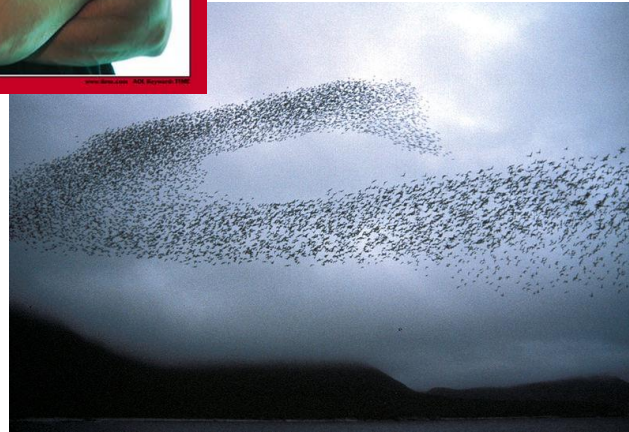
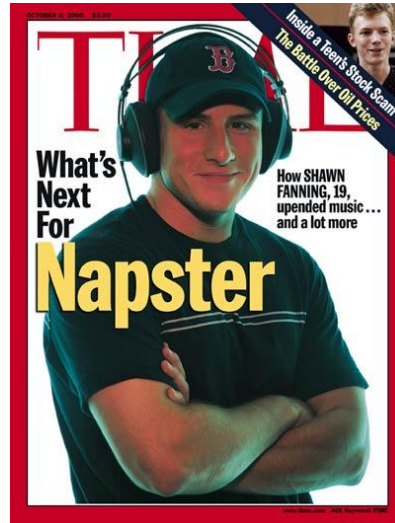
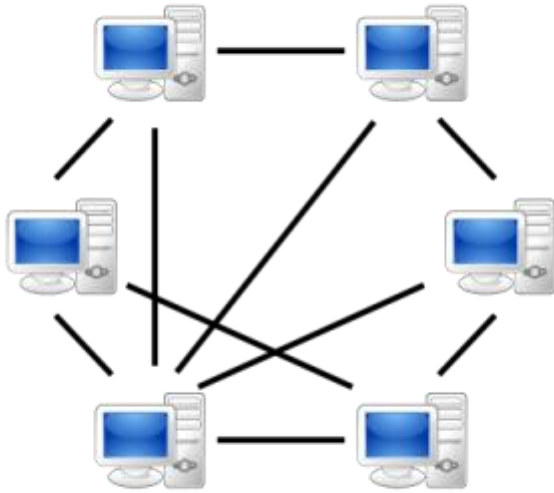


Peer-to-Peer Applications



Computer Networking: A Top Down Approach

6th edition

Jim Kurose, Keith Ross

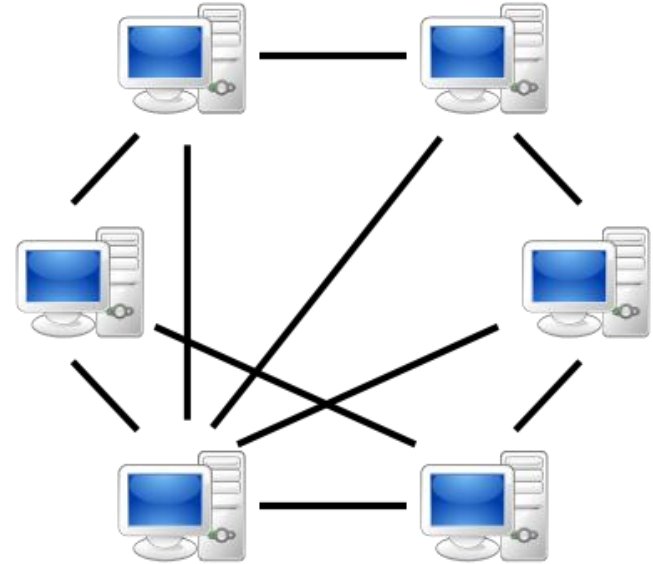
Addison-Wesley

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Overview

- Peer-to-peer applications
 - Motivation, types
 - Overlay networks
 - Napster, the rise and fall
 - Performance analysis
 - P2P vs. client-server
 - BitTorrent
 - Distributed hash tables

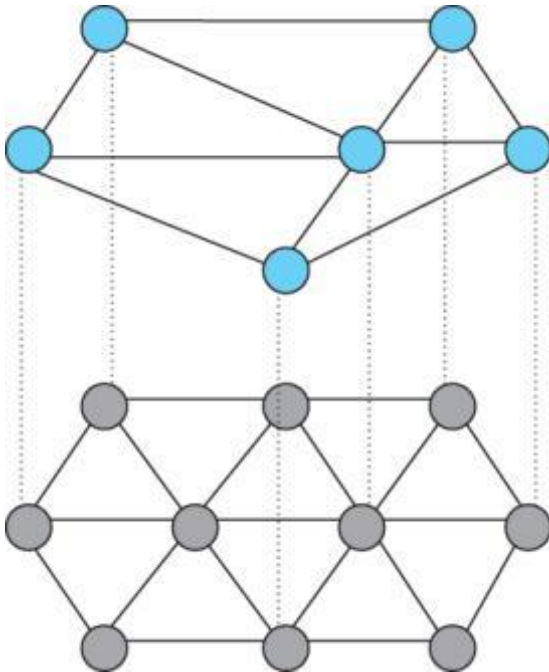


Overlay services: P2P

- Peer-to-peer (P2P) networks
 - Community of users pooling resources (storage space, bandwidth, CPU) to provide a service
 - e.g. Sharing MP3 files, Skype
 - Nodes are hosts willing to share, links are tunnels used to transport objects of interest
- Types:
 - Centralized P2P – central server for indexing
 - Pure P2P – all peers are equals
 - Hybrid P2P – some peers are supernodes

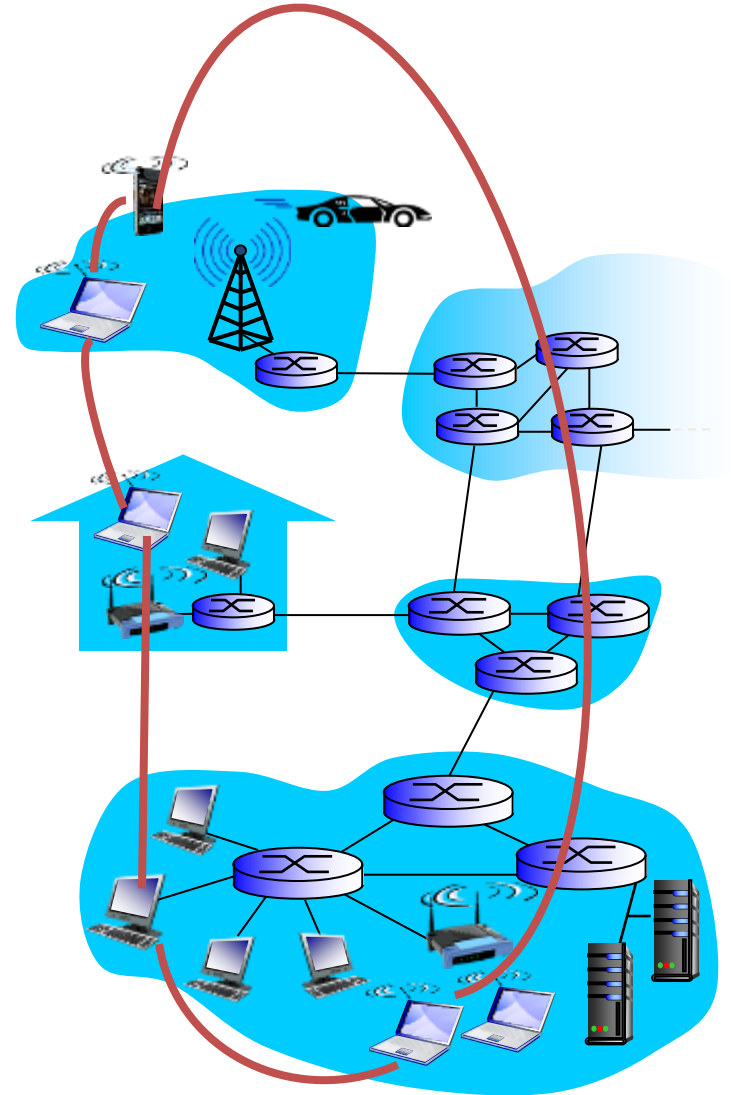
Overlay networks

- Overlay networks
 - Logical network running on top of physical network
 - Support alternate routing strategies
 - Experimental protocols



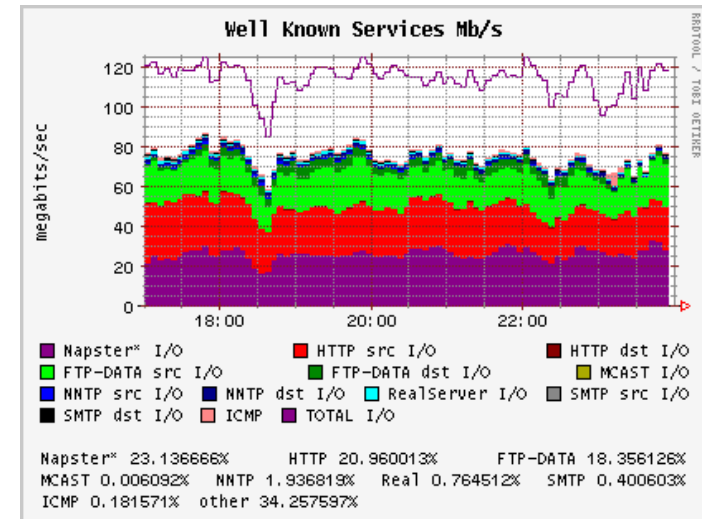
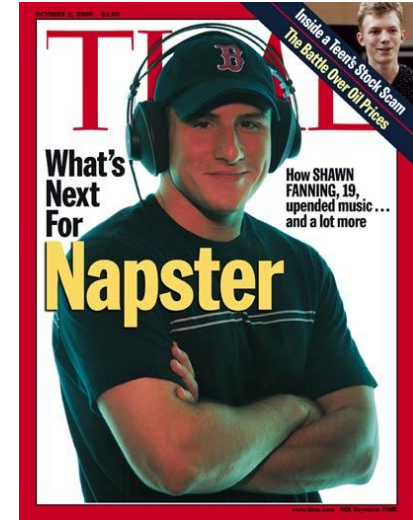
Overlay
network

Physical
network



P2P: Napster

- Napster: the rise
 - Created by Shawn Fanning
 - Christmas break, freshmen year at college
 - Allows search and sharing of MP3s
 - January 1999, Napster version 1.0
 - May 1999
 - Company founded
 - Shawn drops out of school
 - September 1999, 1st lawsuits
 - No such thing as bad publicity
 - By 2000, 80 million users

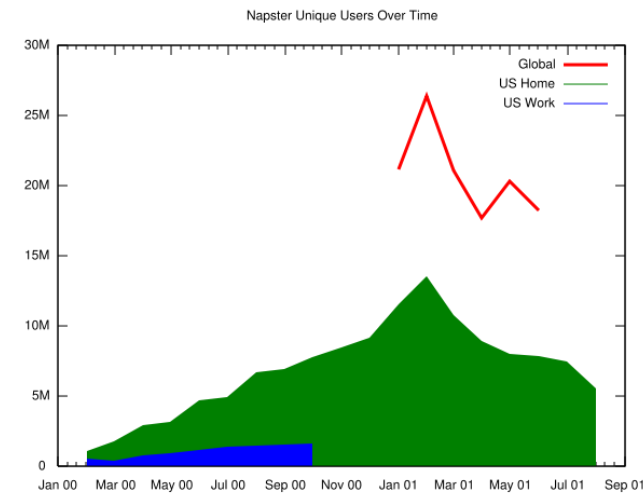


UW-Madison, March 9th, 2000

P2P: Napster

- Napster: the fall

- December 1999, RIAA lawsuit
- Metallica's "I Disappear" circulates
 - Before official release, starts getting radio play
 - 2000 band files a lawsuit
- July 2001, shutdown by lawsuits
- 2002, relaunched as paid service
 - Record labels not keen to license
 - Files bankruptcy
- Gave rise to many P2P alternatives
- Forced industry out of stone age
 - iTunes



Napster users peak, Feb 2001.

Napster technology

- User installs software
 - Registers name, password, local dir with music
- Client contacts central Napster server
 - Connects via TCP
 - Provides list of music in user's directory
 - Napster updates its database
- Client searches for music
 - Napster identifies currently online client with file
 - Provides IP addresses so client can download directly



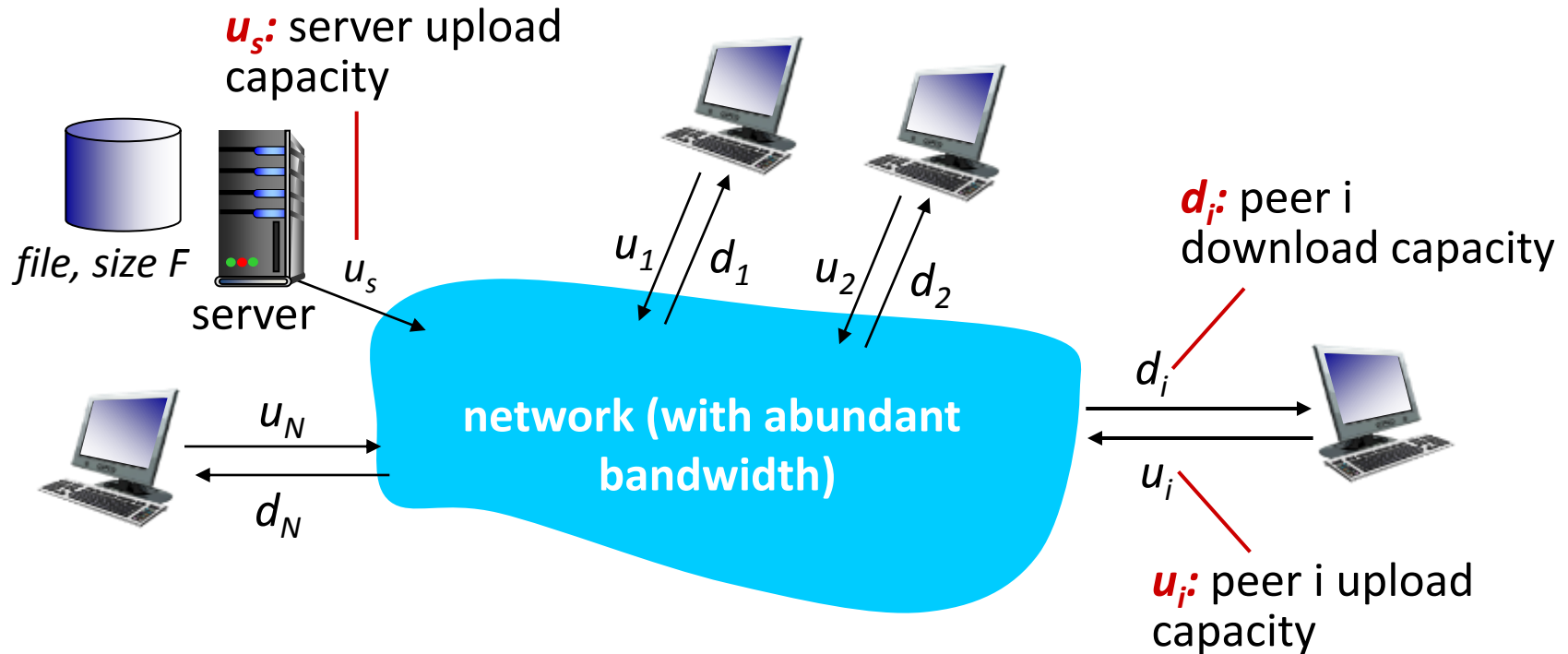
Napster technology

- Central server continually updated
 - Easy to track music currently available and from what peer
 - Good source to prove copyright infringement
 - Single point of failure, performance bottleneck
- Peer-to-peer transfer
 - Key idea of P2P: heavy lifting done between peers
 - No need for Napster to provision lots of capacity
 - Just enough to support indexing/search needs of clients
- Proprietary protocol

File distribution: client-server vs. P2P

Question: Time to distribute file (size F) from one server to N peers?

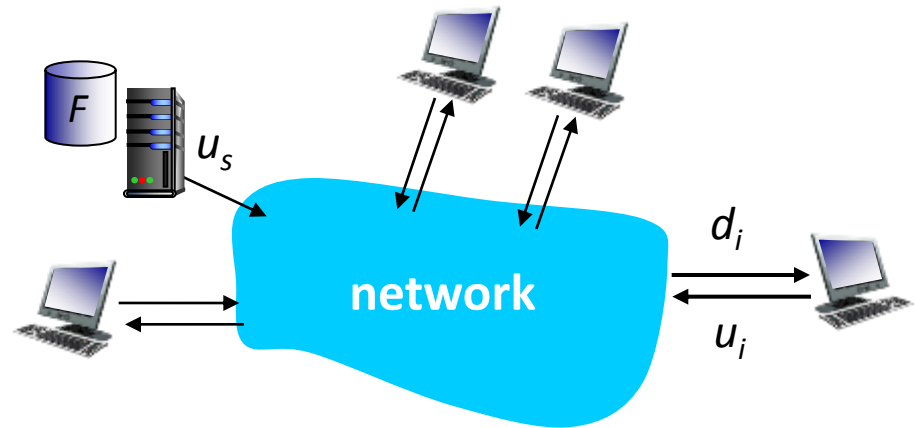
- Peer upload/download capacity is limited resource



File distribution time: client-server

- *Server transmission*: must sequentially send (upload) N file copies:
 - Time to send one copy: F/u_s
 - Time to send N copies: NF/u_s

- ❖ *Client*: each client must download file copy
 - d_{\min} = min client download rate
 - Min client download time: F/d_{\min}



*Time to distribute F
to N clients using
client-server approach*

$$D_{c-s} \geq \max\{NF/u_s, F/d_{\min}\}$$

increases linearly in N

File distribution time: P2P

- **Server transmission:**

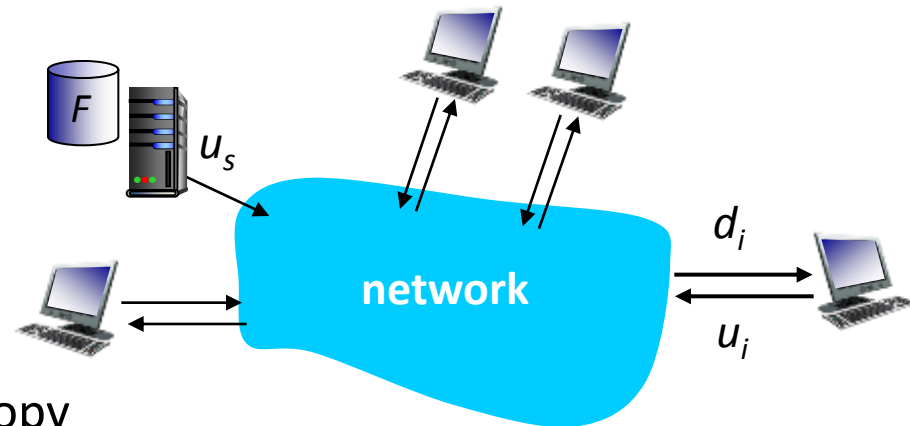
- Must upload at least one copy
- Time to send one copy: F/u_s

- **Client:**

- Each client must download file copy
- Min client download time: F/d_{min}

- **Clients:**

- Aggregate download of NF bits
- Max upload rate (limiting max download rate) is $u_s + \sum u_i$



*Time to distribute F
to N clients using
P2P approach*

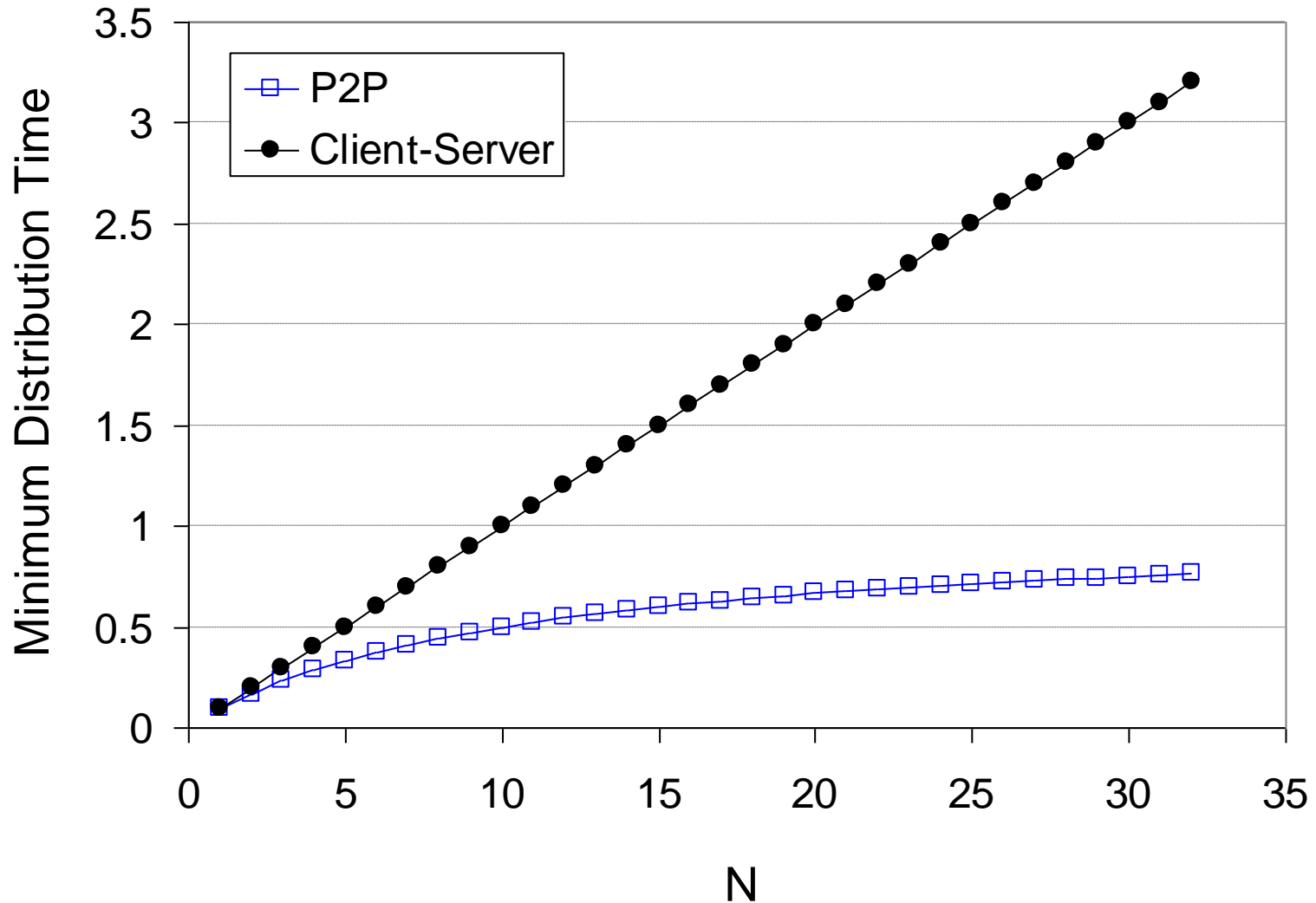
$$D_{P2P} > \underline{\max} \{ F/u_s, F/d_{min}, NF / (u_s + \sum u_i) \}$$

increases linearly in N ...

... but so does this, as each peer brings service capacity

Client-server vs. P2P example

Client upload rate = u , $F/u = 1$ hour, $u_s = 10u$, $d_{min} \geq u_s$



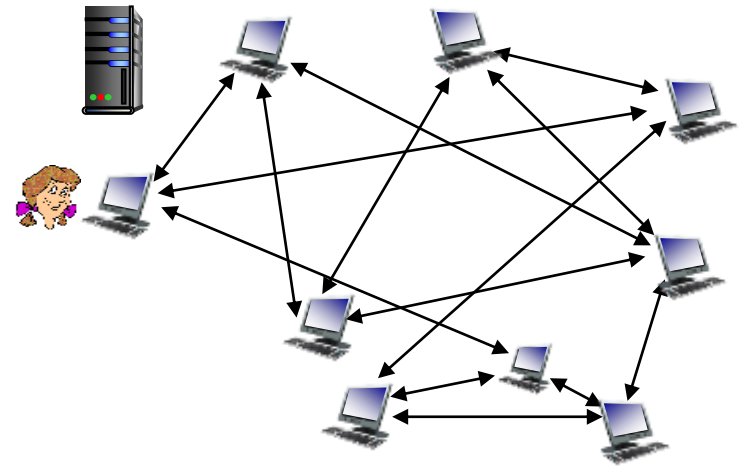
P2P: BitTorrent

- BitTorrent protocol
 - 2001, Bram Cohen releases first implementation
 - Now supported by many different clients
 - 2011, ~100 million users
- Motivations:
 - Serve up popular content fast
 - Popularity exhibits temporal locality
 - Efficient fetching, not searching
 - Distribute same file to many peers
 - Single publisher, many downloaders
 - Measures to prevent free-loading



BitTorrent process

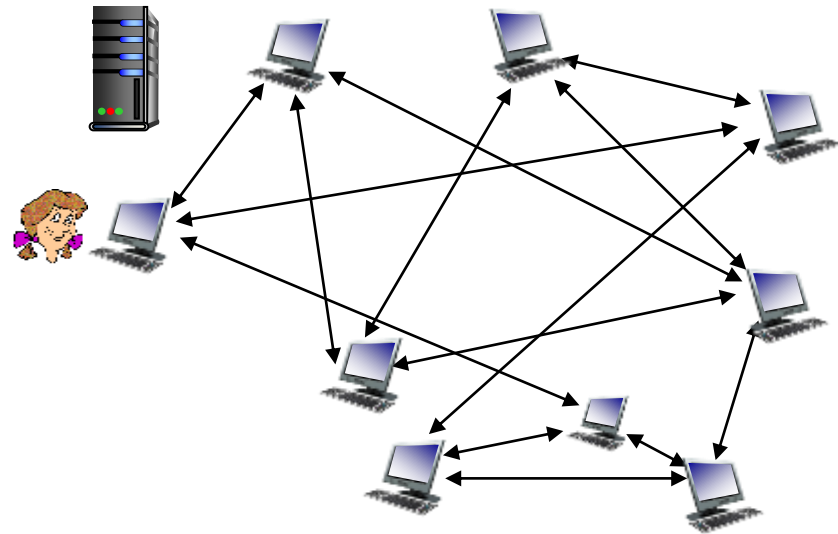
- File divided into many 256KB chunks
 - Peers exchange the pieces by uploading and downloading to each other
 - Seed: peer with entire file
- Process:
 - Users find torrent of interest, open in client
 - Client contacts the *tracker* listed in torrent file
 - Gets list of peers currently transferring the file
 - Joins the *swarm*
 - Peers currently with some/all of the file



BitTorrent process

- Peer joining torrent:

- Has no chunks, but will accumulate them over time from other peers
- Registers with tracker to get list of peers, connects to subset of peers, "neighbors"

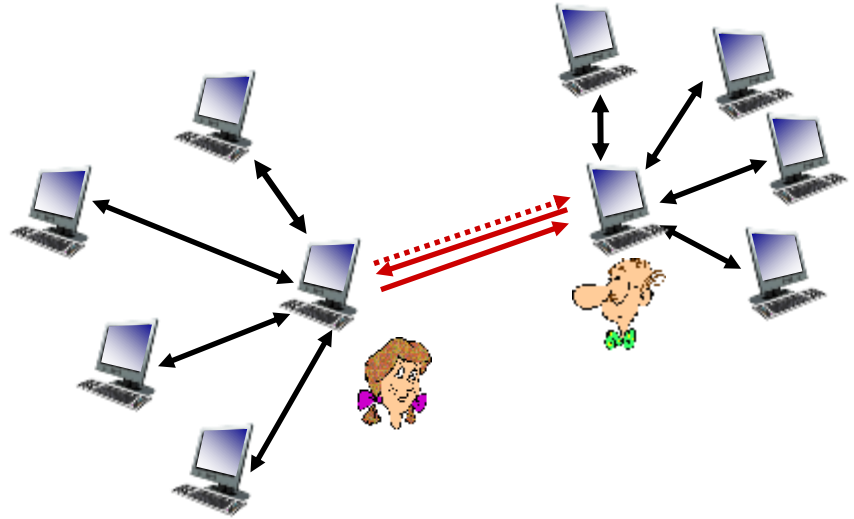


- ❖ While downloading, peer uploads chunks to other peers
- ❖ Peer may change peers with whom it exchanges chunks
- ❖ *Churn*: peers may come and go
- ❖ Once peer has entire file it may (selfishly) leave or (altruistically) remain in torrent

BitTorrent: requesting, sending file chunks

Requesting chunks:

- At any given time, different peers have different subsets of file chunks
- Periodically, Alice asks each peer for list of chunks that they have
- Alice requests missing chunks from peers, rarest first

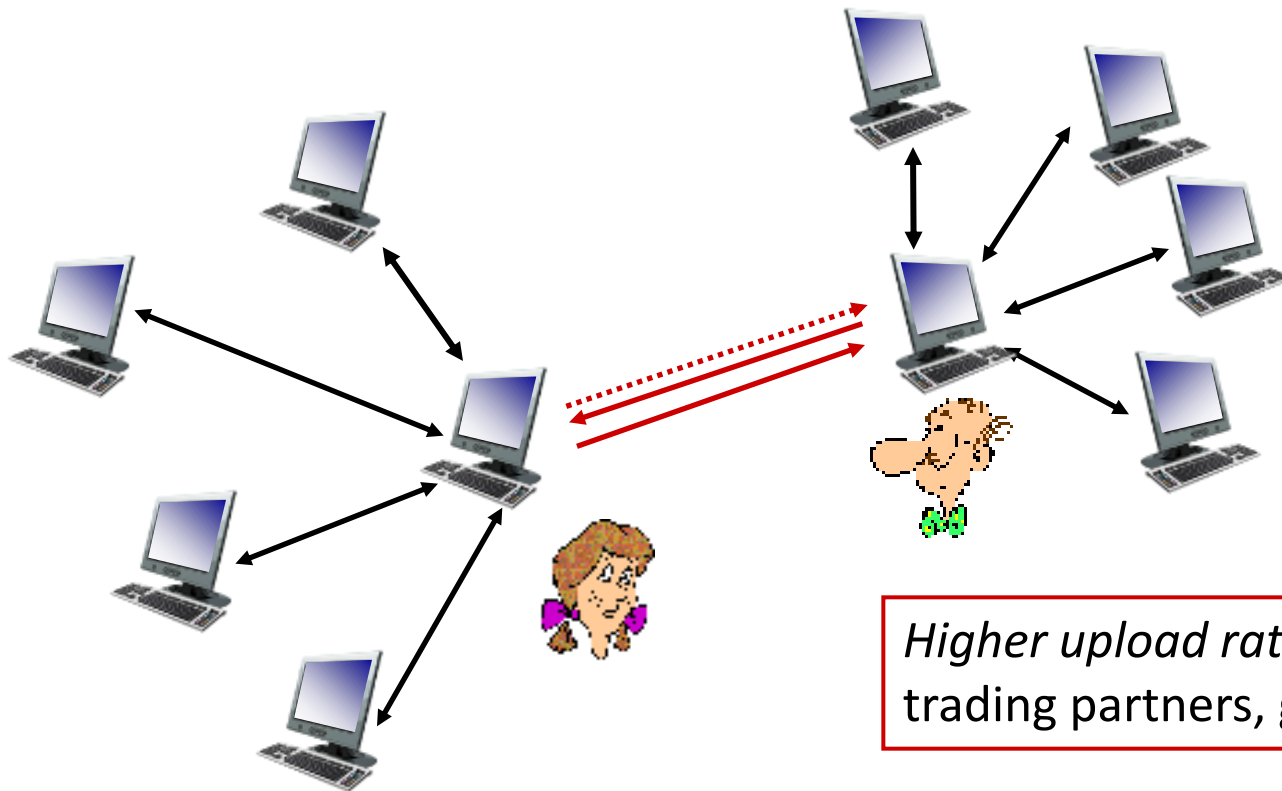


Sending chunks: tit-for-tat

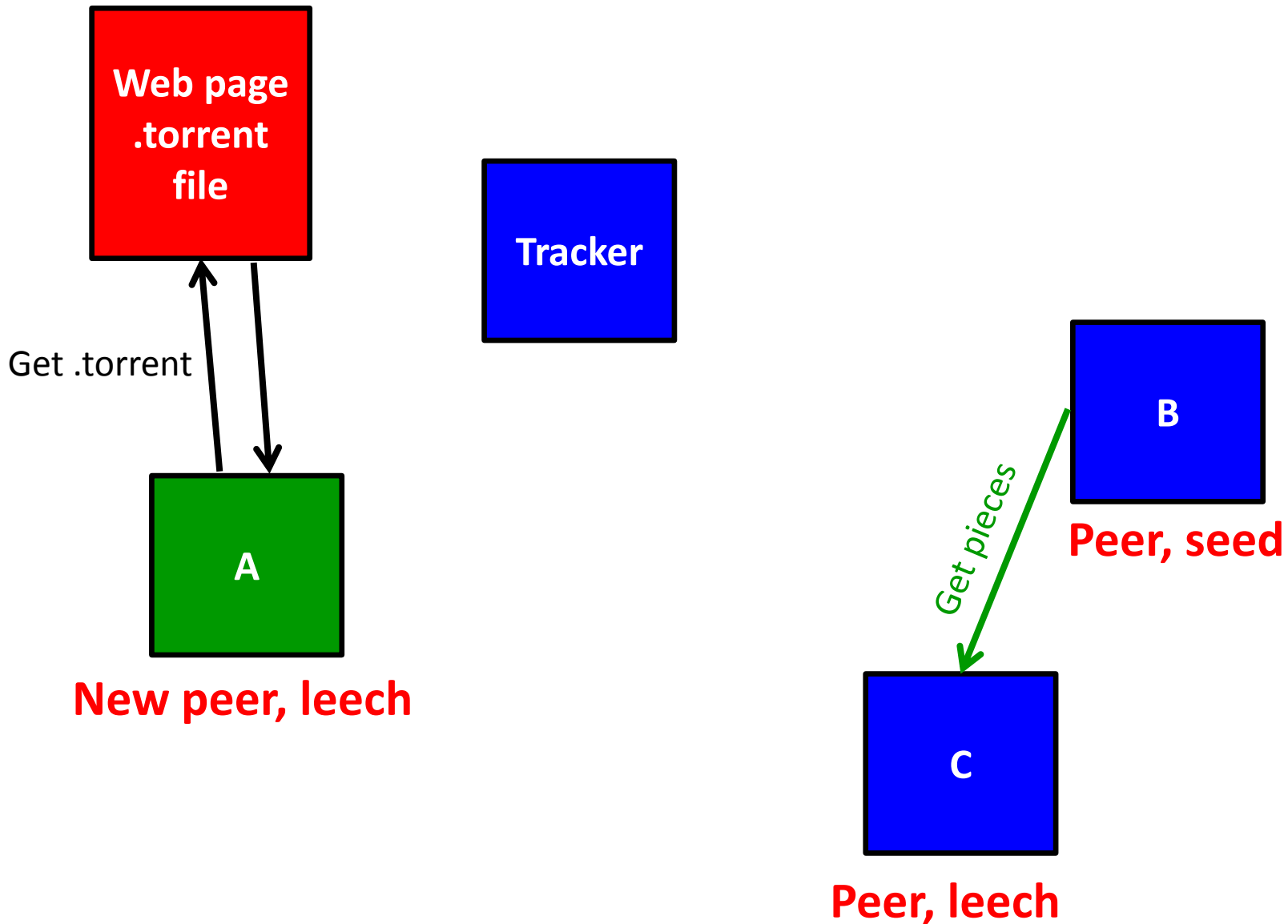
- ❖ Alice sends chunks to 4 peers currently sending her chunks *at highest rate*
 - Other peers are choked by Alice
 - Re-evaluate top 4 every 10 seconds
- ❖ Every 30 secs: randomly select another peer, starts sending chunks
 - "Optimistically unchoke" this peer
 - Newly chosen peer may join top 4

BitTorrent: tit-for-tat

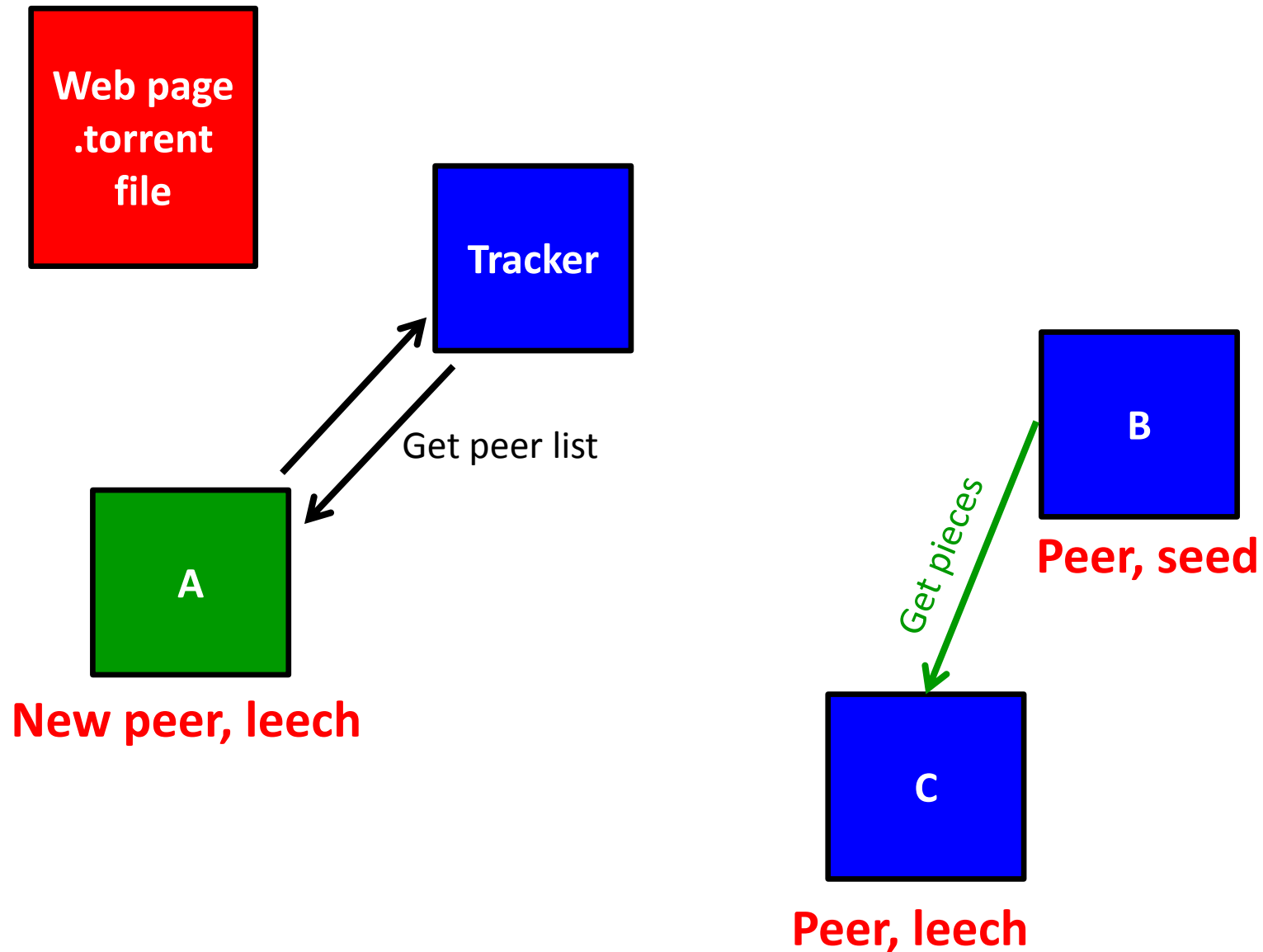
- (1) Alice "optimistically unchokes" Bob
- (2) Alice becomes one of Bob's top-four providers; Bob reciprocates
- (3) Bob becomes one of Alice's top-four providers



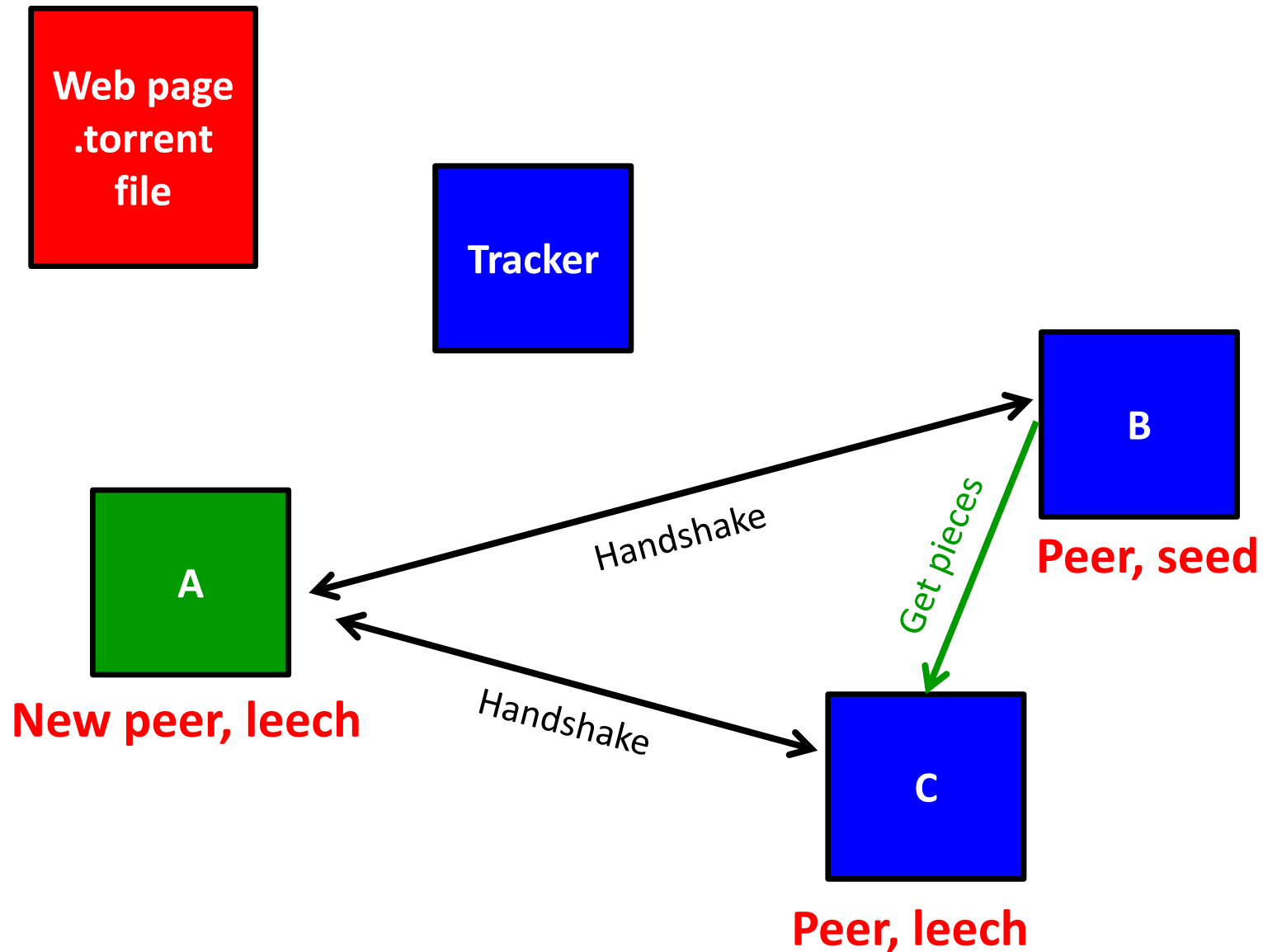
BitTorrent process



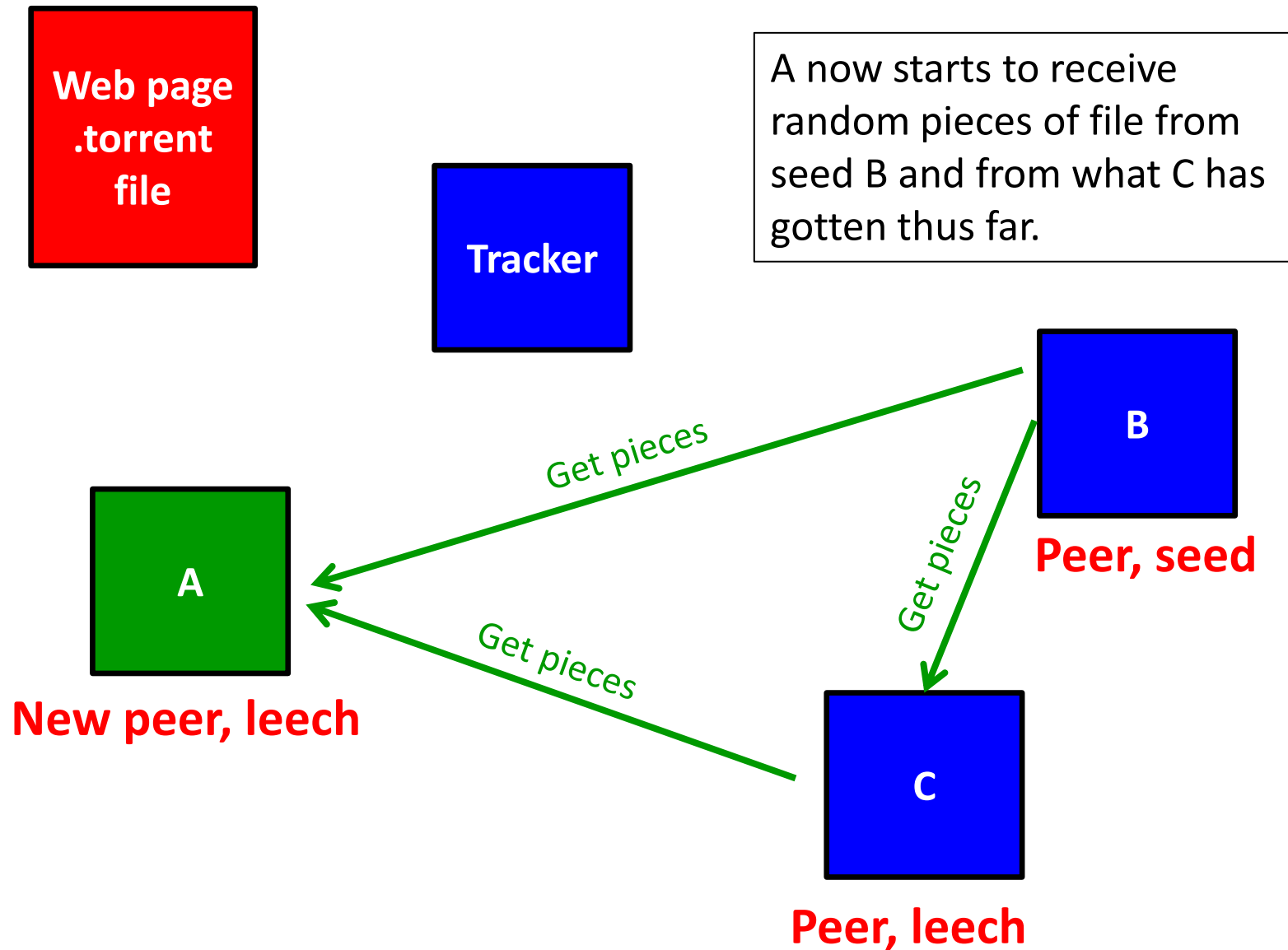
BitTorrent process



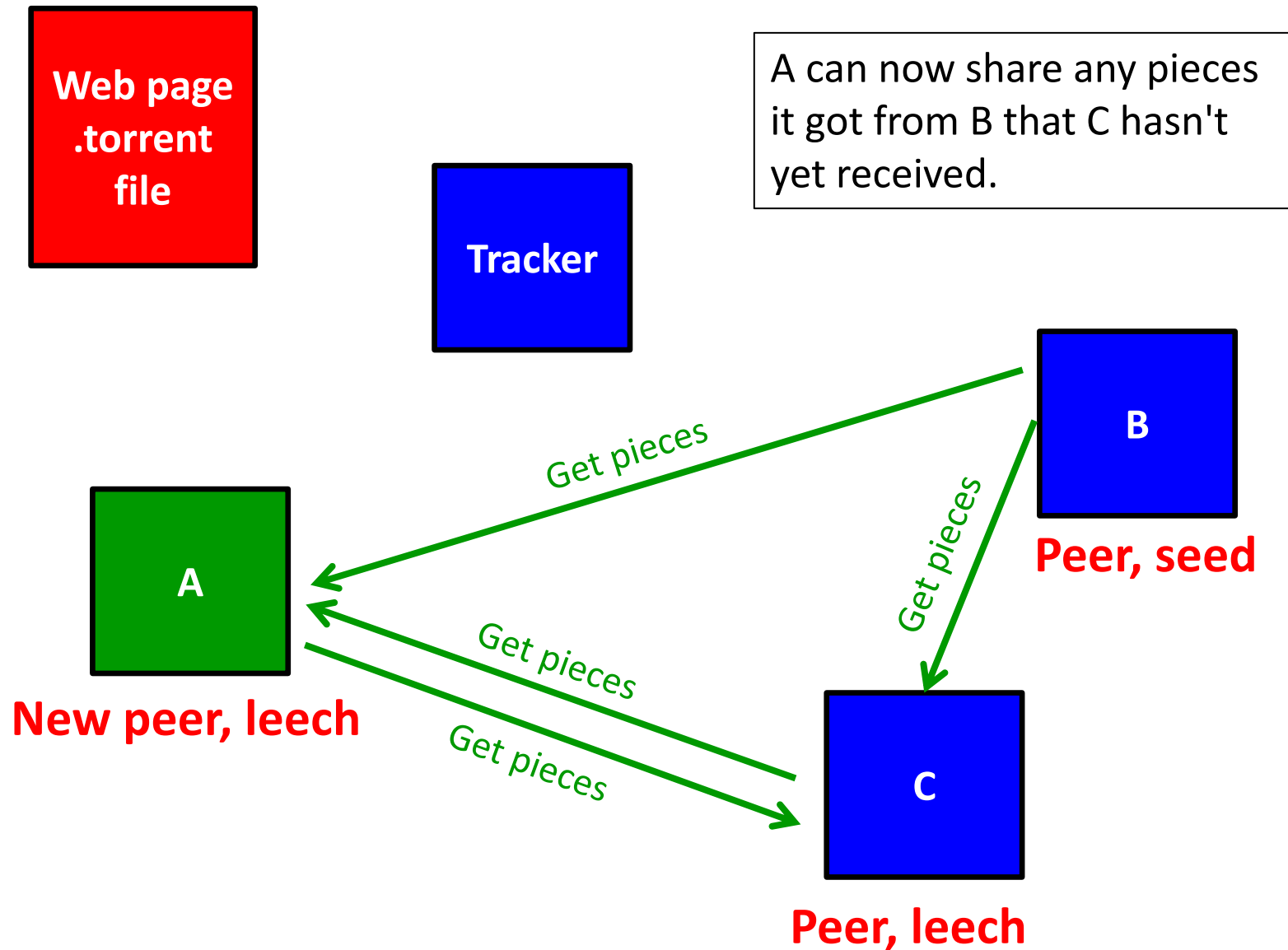
BitTorrent process



BitTorrent process

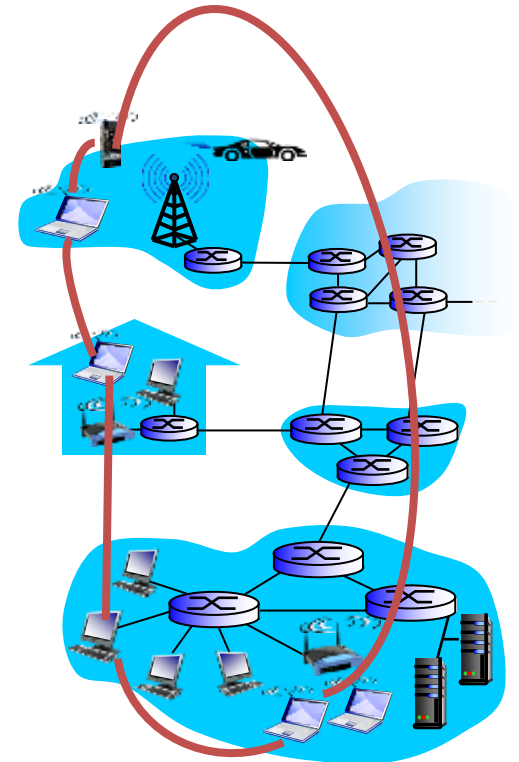
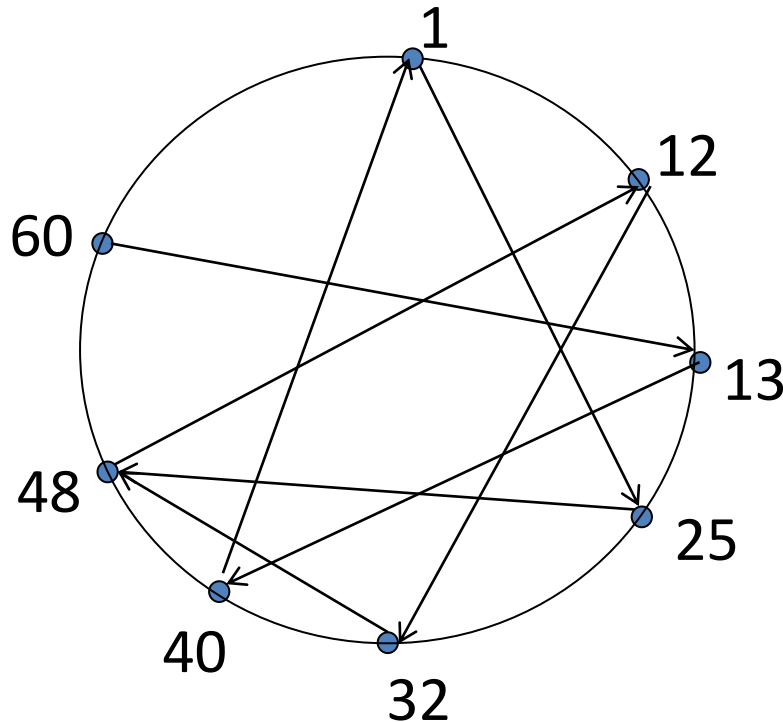


BitTorrent process



Distributed Hash Table (DHT)

- Hash table
- DHT paradigm
- Circular DHT and overlay networks
- Peer churn



Simple database

- Simple database with (key, value) pairs:
 - **Key:** human name
 - **Value:** social security #

Key	Value
John Washington	132-54-3570
Diana Louise Jones	761-55-3791
Xiaoming Liu	385-41-0902
Rakesh Gopal	441-89-1956
Linda Cohen	217-66-5609
.....
Lisa Kobayashi	177-23-0199

- **Key:** movie title
- **Value:** IP address of system storing movie

Hash Table

- More convenient:
 - Store/search on numerical representation of key
 - $\text{Key} = \text{hash}(\text{original key})$

Original Key	Key	Value
John Washington	8962458	132-54-3570
Diana Louise Jones	7800356	761-55-3791
Xiaoming Liu	1567109	385-41-0902
Rakesh Gopal	2360012	441-89-1956
Linda Cohen	5430938	217-66-5609
.....	
Lisa Kobayashi	9290124	177-23-0199

Distributed Hash Table (DHT)

- Distribute (key, value) pairs over millions of peers
 - Pairs are evenly distributed over peers
- Any peer can query database with a key
 - Database returns value for the key
 - To resolve query, small number of messages exchanged among peers
- Peer only knows a small number of other peers
- Robust to peers coming and going, *churn*

Assign key-value pairs to peers

- Rules:

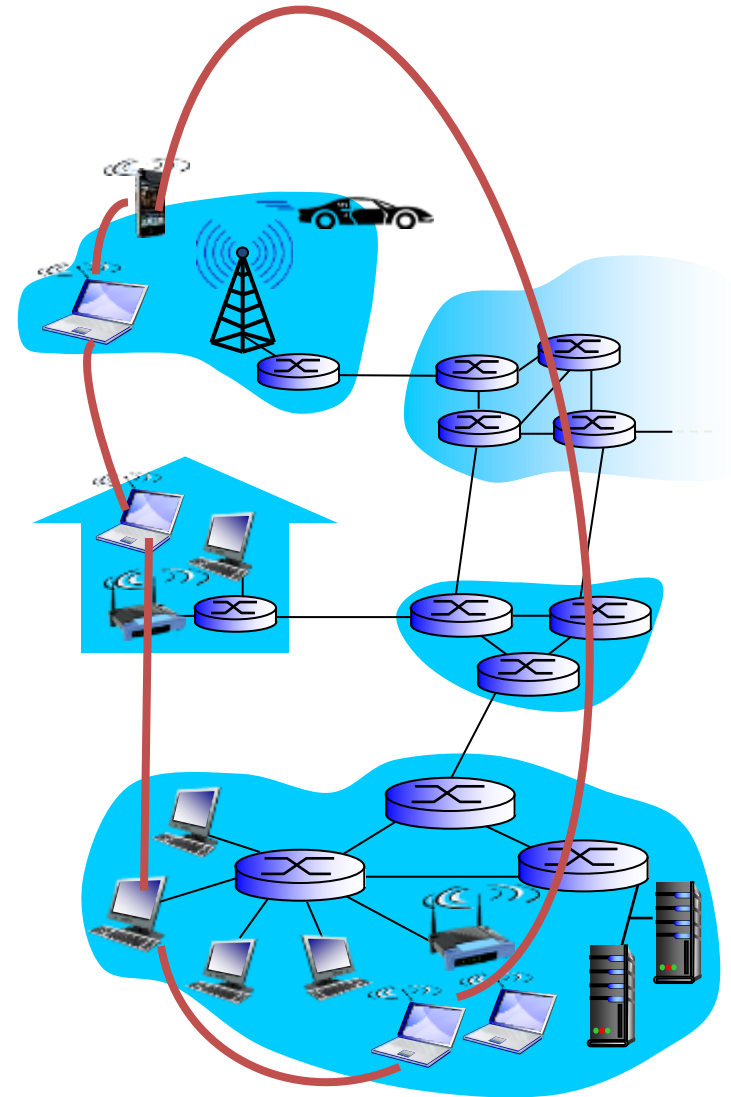
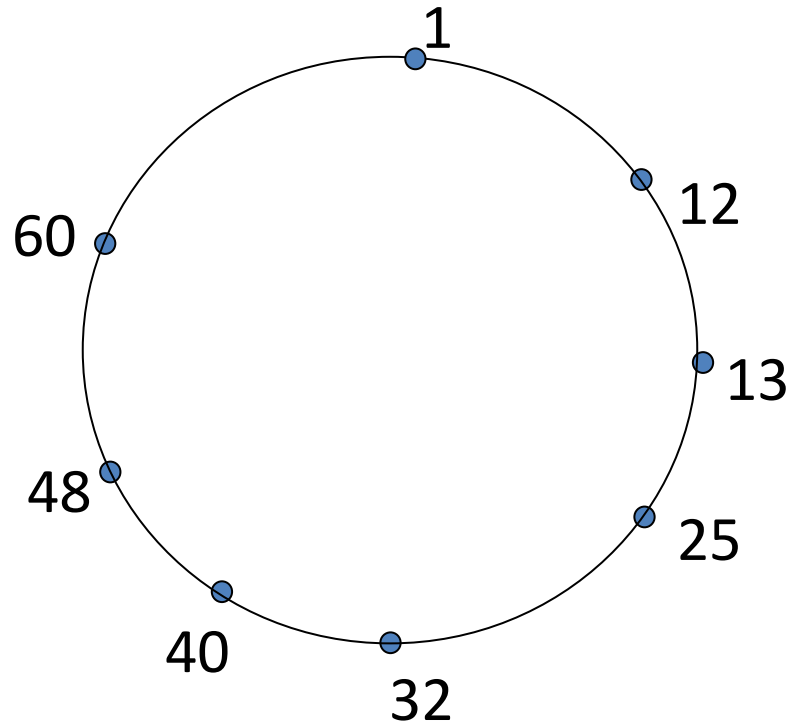
- Assign key-value pair to the peer that has the *closest* ID
- Closest is the *immediate successor* of the key

- Example:

- ID space $\{0,1,2,3,\dots,63\}$
- 8 peers 1, 12, 13, 25, 32, 40, 48, 60
- If key = 51, then assigned to peer 60
- If key = 60, then assigned to peer 60
- If key = 61, then assigned to peer 1

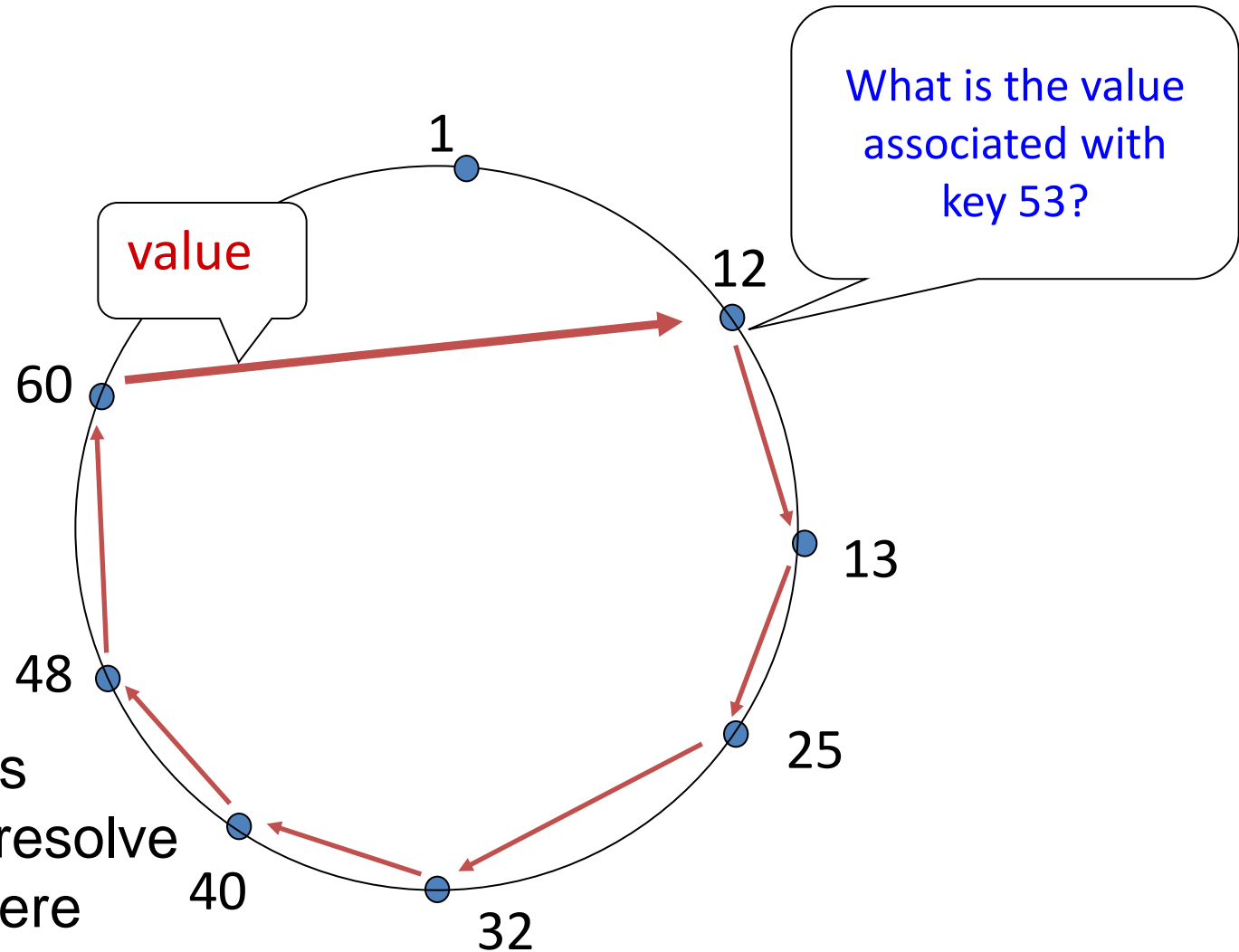
Circular DHT

- Each peer *only* aware of immediate successor and predecessor



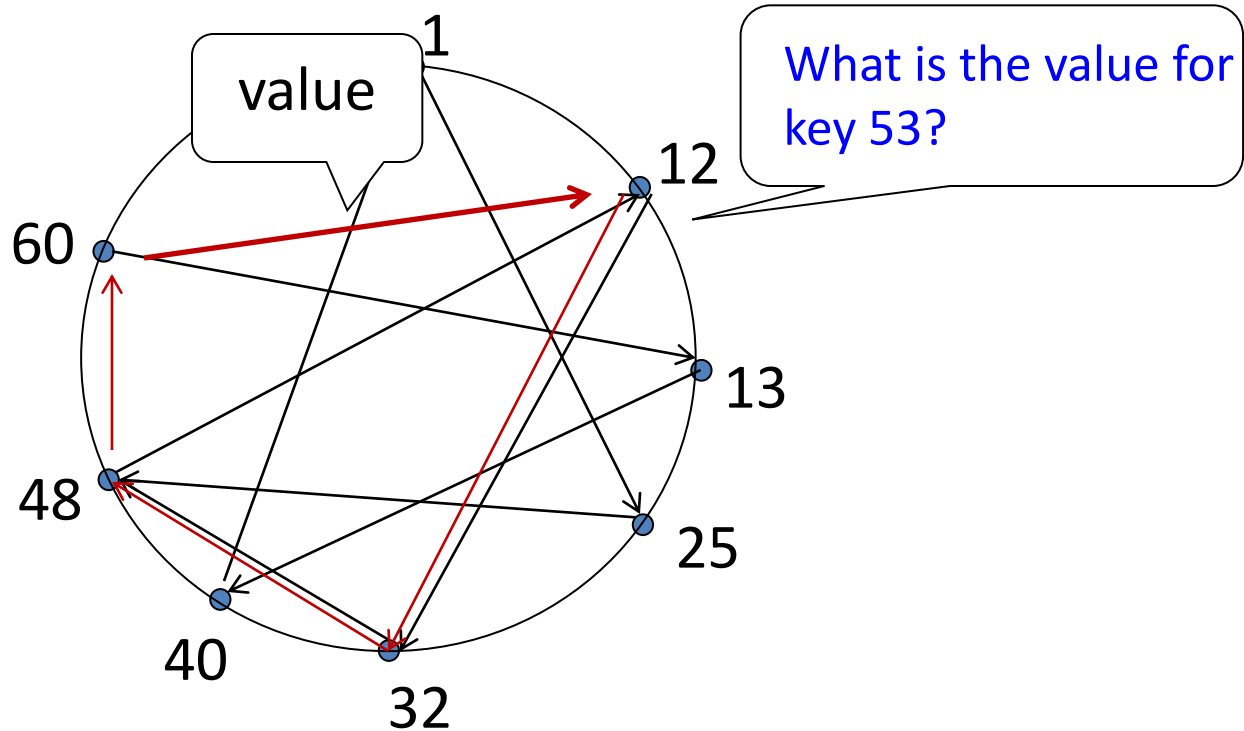
"overlay network"

Resolving a query



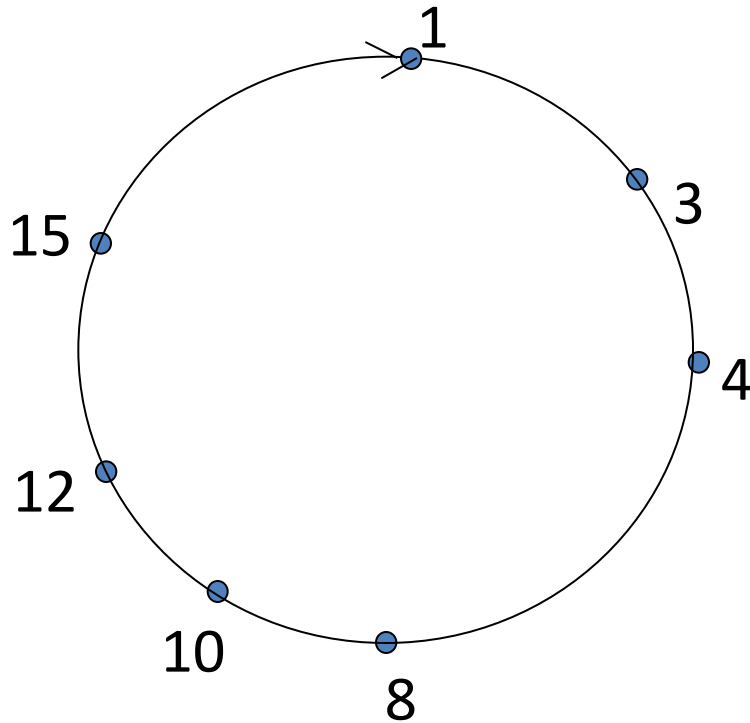
$O(N)$ messages
on average to resolve
query, when there
are N peers

Circular DHT with shortcuts



- Each peer keeps track of IP addresses of predecessor, successor, and short cuts
- Reduced from 6 to 3 messages
- Possible to design shortcuts with $O(\log N)$ neighbors, $O(\log N)$ messages in query

Peer churn



Handling peer churn:

- ❖ Peers may come and go (churn)
- ❖ Each peer knows address of its two successors
- ❖ Each peer periodically pings its two successors to check aliveness
- ❖ If immediate successor leaves, choose next successor as new immediate successor

Example: peer 5 abruptly leaves

- Peer 4 detects peer 5's departure; makes 8 its immediate successor
- 4 asks 8 who its immediate successor is; makes 8's immediate successor its second successor.

Summary

- Peer-to-peer applications
 - Use an overlay network, logical network running on top of existing physical network
 - Can scale with demand better than client-server model
 - Clients share chunks using their upload/download links
 - Finding things:
 - May be centralized (e.g. Napster)
 - Decentralized via a distributed hash table