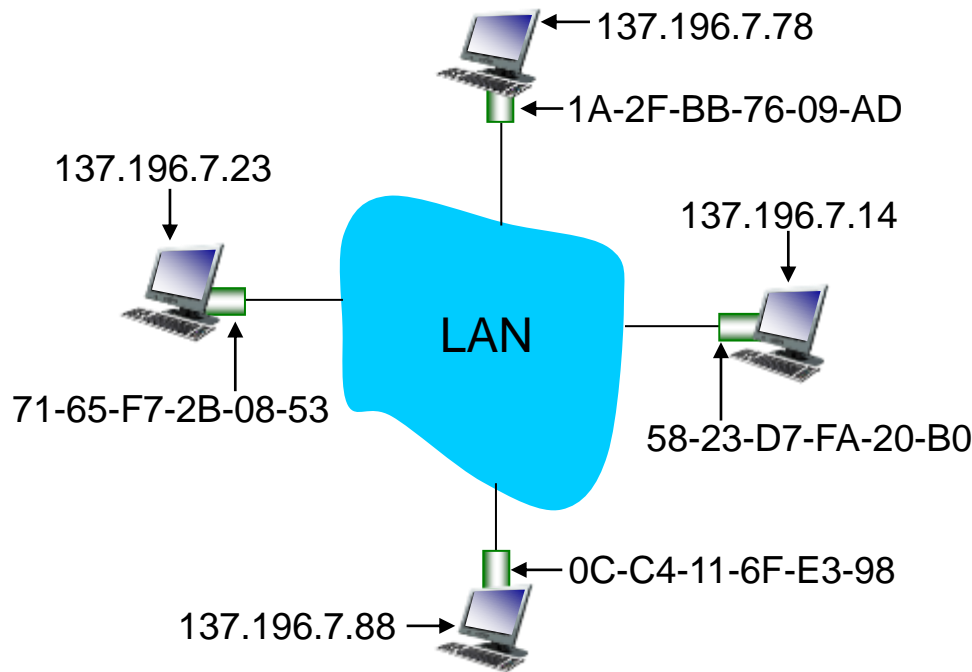


Link-layer addressing, Ethernet, VLANs



Computer Networking: A Top Down Approach

6th edition

Jim Kurose, Keith Ross

Addison-Wesley

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Link layer, LANs: outline

5.1 Introduction,
services

5.2 Error detection,
correction

5.3 Multiple access
protocols

5.4 LANs

- Addressing, ARP
- Ethernet
- Switches
- VLANs

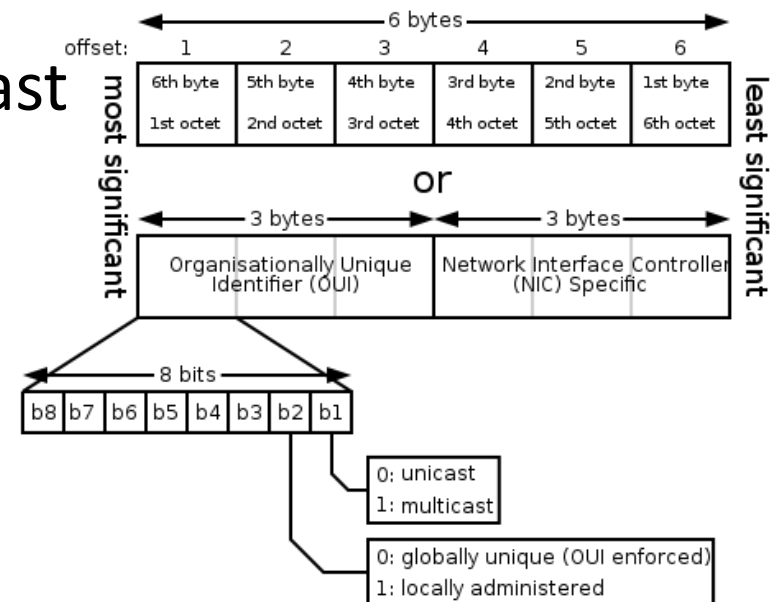
5.5 Link virtualization:
MPLS

5.6 Data center
networking

5.7 A day in the life of
a web request

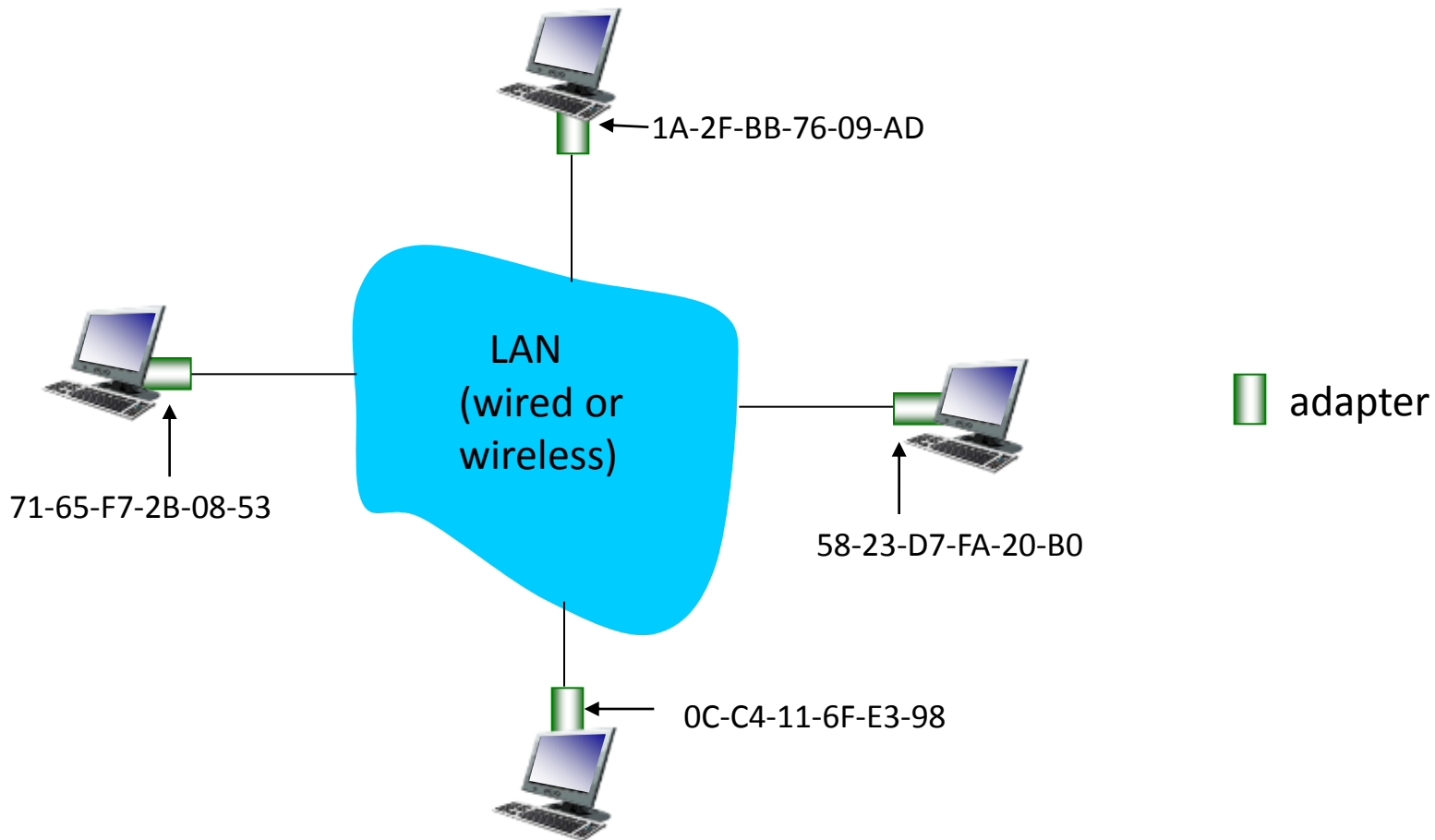
Link-layer addressing

- Media Access Control address (MAC)
 - 48-bit globally unique address
 - 281,474,976,710,656 possible addresses
 - Should last till 2100
 - e.g. 01:23:45:67:89:ab
 - Address of all 1's is broadcast
 - FF:FF:FF:FF:FF:FF



LAN addresses and ARP

Each adapter on LAN has unique *MAC* address
Including routers, NAT boxes, etc.



Address translation

- Problem:
 - How does host send a message to someone on their own network? Their default router?
 - IP address is not the link-level address (e.g. MAC)
- Solution:
 - Host maintains table: IP address -> link address
 - Using the Address Resolution Protocol (ARP)

ARP procedure

- If destination IP in sender's ARP table:
 - Fire off link-layer packet
 - Otherwise send ARP query using broadcast address
- ARP query:
 - IP address you're looking for
 - Your own IP and hardware address
 - Destination responds with hardware address
 - Other hosts can ignore or refresh their ARP tables
- Plug-and-play, no intervention from admin

```
vertanen@katie:/usr/sbin$ arp
Address          HWtype  HWaddress           Flags Mask    Iface
10.33.73.39      ether   c8:2a:14:53:da:e5   C             eth0
lugnut           ether   00:16:3e:c2:af:f0   C             eth0
10.33.73.254     ether   00:12:f2:81:ca:74   C             eth0
10.33.73.120     ether   78:2b:cb:ac:9c:0a   C             eth0
mtcs32.mtech.edu ether   b8:ac:6f:45:56:ef   C             eth0
10.33.73.121     ether   78:2b:cb:ac:9c:1b   C             eth0
vertanen@katie:/usr/sbin$
```

Addressing: routing to another LAN

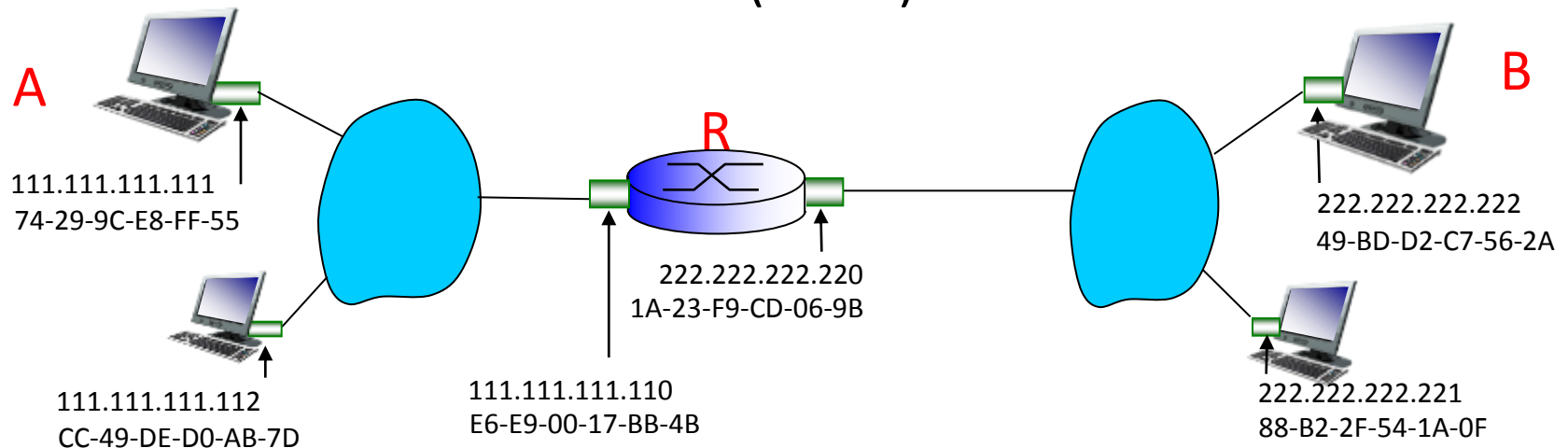
Walkthrough: Send datagram from A to B via R

– Focus on addressing

- At IP (datagram) and MAC layer (frame)

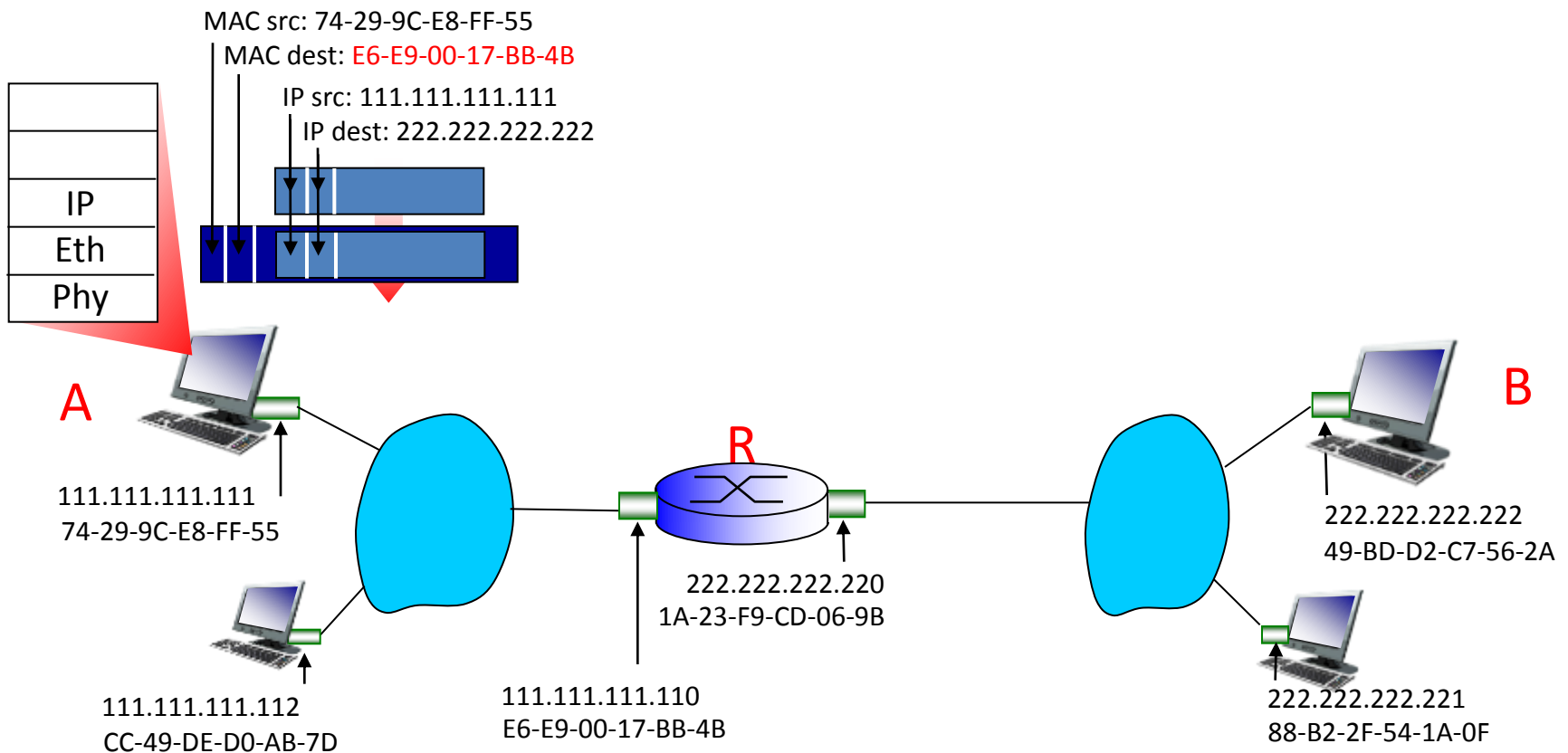
– Assume:

- A knows B's IP address
- A knows IP address of first hop router, R (how?)
- A knows R's MAC address (how?)



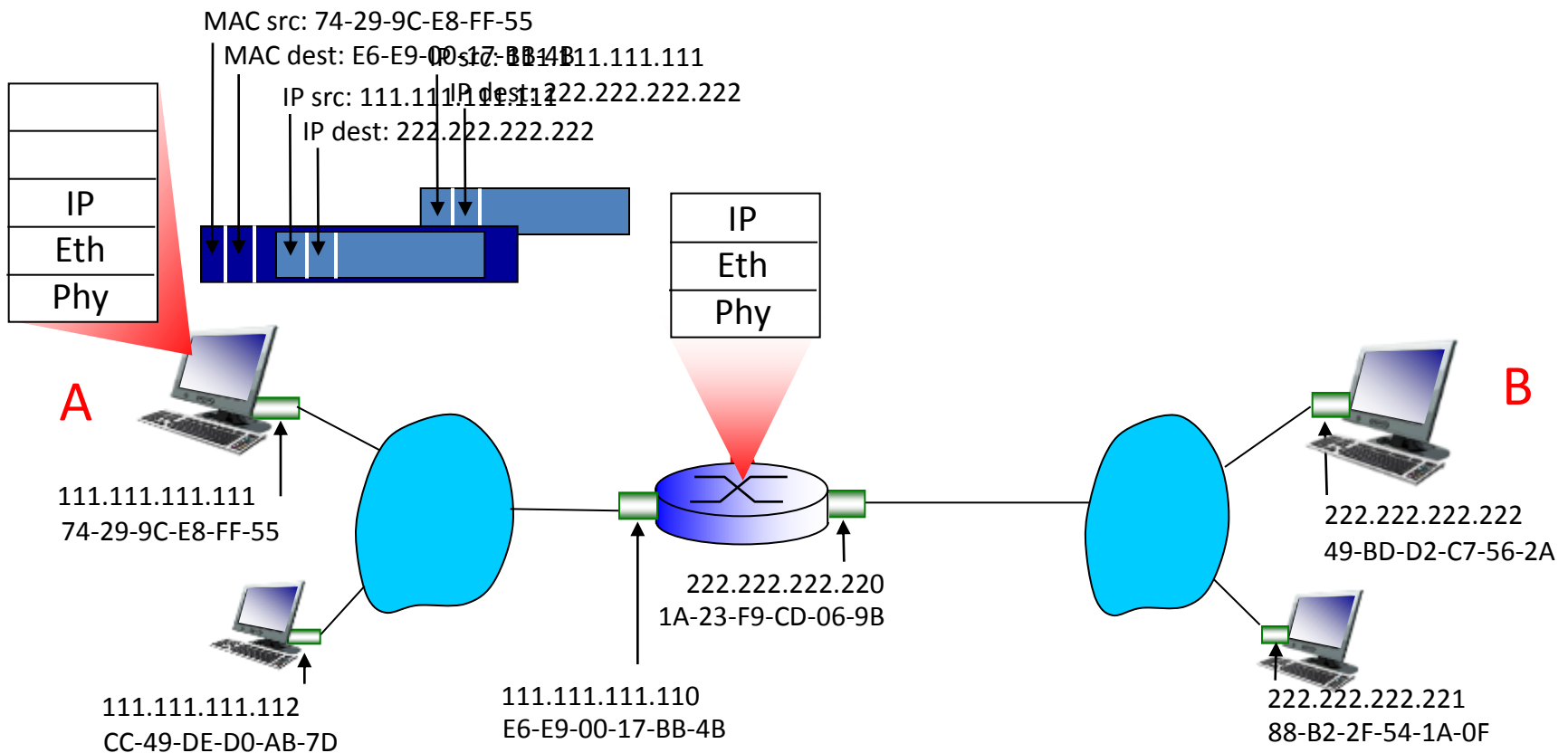
Addressing: routing to another LAN

- ❖ A creates IP datagram with IP source A, destination B
- ❖ A creates link-layer frame with R's MAC address as dest, frame contains A-to-B IP datagram



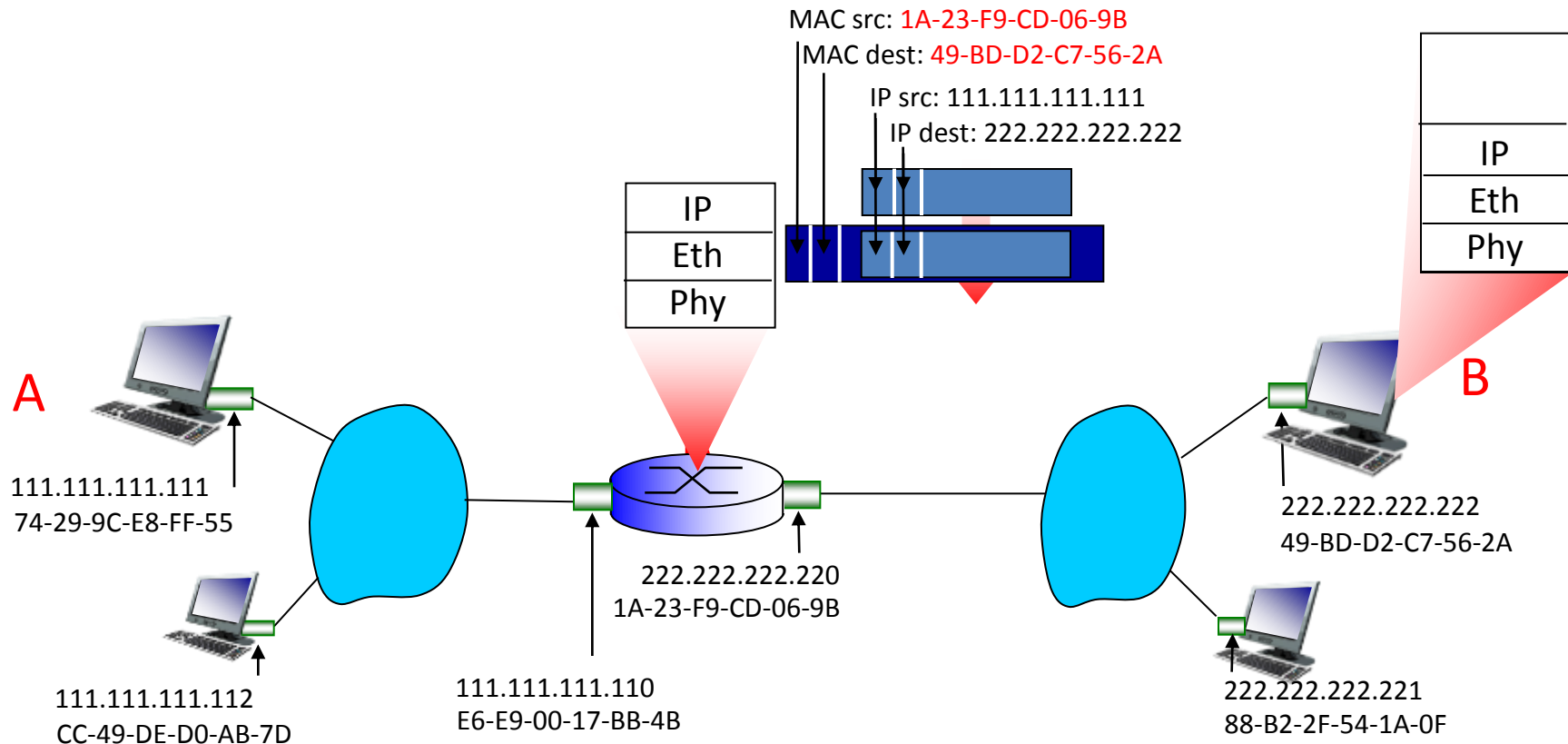
Addressing: routing to another LAN

- ❖ Frame sent from A to R
- ❖ Frame received at R, datagram removed, passed up to IP



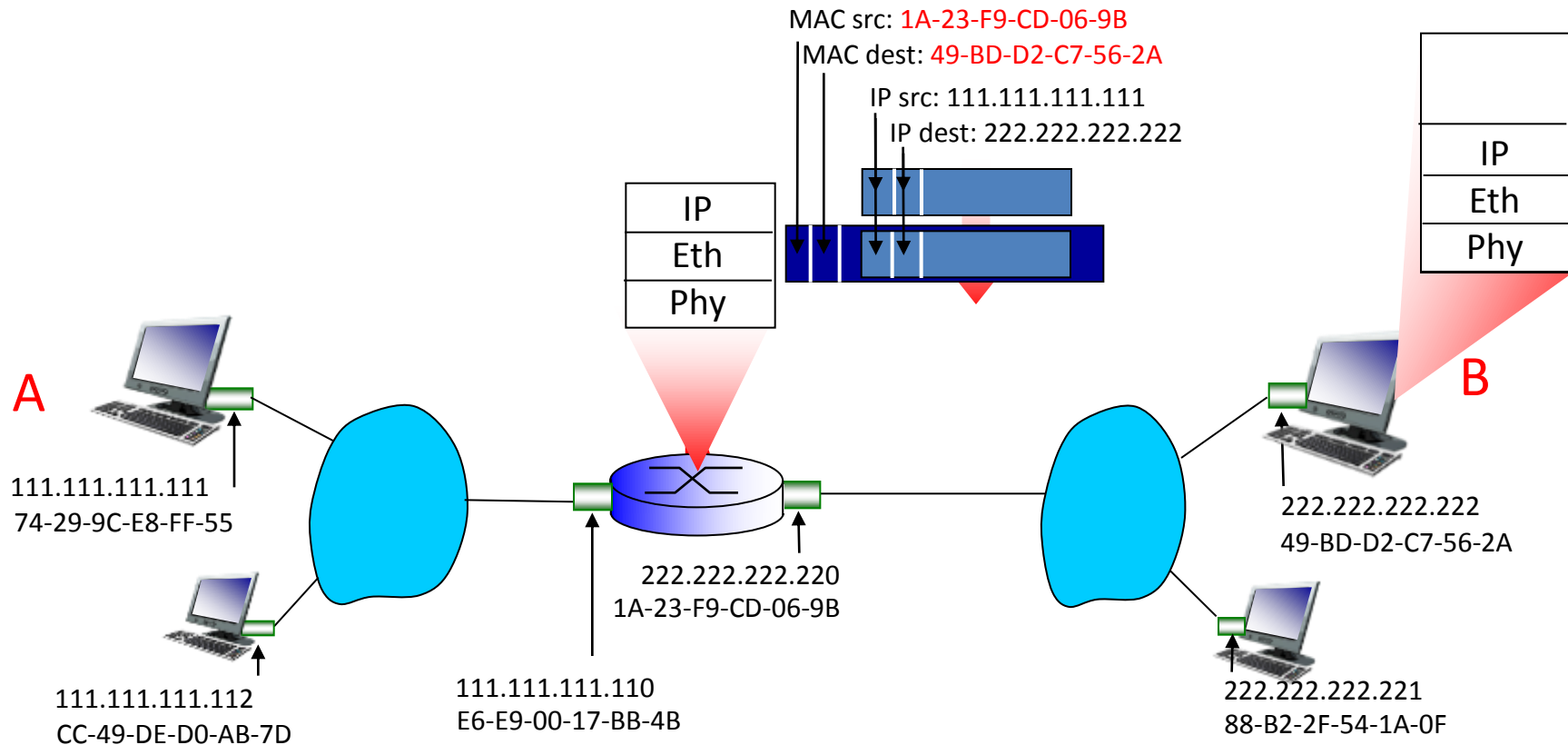
Addressing: routing to another LAN

- ❖ R forwards datagram with IP source A, destination B
- ❖ R creates link-layer frame with B's MAC address as dest, frame contains A-to-B IP datagram



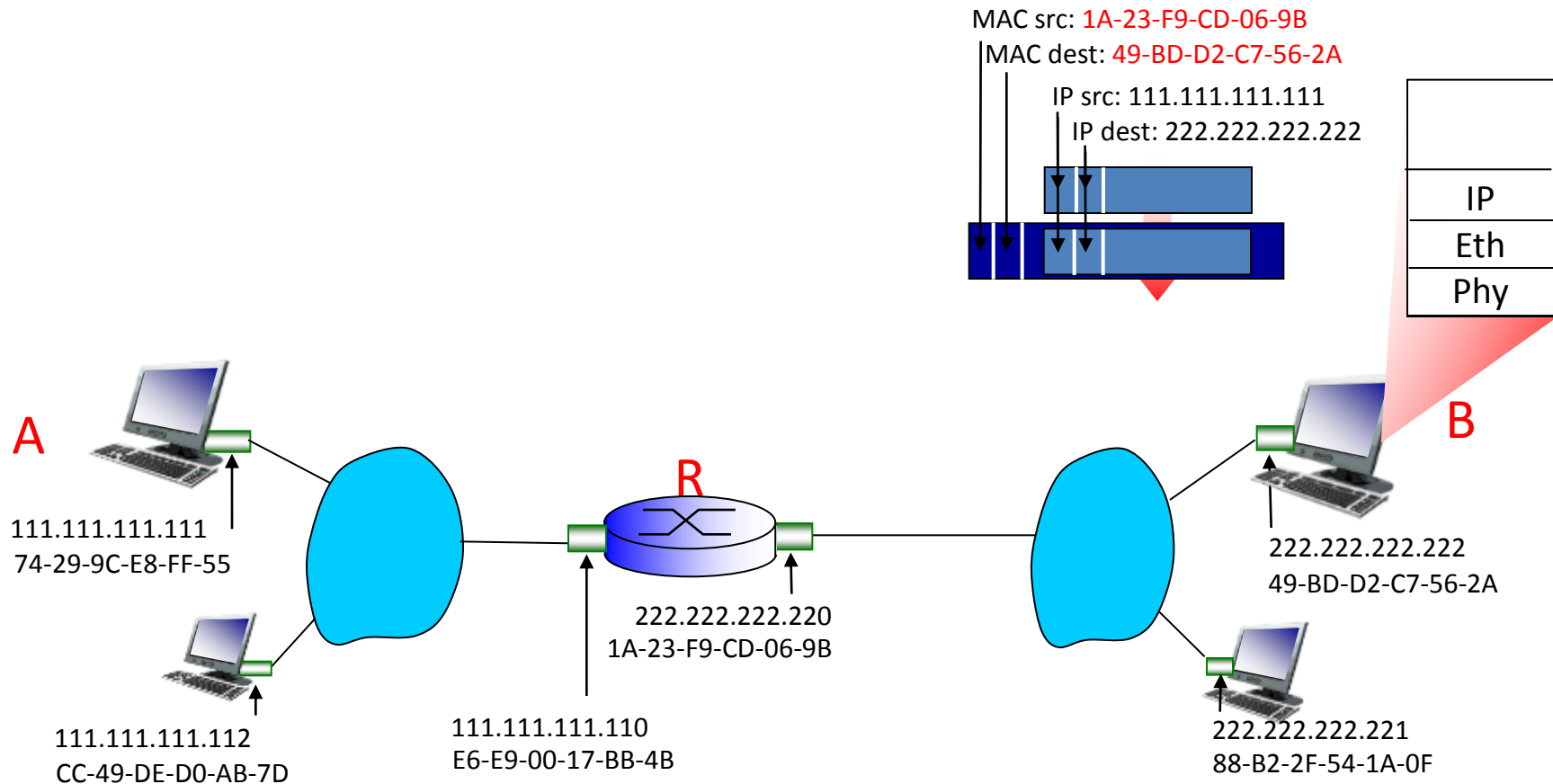
Addressing: routing to another LAN

- ❖ R forwards datagram with IP source A, destination B
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Addressing: routing to another LAN

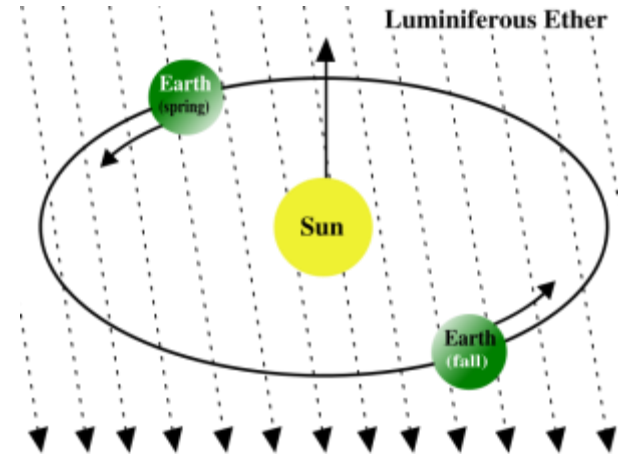
- ❖ R forwards datagram with IP source A, destination B
- ❖ R creates link-layer frame with B's MAC address as dest, frame contains A-to-B IP datagram



Classic Ethernet

- Ethernet

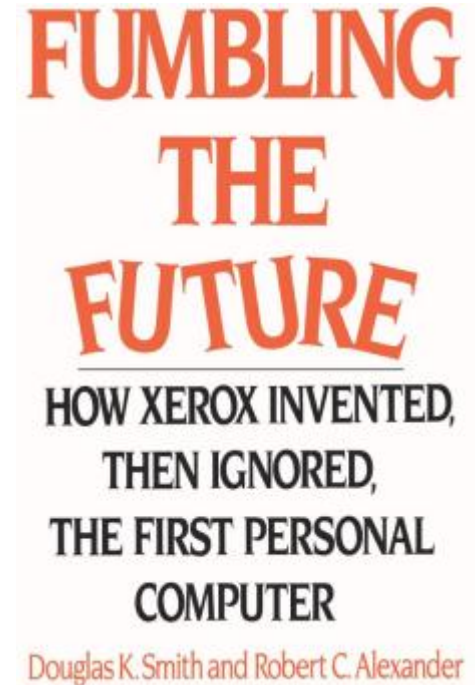
- luminiferous ether through which electromagnetic radiation once thought to propagate
- Carrier Sense, Multiple Access with Collision Detection (CSMA/CD)
- IEEE 802.3



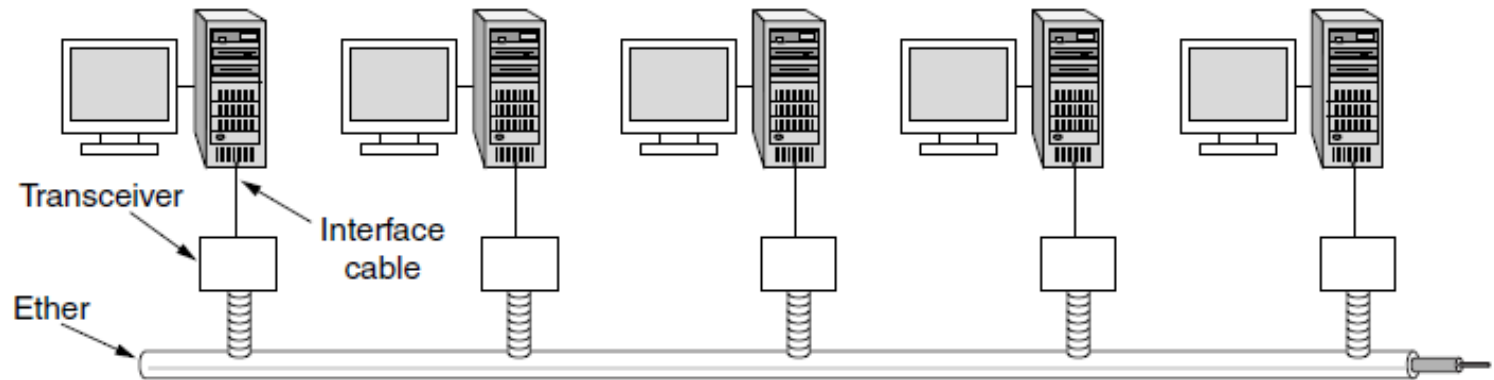
Robert Metcalfe, co-inventor of Ethernet

Classic Ethernet

- Ethernet
 - Xerox Ethernet standardized as IEEE 802.3 in 1983
 - Xerox not interested in commercializing
 - Metcalfe leaves and forms 3Com



Classic Ethernet connectivity



- Shared medium
 - All hosts hear all traffic on cable
 - Hosts tapped the cable
 - 2500m maximum length
 - May include repeaters amplifying signal
 - 10 Mbps bandwidth

Classic Ethernet cabling



Thick Ethernet cable (yellow), 10BASE-5 transceivers, cable tapping tool (orange), 500m maximum length.



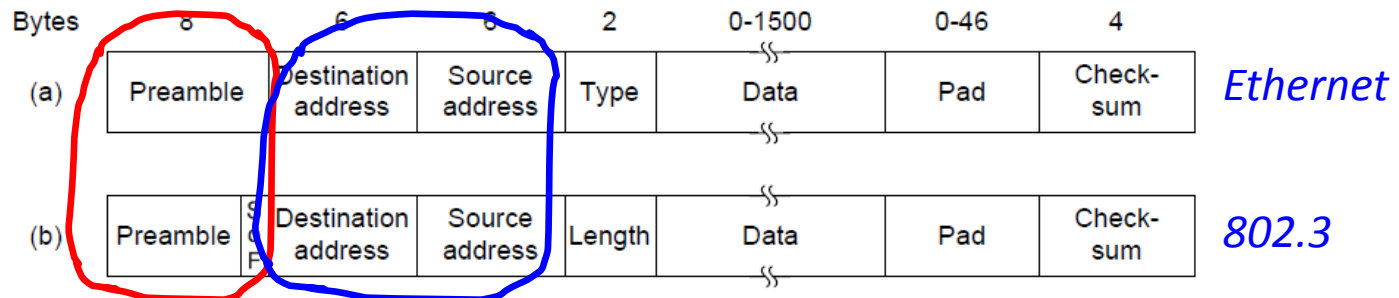
Cable after being "vampire" tapped.



Thin Ethernet cable (10BASE2) with BNC T-connector, 185m maximum length.

Ethernet frame format

- Frame format
 - Manchester encoded
 - Preamble products 10-Mhz square wave
 - Allows clock synch between sender & receiver
 - Pad to at least 64-bytes (collision detection)



Alternating 0's and 1's (except SoF of 11)

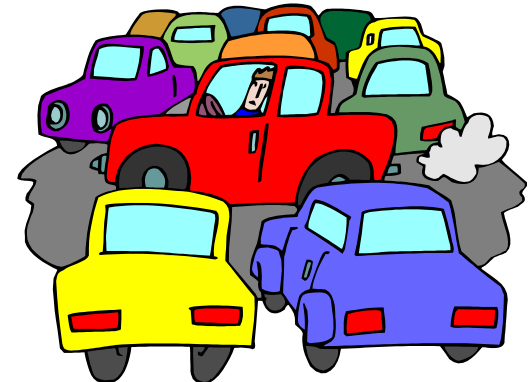
48-bit MAC addresses

Ethernet receivers

- Hosts listens to medium
 - Deliver to host:
 - Any frame with host's MAC address
 - All broadcast frames (all 1's)
 - Multicast frames (if subscribed to)
 - Or all frames if in promiscuous mode

MAC sublayer

- Media Access Control (MAC) sublayer
 - Who goes next on a shared medium
 - Ethernet hosts can sense if medium in use
 - Algorithm for sending data:
 1. Is medium idle? If not, wait.
 2. Start transmitting data, listen for collision.
 3. If collision detected, transmit 32-bit jamming sequence. Stop transmitting and go to backoff procedure.

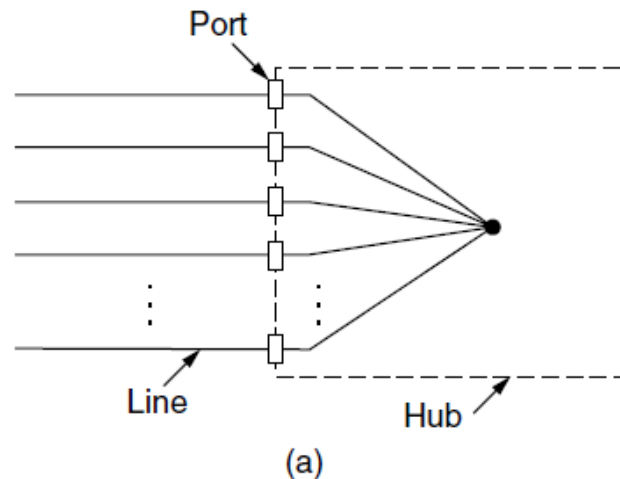


Backoff procedure

- Binary exponential backoff
 - First collision
 - Wait 0-1 timeslots (chosen at random)
 - Second collision
 - Wait 0-3 timeslots
 - In general, i^{th} collision
 - Wait a random number of timeslots between 0 and $2^i - 1$ (max of 1023 slots)
 - Give up after 16 or so retries
 - Timeslot = 51.2 μs

Ethernet hubs

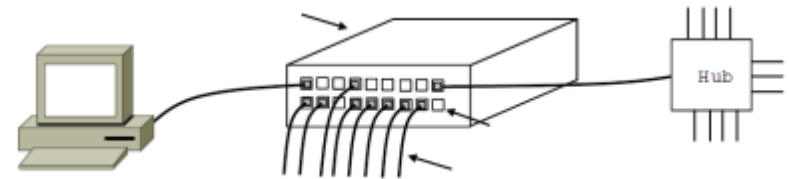
- Long single cable
 - Hard to find breaks or loose connections
- Different wiring pattern
 - Each host wired straight to hub
 - Hub simply connected all wires together
 - Using existing office twisted pair phone lines



Switched Ethernet

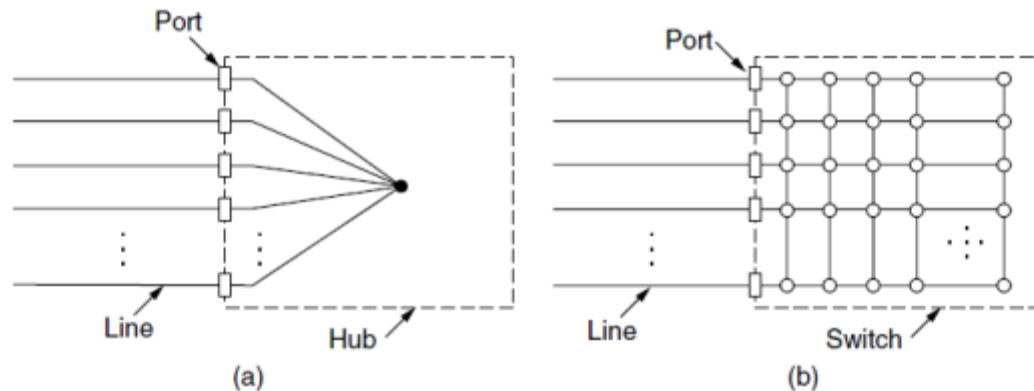
- Hubs

- Made network easier to manage
- But did not address capacity problem



- Switches

- High-speed backplane connecting all ports
- Only output frame to destination port
- Isolates traffic, no collisions, better security



Switch forwarding table

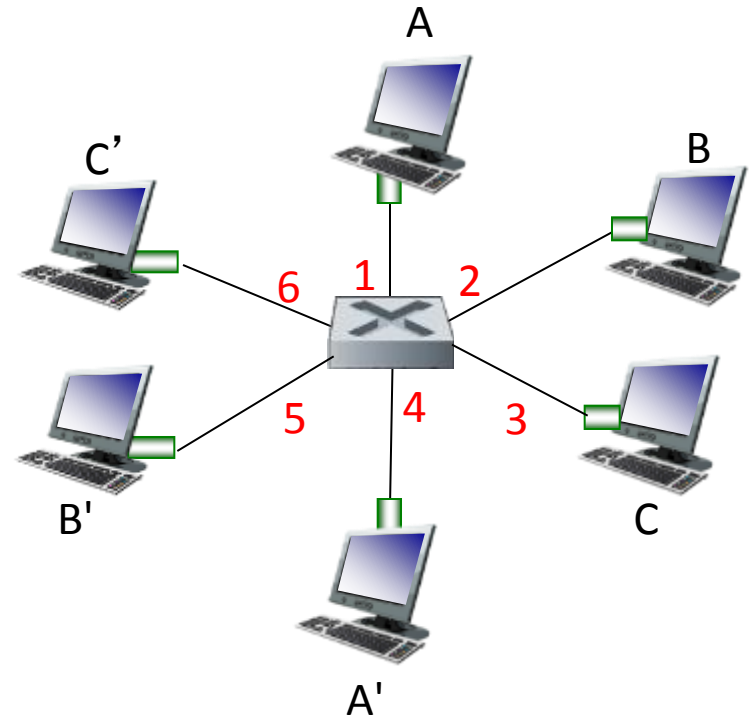
Q: How does switch know A' reachable via interface 4, B' reachable via interface 5?

❖ A: Each switch has a **switch table**, each entry:

- (MAC address of host, interface to reach host, time stamp)
- Looks like a routing table!

Q: How are entries created, maintained in switch table?

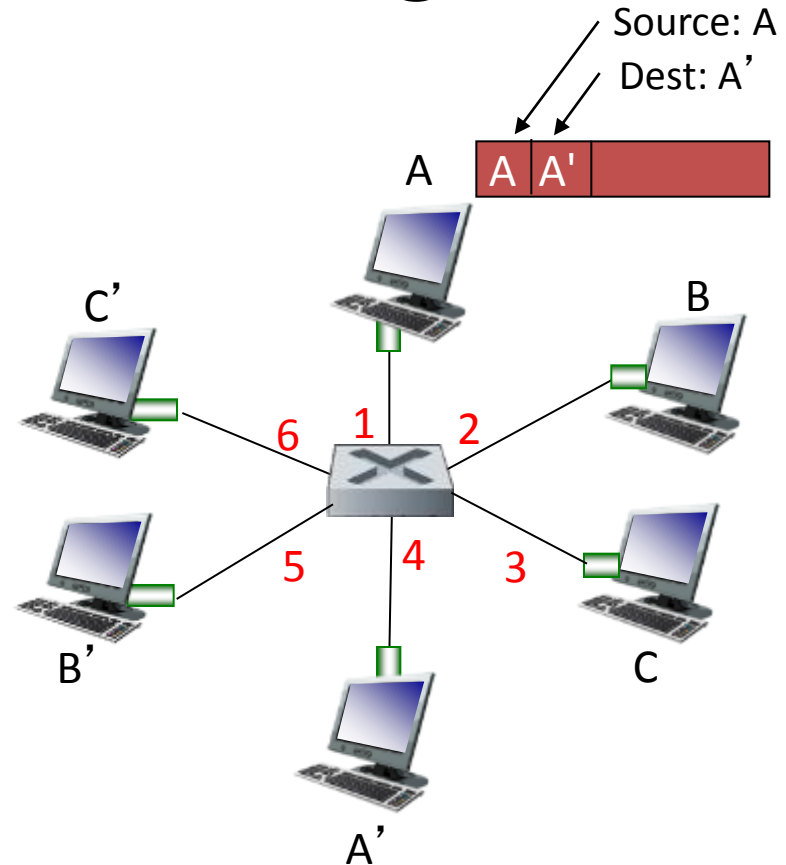
- Something like a routing protocol?



*Switch with six interfaces
(1,2,3,4,5,6)*

Switch: self-learning

- Switch *learns* which hosts can be reached through which interfaces
 - When frame received, switch learns location of sender: incoming LAN segment
 - Records sender/location pair in switch table



MAC addr	interface	TTL
A	1	60

*Switch table
(initially empty)*

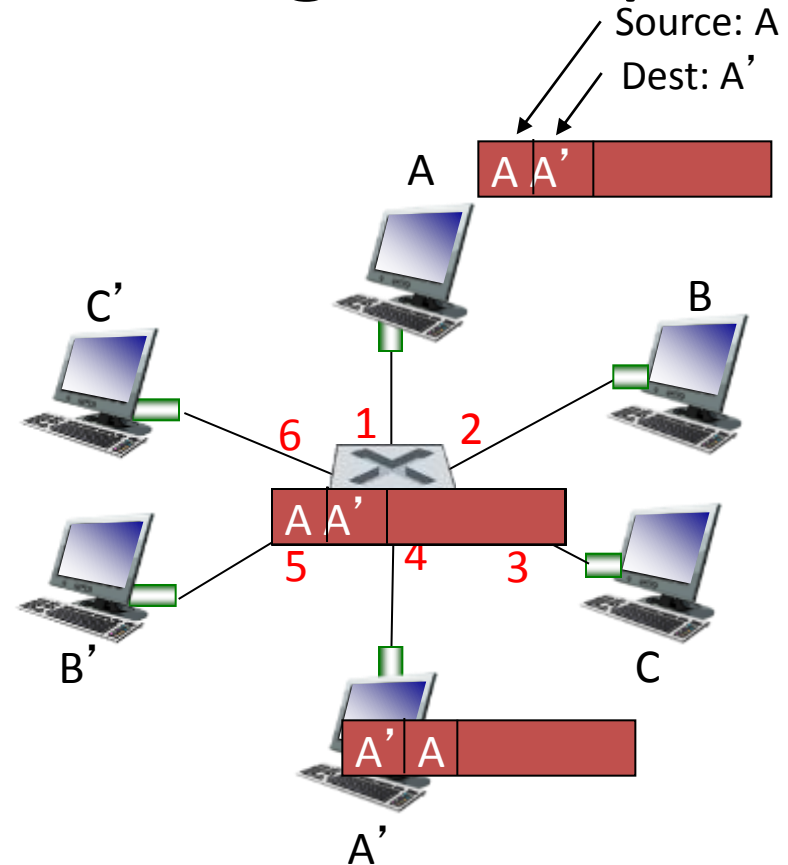
Switch: frame filtering/forwarding

When frame received at switch:

1. Record incoming link, MAC address of sending host
2. Index switch table using MAC destination address
3. if entry found for destination
 then {
 if destination on segment from which frame arrived
 then drop frame
 else forward frame on interface indicated by entry
 }
 else flood /* fwd on all ports except arriving interface */

Self-learning, forwarding: example

- Frame destination, A', location unknown: *flood*
- ❖ Destination A location known: *selectively send on just one link*

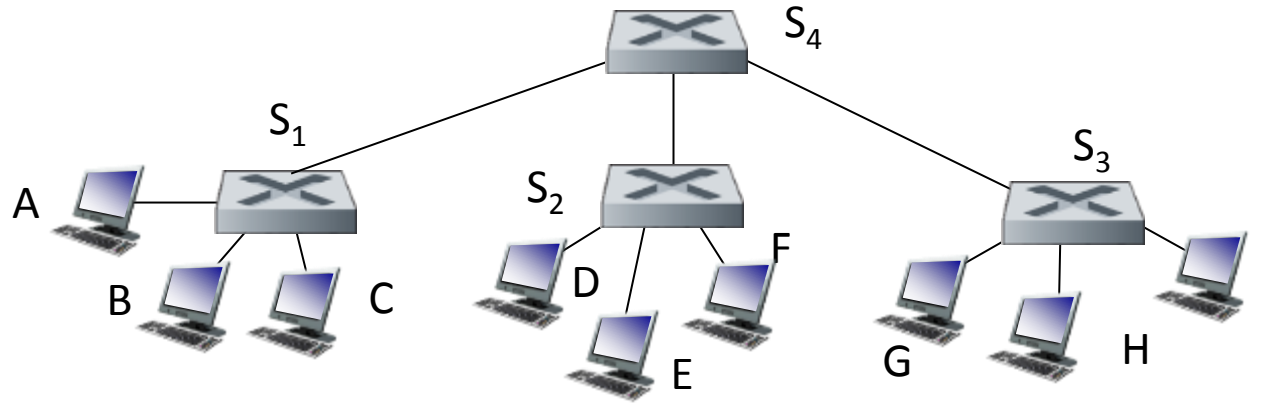


MAC addr	interface	TTL
A	1	60
A'	4	60

*switch table
(initially empty)*

Interconnecting switches

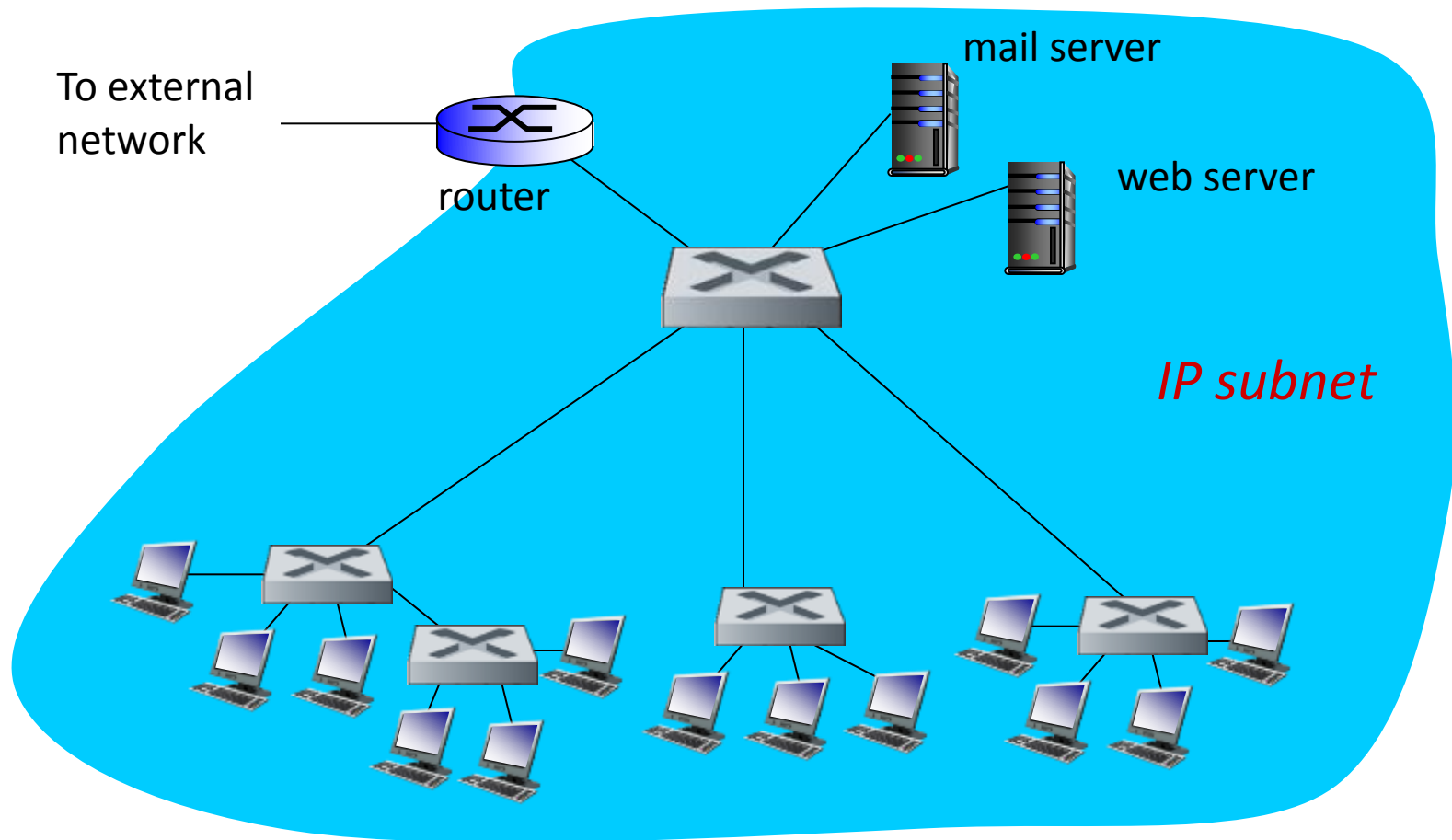
❖ Switches can be connected together



Q: Sending from A to G - how does S₁ know to forward frame destined to F via S₄ and S₃?

❖ A: Self-learning! (works exactly the same as in single-switch case!)

Institutional network



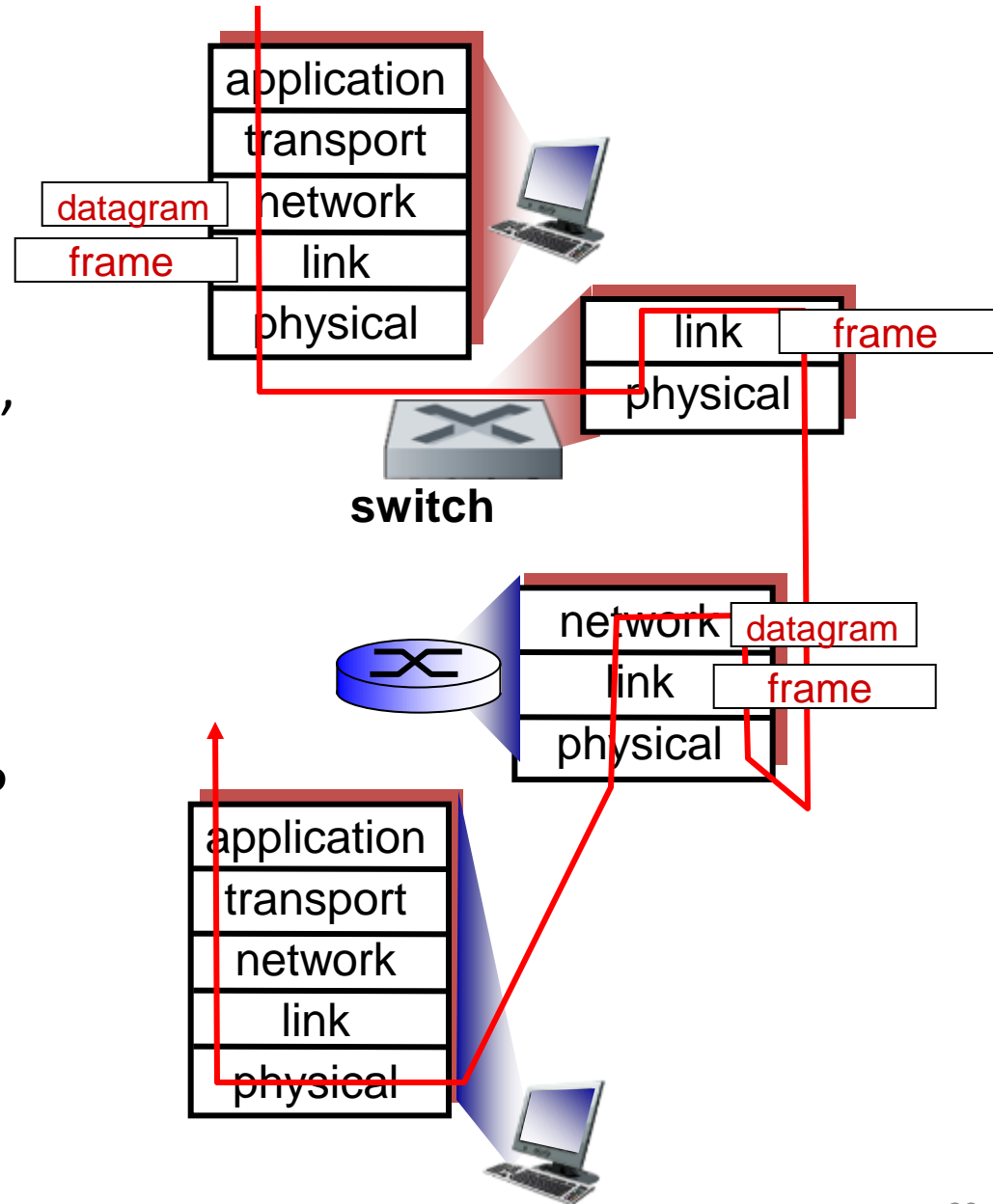
Switches vs. routers

Both are store-and-forward:

- **Routers:** network-layer devices, layer-3 packet switch
- **Switches:** link-layer devices, layer-2 packet switch

Both have forwarding tables:

- **Routers:** compute tables using routing algorithms, IP addresses
- **Switches:** learn forwarding table using flooding, learning, MAC addresses



Fast Ethernet

- Fast Ethernet

- IEEE 802.3u
- Keep all the classic Ethernet frame formats, etc.
- Reduce the bit time from 100 nsec to 10nsec
- 100 Mbps
- No more multidrop cables or vampire taps

Name	Cable	Max. segment	Advantages
100Base-T4	Twisted pair	100 m	Uses category 3 UTP
100Base-TX	Twisted pair	100 m	Full duplex at 100 Mbps (Cat 5 UTP)
100Base-FX	Fiber optics	2000 m	Full duplex at 100 Mbps; long runs

Gigabit Ethernet

- Gigabit Ethernet
 - IEEE 802.3ab
 - 1000 Mbps
 - Unacknowledged datagram service
 - Addition of flow control
 - Unofficial support for jumbo frames
 - Up to 9KB (instead of limit of 1500 bytes)

Name	Cable	Max. segment	Advantages
1000Base-SX	Fiber optics	550 m	Multimode fiber (50, 62.5 microns)
1000Base-LX	Fiber optics	5000 m	Single (10 μ) or multimode (50, 62.5 μ)
1000Base-CX	2 Pairs of STP	25 m	Shielded twisted pair
1000Base-T	4 Pairs of UTP	100 m	Standard category 5 UTP

Even faster

- 10-Gigabit Ethernet
 - 1000x faster than original Ethernet
 - Inside data centers, long haul trunks

Name	Cable	Max. segment	Advantages
10GBase-SR	Fiber optics	Up to 300 m	Multimode fiber (0.85μ)
10GBase-LR	Fiber optics	10 km	Single-mode fiber (1.3μ)
10GBase-ER	Fiber optics	40 km	Single-mode fiber (1.5μ)
10GBase-CX4	4 Pairs of twinax	15 m	Twinaxial copper
10GBase-T	4 Pairs of UTP	100 m	Category 6a UTP

- 40 and 100-Gigabit Ethernet
 - starting to be deployed



[Extreme Networks BlackDiamond X8 Chassis Switch](#)

187.4 lbs, 18" x 30" x 25"

Designed for large, virtualized data centers and clouds, the Extreme Networks BlackDiamond X8 switch provides high density 10GbE and 40GbE ...

\$20,537

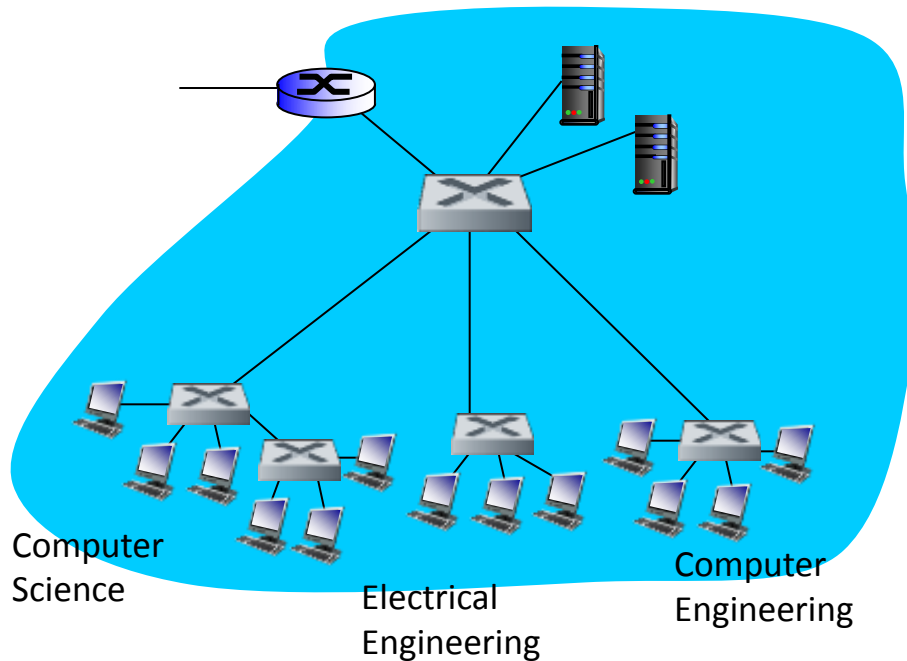
from 3 stores

[Compare prices](#)

Ethernet retrospective

- Why so popular?
 - Easy to administer, no routing or config tables
 - Cheap hardware and wiring
 - Plays nice with TCP/IP
 - Ethernet and IP are connectionless protocols
 - Alternates like ATM were not
 - Periodic speed increases
 - Order of magnitude every few years without throwing away existing infrastructure
 - Borrowed good ideas from other (failed) networking technologies (FDDI, Fiber Channel)

VLANs: motivation



Consider:

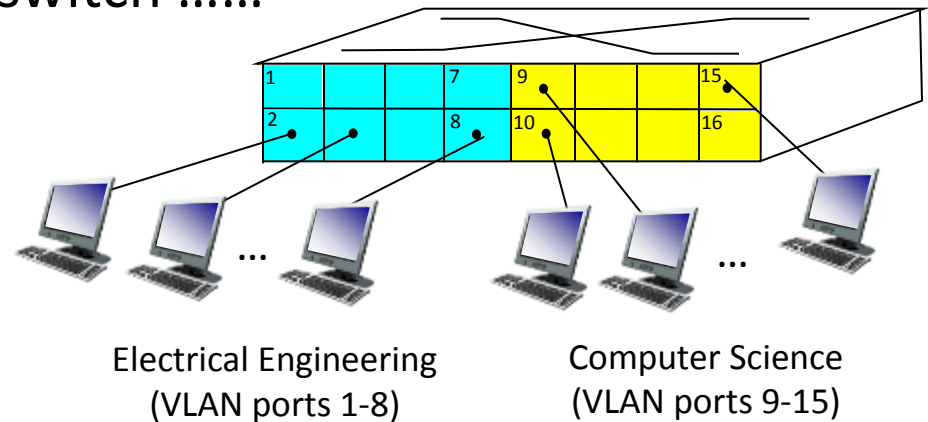
- ❖ CS user moves office to EE, but wants connect to CS switch?
- ❖ Single broadcast domain:
 - All layer-2 broadcast traffic (ARP, DHCP, unknown location of destination MAC address) must cross entire LAN
 - Security/privacy, efficiency issues

VLANs

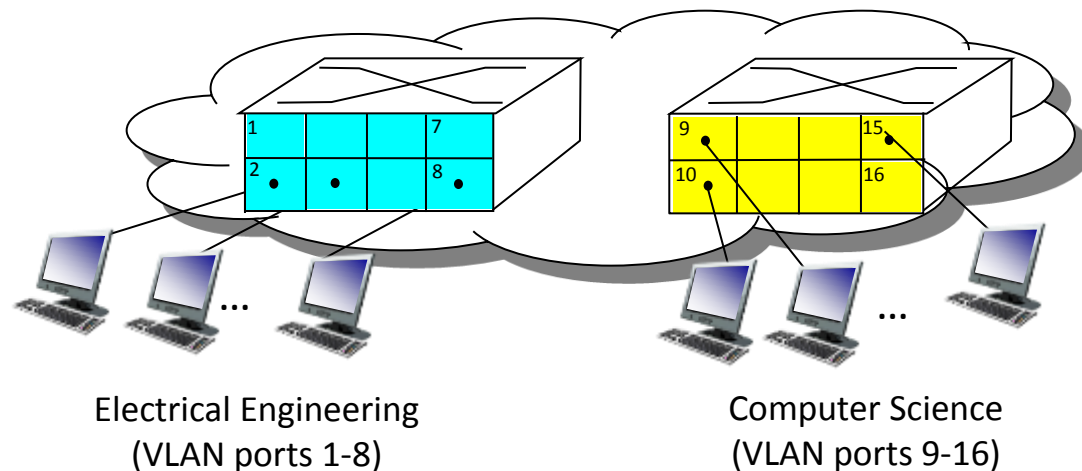
Virtual Local Area Network

Switches supporting VLAN capabilities can be configured to define multiple **virtual** LANS over single physical LAN infrastructure.

Port-based VLAN: Switch ports grouped (by switch management software) so that **single** physical switch



... operates as **multiple** virtual switches



Port-based VLAN

❖ *Traffic isolation:*

Frames to/from ports 1-8
can *only* reach ports 1-8

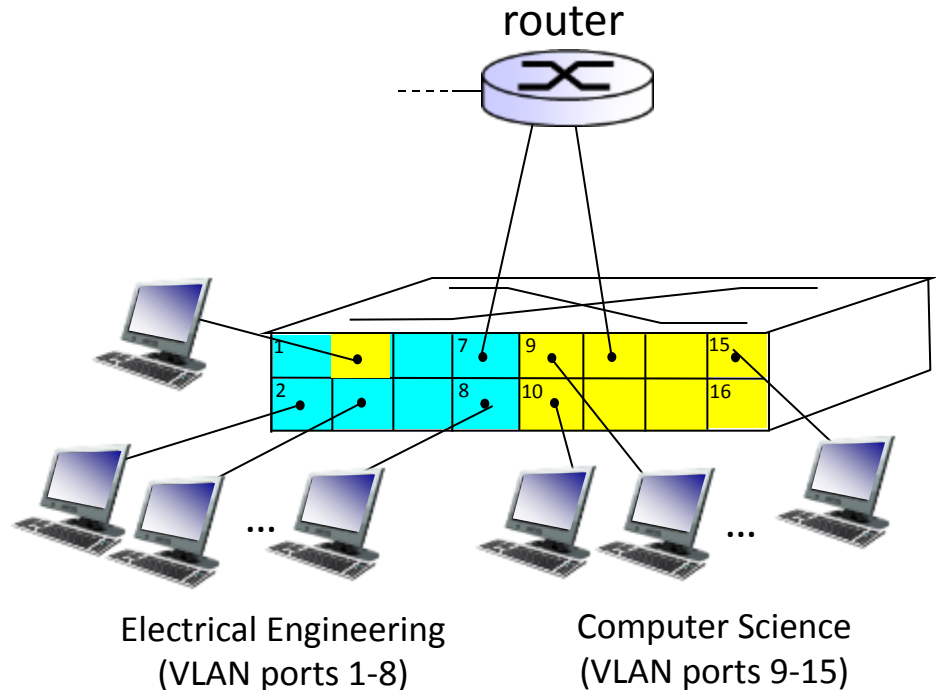
- Can also define VLAN based on MAC addresses of endpoints, rather than switch port

❖ *Dynamic membership:*

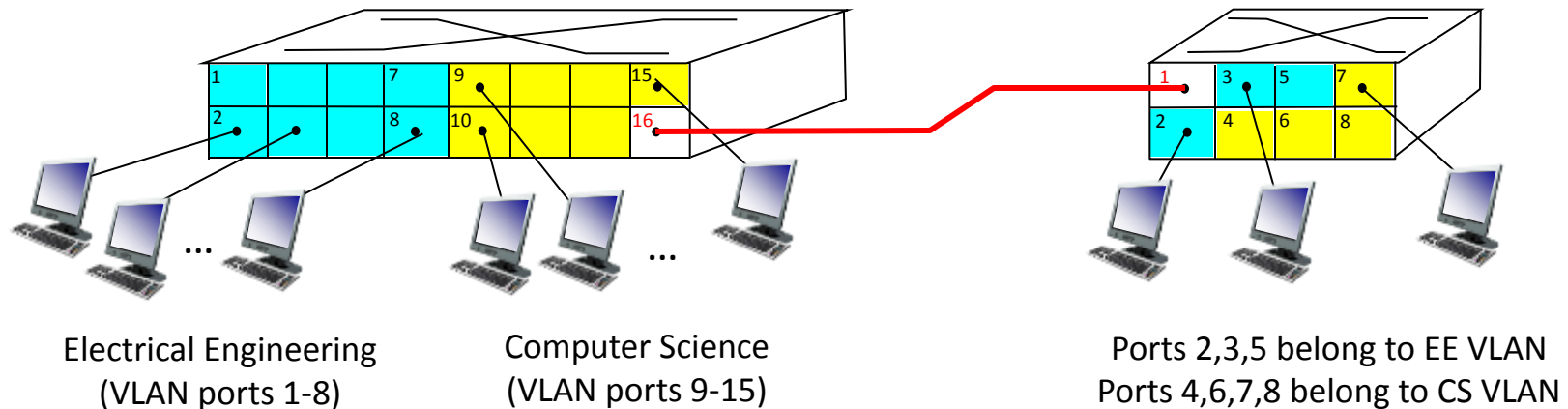
Ports can be dynamically assigned among VLANs

❖ *Forwarding between VLANs:* done via routing

- Just as with separate switches
- In practice vendors sell combined switches plus routers

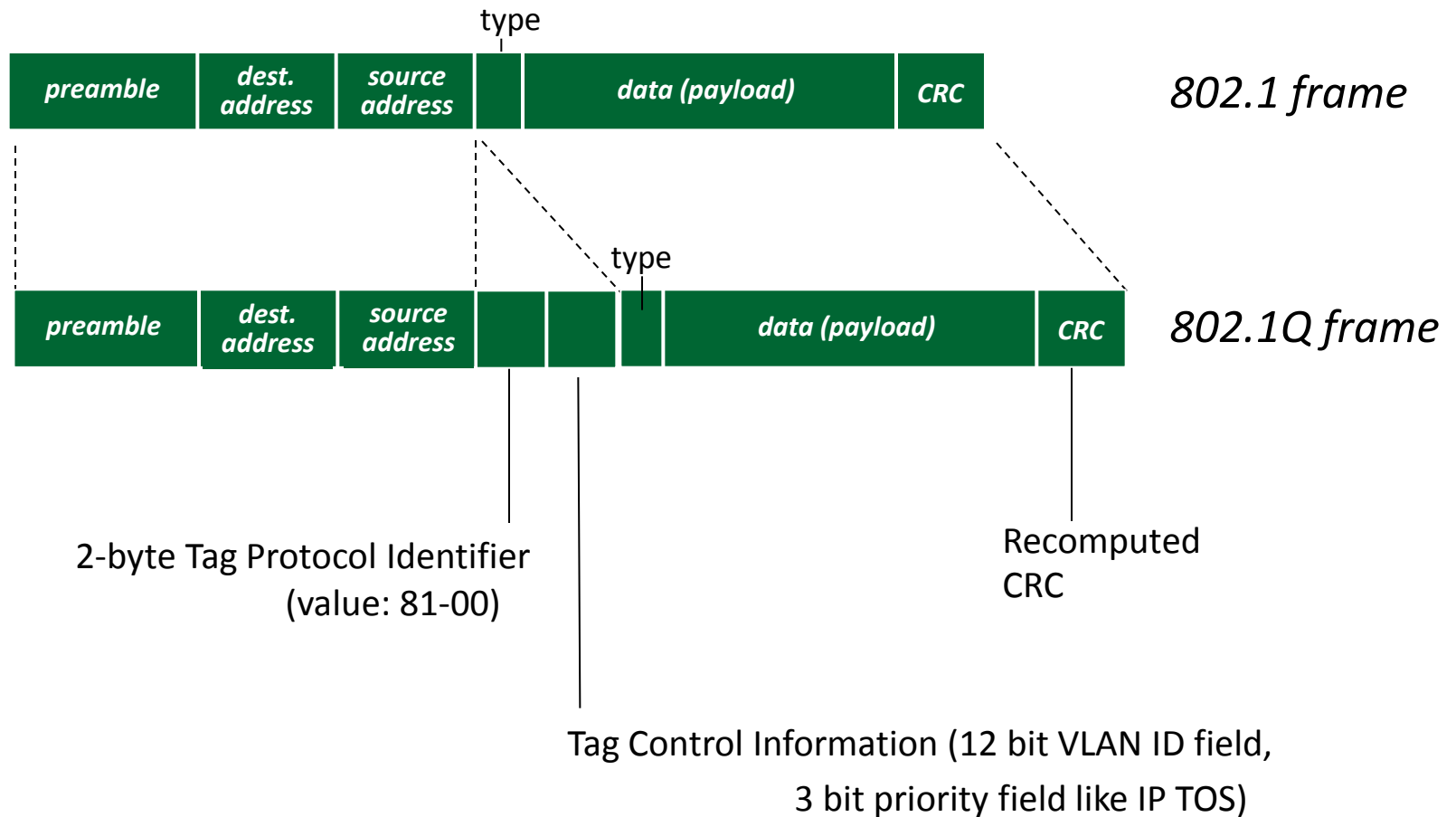


VLANs spanning multiple switches



- **Trunk port:** Carries frames between VLANs defined over multiple physical switches
 - Frames forwarded within VLAN between switches can't be vanilla 802.1 frames, **must carry VLAN ID info**
 - **802.1q protocol** adds/removed additional header fields for frames forwarded between trunk ports

802.1Q VLAN frame format



Summary

- Address Reservation Protocol (ARP)
 - Mapping between link-layer addresses (MAC) and network-layer addresses (IP)
 - Cached table in operating system
 - Broadcast queries for IP destinations with unknown MAC
- Wired Ethernet
 - Long history and widely adopted Hubs vs. switches vs. routers
 - Order of magnitude bit rate increase every few years
 - Careful attention to backwards compatibility
- VLANs
 - Allows virtual segregation of hosts into isolated groups