

# More on recursion

# Overview

- Recursion

- A method calling itself

- A new way of thinking about a problem
- A powerful programming paradigm

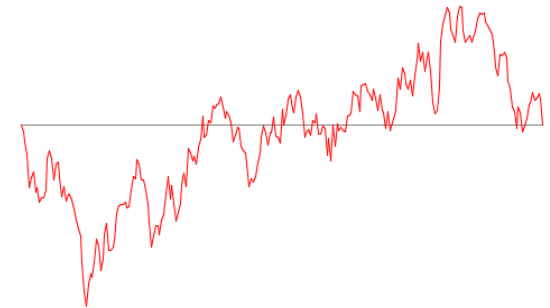
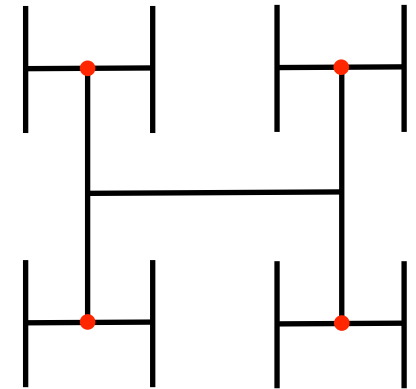
- Examples:

- Last time:

- Factorial, binary search, H-tree, Fibonacci

- Today:

- Greatest Common Divisor (GCD)
- Brownian Motion
- Sorting things



# Greatest Common Divisor

- GCD

- Find largest integer  $d$  that evenly divides  $p$  and  $q$

- e.g.  $\text{gcd}(4032, 1272) = 24$

- $4032 = 2^6 \times 3^2 \times 7^1$

- $1272 = 2^3 \times 3^1 \times 53^1$

- $\text{gcd} = 2^3 \times 3^1 = 24$

- Applications

- Simplify fractions:

$$1272/4032 = 53/168$$

- RSA cryptography

# Simple GCD algorithm

- GCD

- Find largest integer **d** that evenly divides **p** and **q**
  - Assume  $p > q$ , **p** and **q** are positive integers

- Simple algorithm:

- Set  $i = q$
- See if  $i$  evenly divides both **p** and **q**
  - If yes,  $i$  is the GCD
- Decrement  $i$
- Repeat until  $i = 1$

```
public static long gcd(long p, long q)
{
    for (long i = q; i > 1; i--)
    {
        if ((p % i == 0) && (q % i == 0))
            return i;
    }
    return 1;
}
```

# Euclid's GCD algorithm

- GCD

- Find largest integer **d** that evenly divides **p** and **q**

- Assume **p** > **q**, **p** and **q** are positive integers

- Euclid's algorithm (300 BC)



$$\text{gcd}(p, q) = \begin{cases} p & \text{if } q = 0 \quad \leftarrow \text{base case} \\ \text{gcd}(q, p \% q) & \text{otherwise} \quad \leftarrow \text{reduction step, converges to base case} \end{cases}$$

$\text{gcd}(4032, 1272) =$	
$= \text{gcd}(1272, 216)$	$4032 = 3 \times 1272 + 216$
$= \text{gcd}(216, 192)$	$1272 = 5 \times 216 + 192$
$= \text{gcd}(192, 24)$	$216 = 1 \times 192 + 24$
$= \text{gcd}(24, 0)$	$192 = 8 \times 24 + 0$
$= 24$	

# Greatest Common Divisor

- GCD

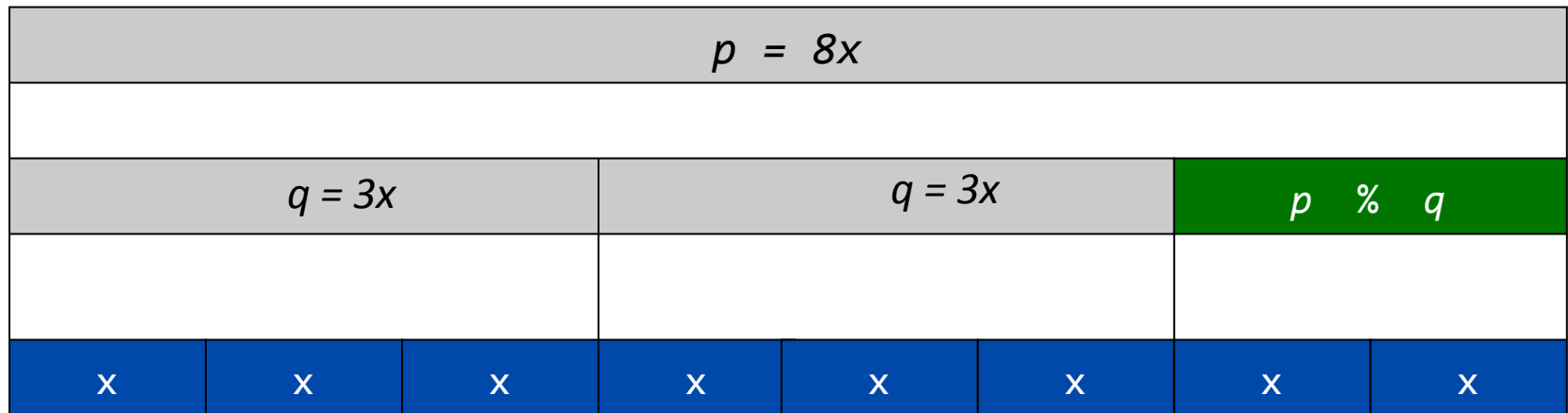
- Find largest integer **d** that evenly divides **p** and **q**

- Assume **p** > **q**, **p** and **q** are positive integers

$$\text{gcd}(p, q) = \begin{cases} p & \text{if } q = 0 \\ \text{gcd}(q, p \% q) & \text{otherwise} \end{cases}$$

← base case

← reduction step,  
converges to base case



↑  
gcd

# Greatest Common Divisor

- GCD

- Find largest integer **d** that evenly divides **p** and **q**
  - Assume **p** > **q**, **p** and **q** are positive integers

$$\text{gcd}(p, q) = \begin{cases} p & \text{if } q = 0 \\ \text{gcd}(q, p \% q) & \text{otherwise} \end{cases}$$

← base case  
← reduction step,  
converges to base case

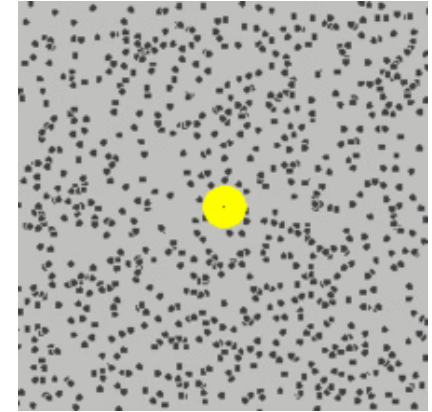
```
public static long gcd(long p, long q)
{
    if (q == 0)
        return p;
    else
        return gcd(q, p % q);
}
```

← base case  
← reduction step

# Brownian motion

- Models many natural and artificial phenomenon

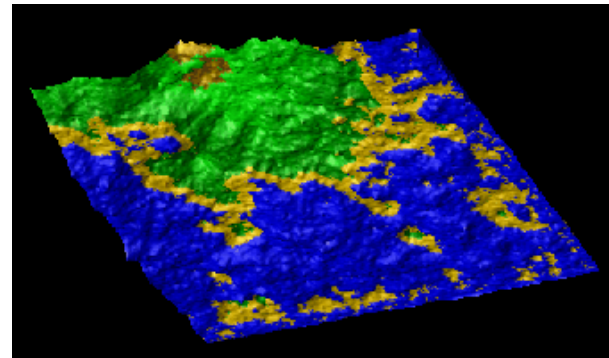
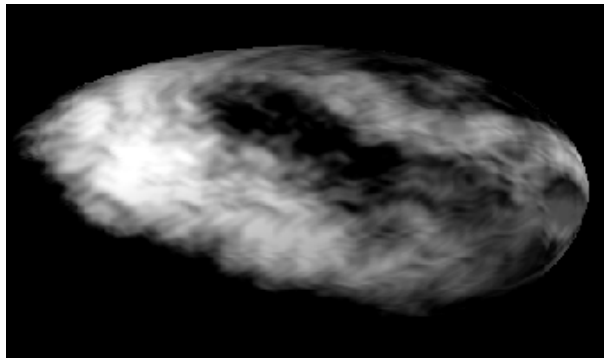
- Motion of pollen grains in water



- Price of stocks



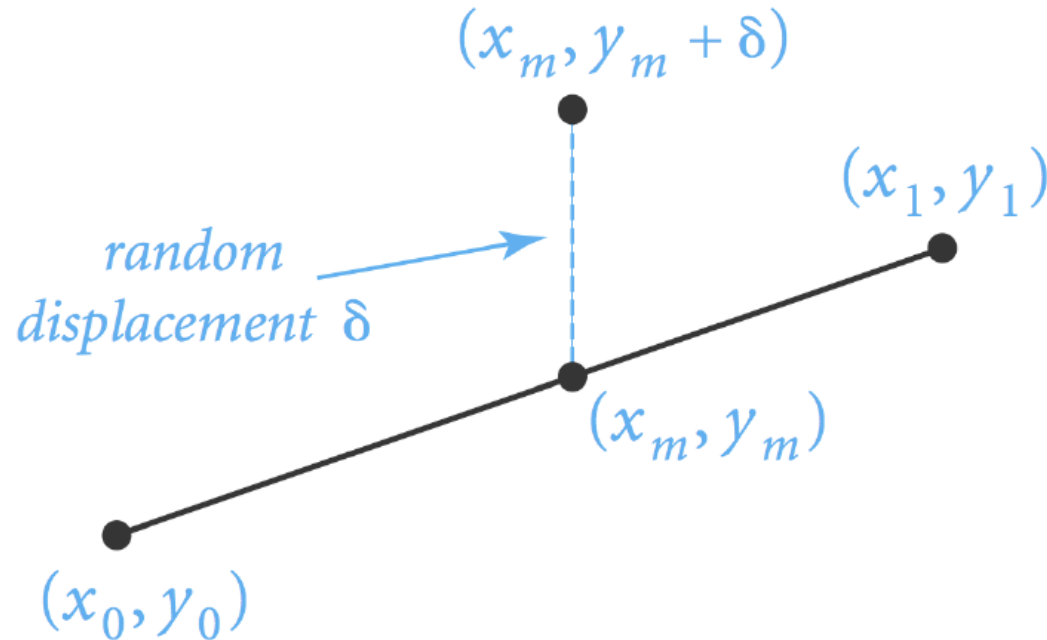
- Rugged shapes of mountains and clouds



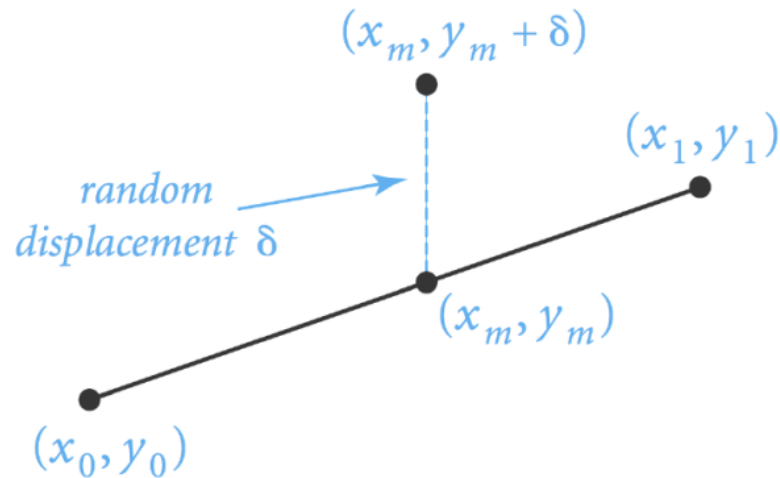


# Simulating Brownian Motion

- **Midpoint displacement method:**
  - Track interval  $(x_0, y_0)$  to  $(x_1, y_1)$
  - Choose  $\delta$  displacement randomly from Gaussian
  - Divide in half,  $x_m = (x_0 + x_1)/2$  and  $y_m = (y_0 + y_1)/2 + \delta$
  - Recur on the left and right intervals



# Recursive midpoint displacement algorithm



```
void curve(double x0, double y0, double x1, double y1, double var)
{
    if (x1 - x0 < .005)
    {
        StdDraw.Line(x0, y0, x1, y1); ← base case
        return;
    }

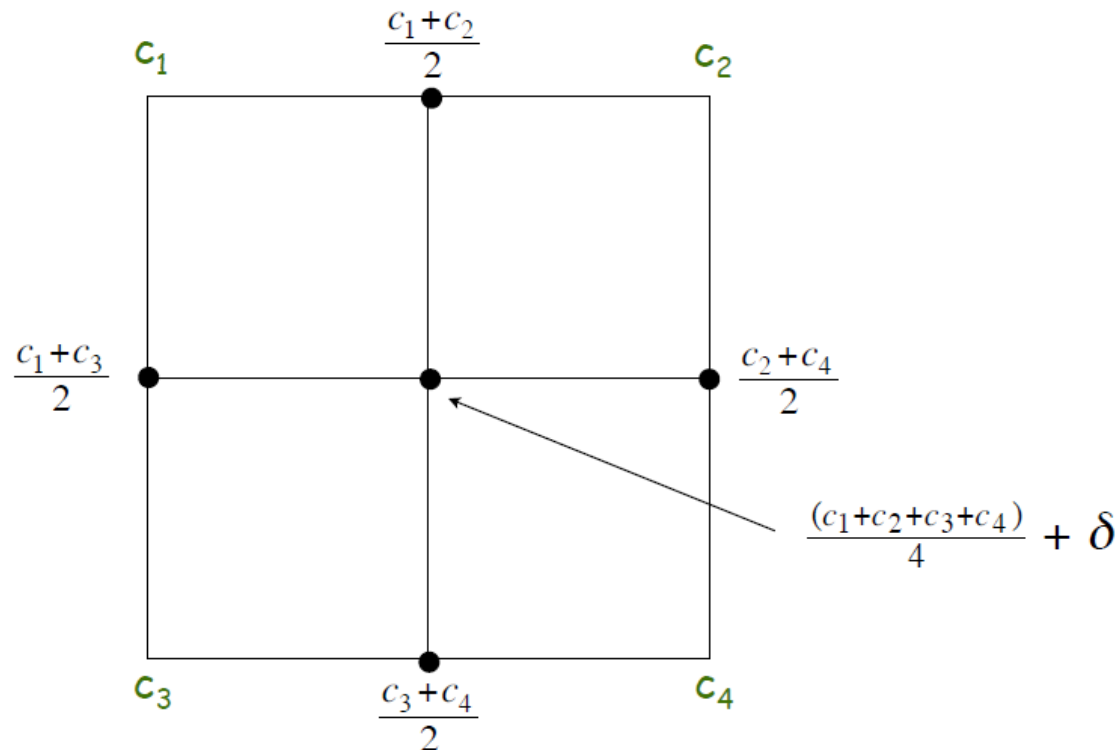
    double xm = (x0 + x1) / 2.0;
    double ym = (y0 + y1) / 2.0;

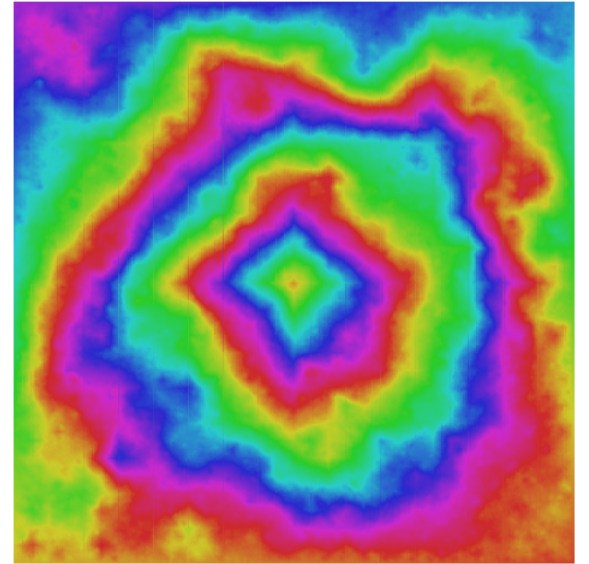
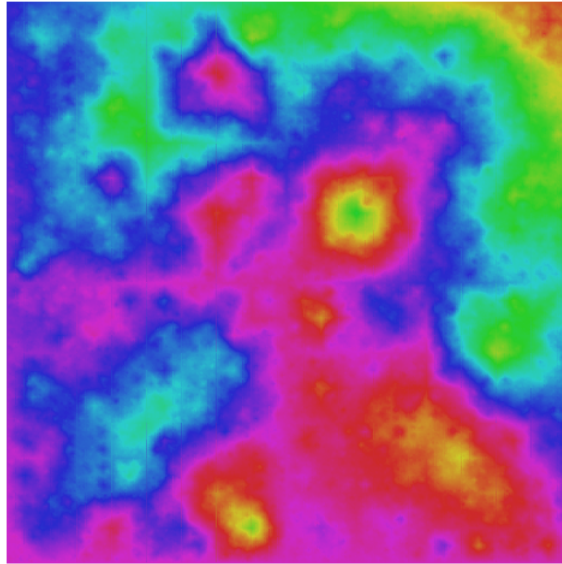
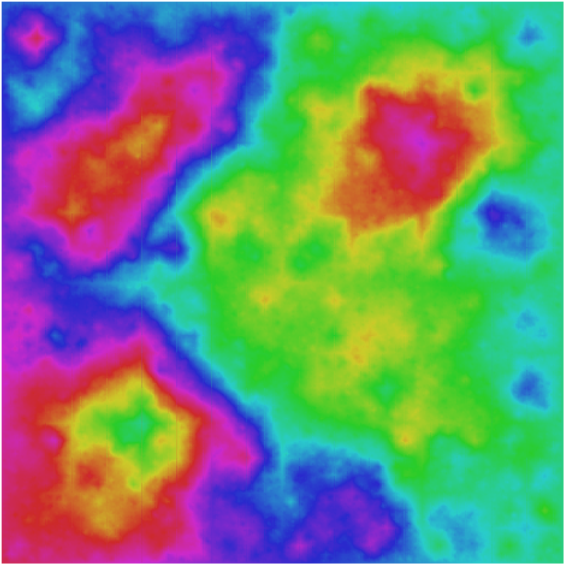
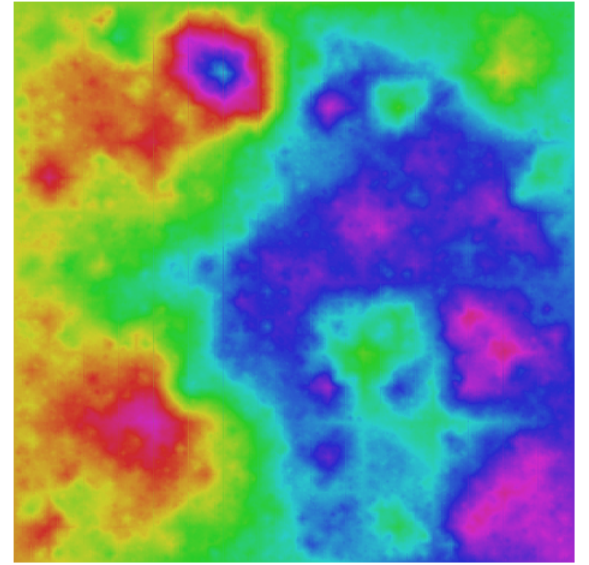
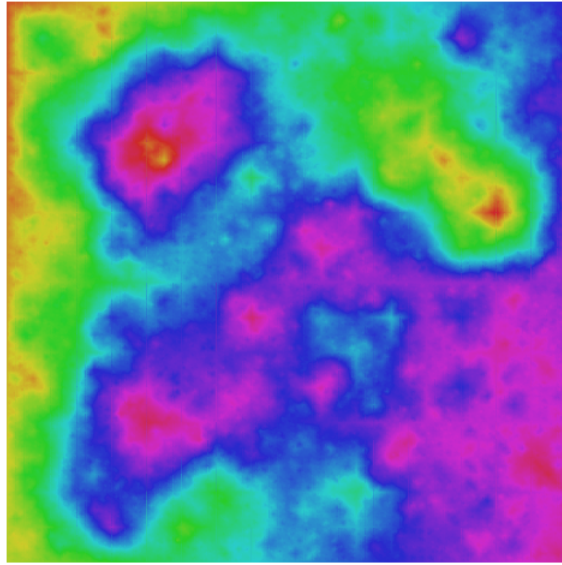
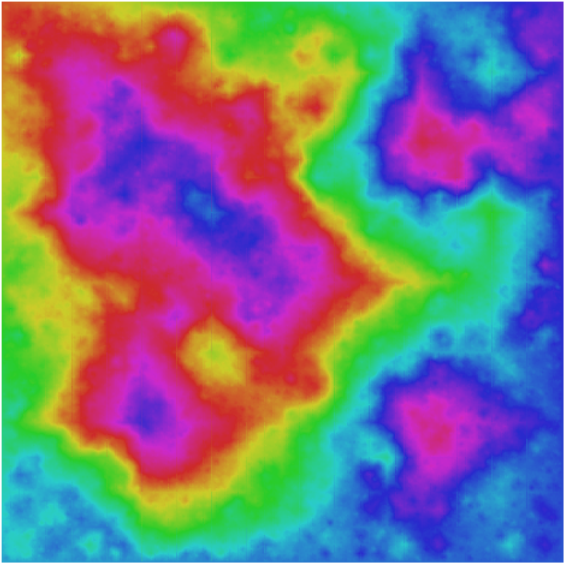
    ym = ym + StdRandom.gaussian(0, Math.sqrt(var));

    curve(x0, y0, xm, ym, var / 2.0); ← reduction step
    curve(xm, ym, x1, y1, var / 2.0);
}
```

# Plasma cloud

- Same idea, but in 2D
  - Each corner of square has some greyscale value
  - Divide into four sub-squares
  - New corners: avg of original corners, or all 4 + random
  - Recur on four sub-squares







# Brownian landscape



# Divide and conquer

- **Divide and conquer paradigm**

- Break big problem into small sub-problems
- Solve sub-problems recursively
- Combine results

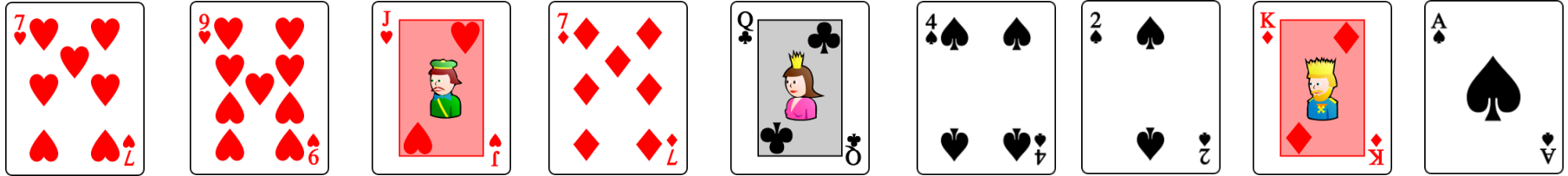
“Divide et impera. Vendi, vidi, vici.”  
-Julius Caesar

- **Used to solve many important problems**

- Sorting things, mergesort:  $O(N \log N)$
- Parsing programming languages
- Discrete FFT, signal processing
- Multiplying large numbers
- Traversing multiply linked structures (stay tuned)

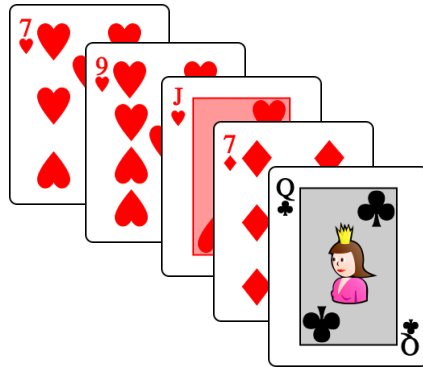
# Divide and conquer: sorting

- Goal: Sort by number, ignore suit, aces high

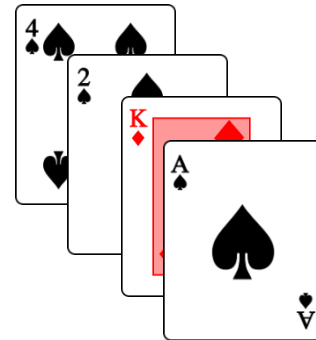


## Approach

- 1) Split in half (or as close as possible)
- 2) Give each half to somebody to sort
- 3) Take two halves and merge together



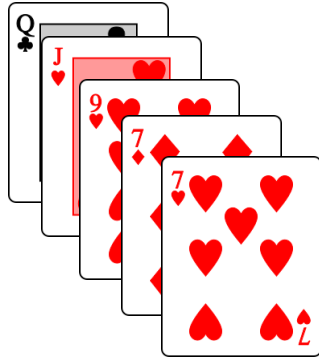
Unsorted pile #1



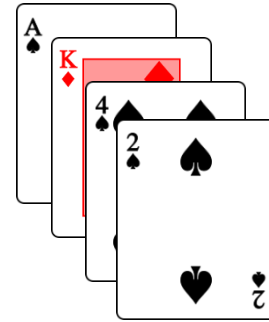
Unsorted pile #2

## Approach

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Sorted pile #1



Sorted pile #2

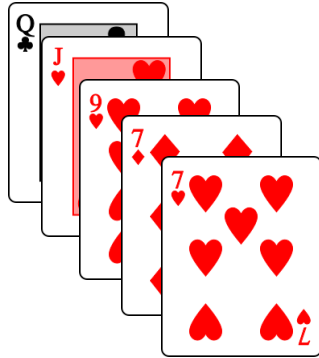
## Merging

Take card from whichever pile has lowest card

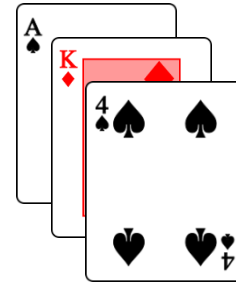


## Approach

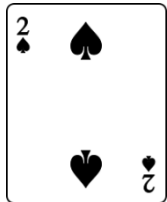
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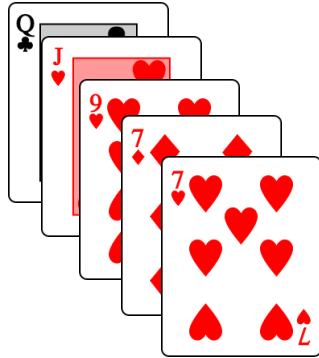


Sorted pile #2

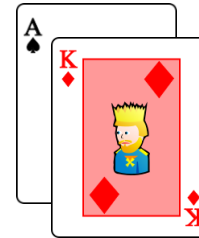


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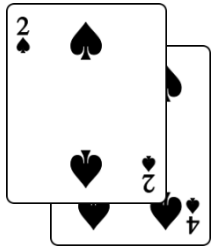
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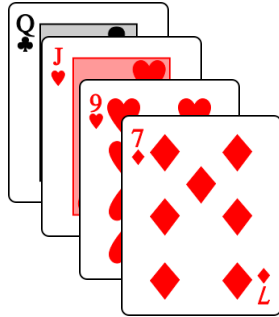


Sorted pile #2

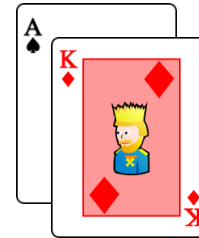


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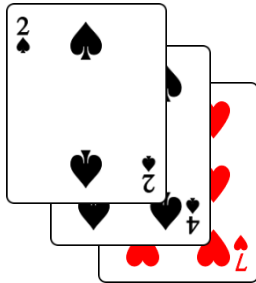
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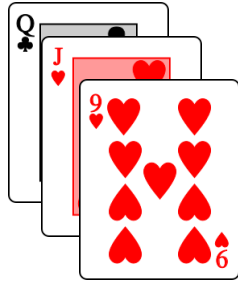


Sorted pile #2

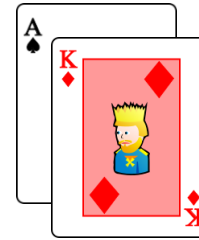


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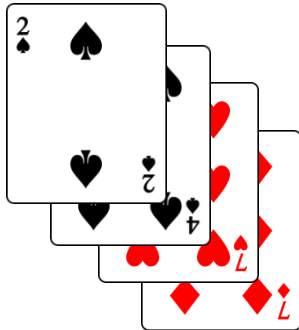
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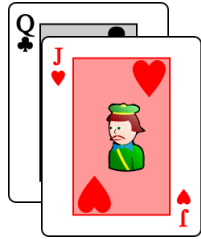


Sorted pile #2

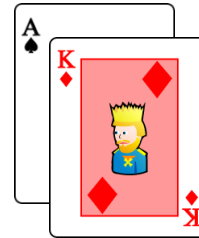


## Approach

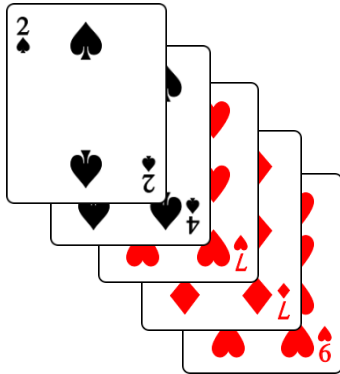
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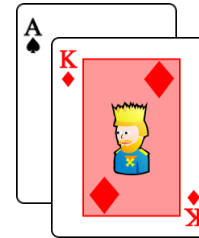
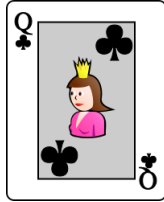


Sorted pile #2



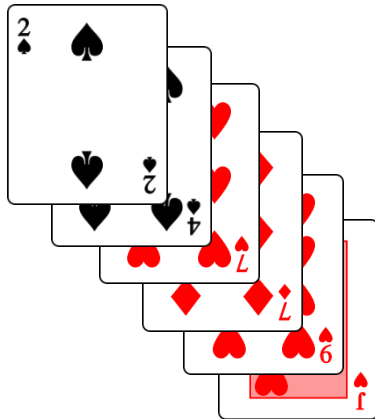
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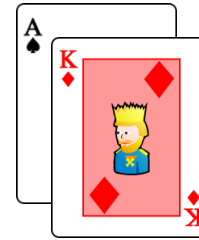
Sorted pile #1

Sorted pile #2



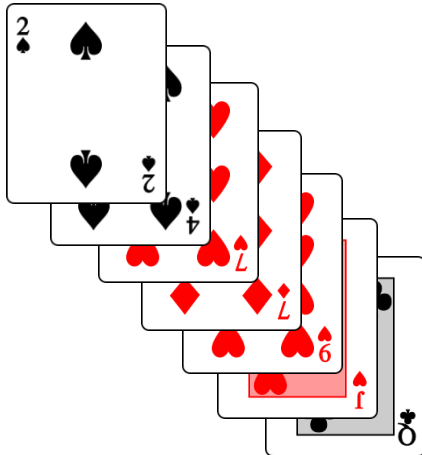
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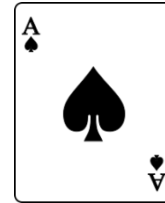
Sorted pile #1

Sorted pile #2

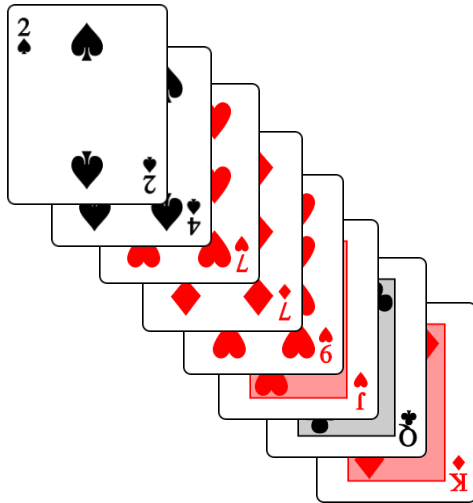


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Sorted pile #1



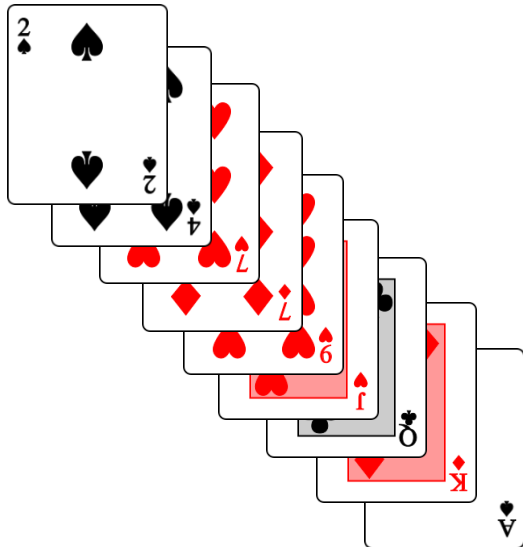
Sorted pile #2



## Approach

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Sorted pile #1



Sorted pile #2

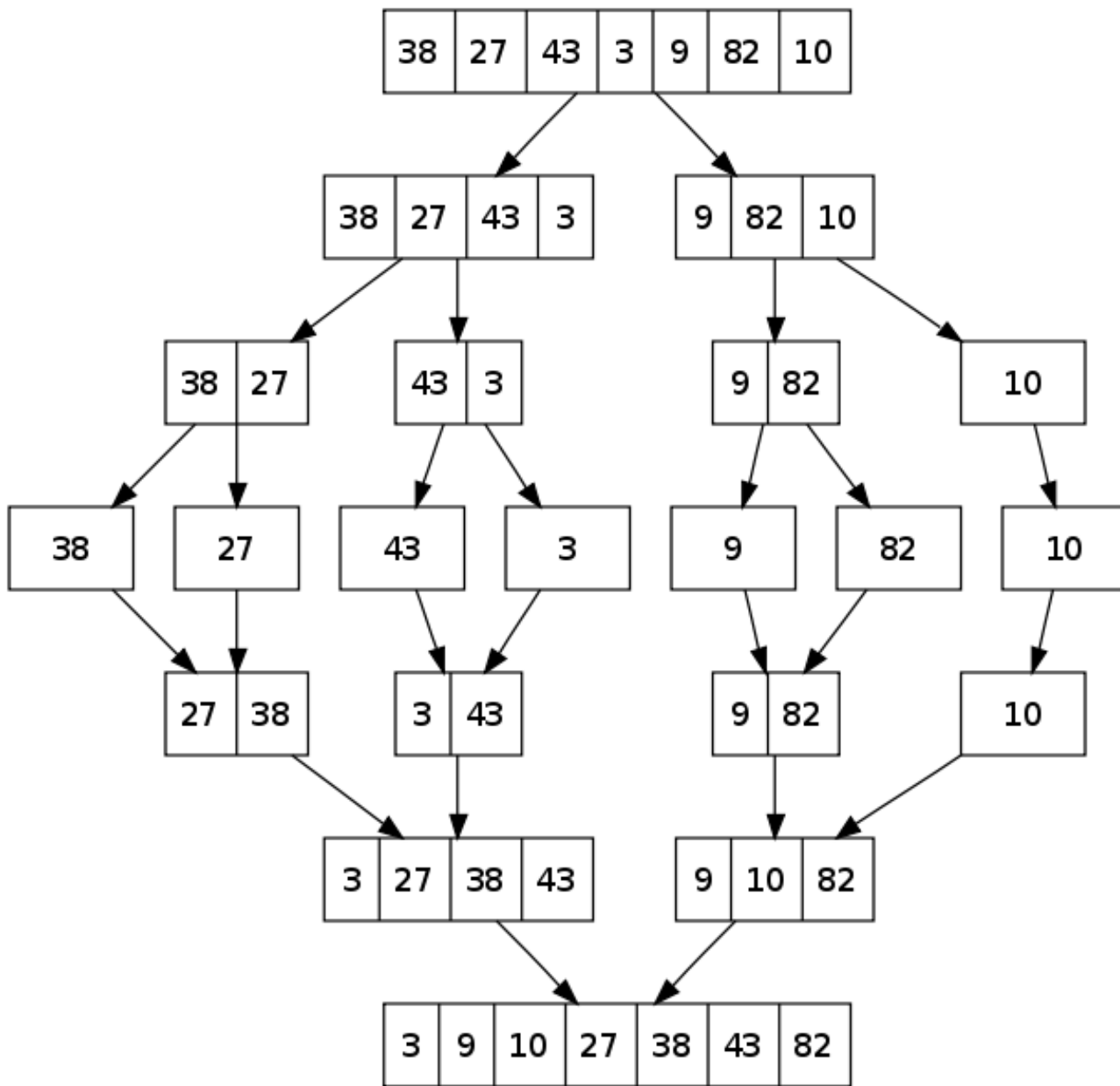
How many operations to do the merge?

Linear in the number of cards,  $O(N)$

But how did pile 1 and 2 get sorted?

**Recursively of course!**

Split each pile into two halves, give to different people to sort.



How many split levels?  
 $O(\log_2 N)$   
 How many merge levels?  
 $O(\log_2 N)$   
 Operations per level?  
 $O(N)$   
 Total operations?  
 $O(N \log_2 N)$

# Summary

- Recursion

- A method calling itself:

- Sometimes just once, e.g. binary search
- Sometimes twice, e.g. mergesort
- Sometimes multiple times, e.g. H-tree

- All good recursion must come to an end:

- Base case that does NOT call itself recursively

- A powerful tool in computer science:

- Allows elegant and easy to understand algorithms
- (Once you get your head around it)

