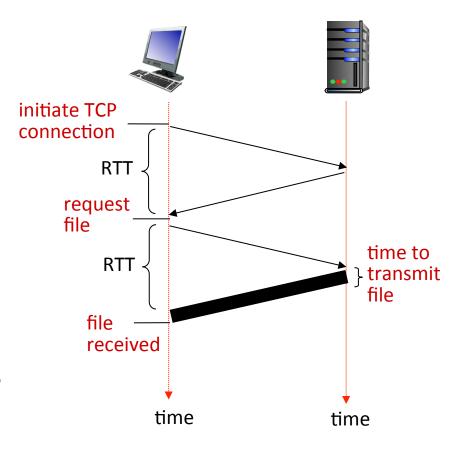
#### RTT:

Time for small packet to travel from client to server and back

#### HTTP response time:

- One RTT to initiate TCP connection
- One RTT for HTTP request and first few bytes of HTTP response to return
- File transmission time
- Non-persistent HTTP response time:

2RTT+ file transmission time



### Persistent HTTP

### Non-persistent issues:

- Requires 2 RTTs per object
- OS overhead for each TCP connection
- Browsers often open parallel TCP connections to fetch referenced objects

#### Persistent HTTP:

- Server leaves connection open after sending response
- Subsequent HTTP
   messages between same
   client/server sent over
   open connection
- Client sends requests as soon as it encounters a referenced object
- As little as one RTT for all the referenced objects

### User-server state: cookies

#### Many Web sites use cookies

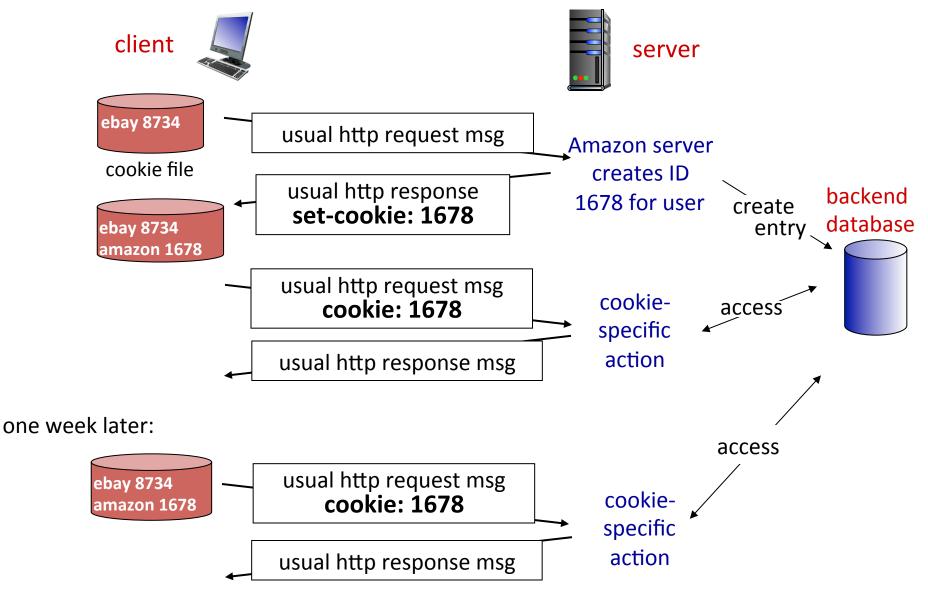
#### Four components:

- 1) Cookie header line of HTTP *response* message
- 2) Cookie header line in next HTTP *request* message
- Cookie file kept on user's host, managed by user's browser
- 4) Back-end database at Web site

#### Example:

- Susan always access Internet from PC
- Visits specific e-commerce site for first time
- When initial HTTP requests arrives at site, site creates:
  - Unique ID
  - Entry in backend database for ID

# Cookies keeping state



### Cookies

#### What cookies can be used for:

- Authorization
- Shopping carts
- Recommendations
- User session state
  - e.g. web e-mail

#### aside

#### Cookies and privacy:

- Cookies permit sites to learn a lot about you
- You may supply name and e-mail to sites

### How to keep "state":

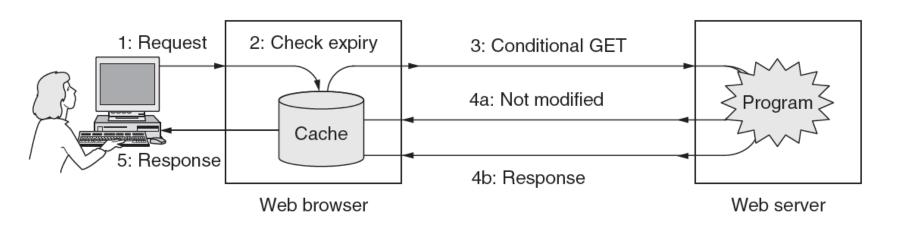
- Protocol endpoints: maintain state at sender/receiver over multiple transactions
- Cookies: HTTP messages carry state

# Improving performance

- How do we make things faster?
- Minimize traffic/latency between client/server
  - Conditional requests
  - Caching
  - Compression
- Speed up server's response
  - Multiple servers
  - Geographically distributed servers
  - Content delivery networks

# **HTTP Caching**

- Clients often cache documents
  - How and when should they check for changes?
- HTTP has cache related headers
  - HTTP/1.0: "Expires: <date>"; "Pragma: no-cache"
  - HTTP/1.1:
    - Cache-Control: No-Cache, Private, Max-age: <seconds>
    - E-tag: <some value>



### **Conditional GETs**

#### Conditional GET

- Fetch resource only if it has changed
- Server avoids wasting resources to send again
- Client sets "if-modified-since" header field
  - Server inspects "last modified" time of object
  - Returns "304 Not Modified" if unchanged, otherwise "200 OK" and new version.
- Client sets "if-no-match" using previous received
   ETag for the desired object
  - Server compares with current "hash" of object

### HTTP conditional GET

GET / HTTP/1.1

Host: katie.mtech.edu Connection: keep-alive User-Agent: Mozilla/5.0

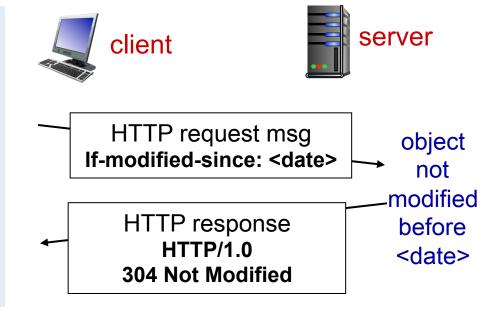
Accept: text/html, application/xhtml+xml

Accept-Encoding: gzip,deflate,sdch

If-None-Match: "c-221-4ace9c0b09cc0"

If-Modified-Since: Wed, 14 Sep 2011

17:04:27 GMT



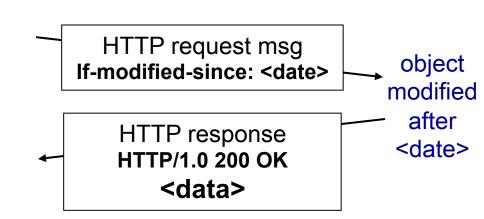
HTTP/1.1 304 Not Modified

Date: Thu, 17 Nov 2011 16:57:53 GMT

Server: Apache/2.2.16 (Debian) Keep-Alive: timeout-15, max=100

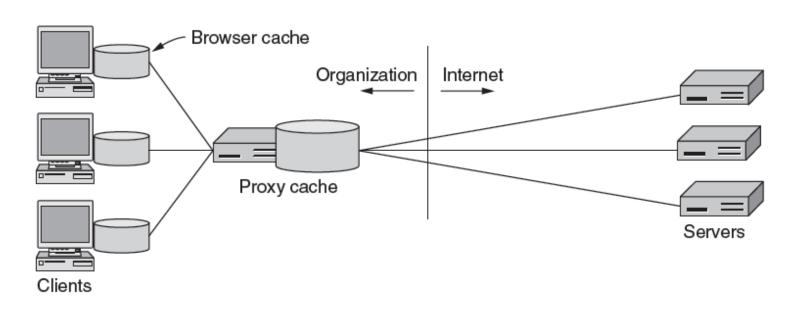
Connection: Keep-Alive

ETag: " c-221-4ace9c0b09cc0"



# Levels of caching

- Caching can occur at many levels:
  - In the client's browser
  - Client configures browser to use web proxy
  - Proxy at the ISP
  - "voluntary proxy" versus intercepting/forced/ transparent



# Summary

- HTTP protocol
  - Statelessness
    - Helps keep things simple and scale
  - Non-persistent vs. persistent connections
    - Avoid 2RTT for each object retrieved
  - Cookies
    - Adding state to web interactions
  - Caching
    - Improve performance
    - Reduce bandwidth requirements